

Week 9 Tutorial  
Hoare Logic

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## 1 Hoare Notation — Warm-up Exercises

These questions are meant to foster the feeling that Hoare notation is a compact way of saying things that you want to say when discussing a program. Each of the following parts consists of a statement in Hoare Logic — each is a simple assertion about a piece of code. In each case you should just say whether the statement is true or false.

All that is wanted for this group of problems is just a true-or-false answer based on your intuitions. Can your intuitions be programmed, though?

(The program variables in this and later questions are typed integer.)

- a)  $\{j = a\} j := j + 1 \{a = j + 1\}$
- b)  $\{i = j\} i := j + i \{i > j\}$
- c)  $\{j = a + b\} i := b; j := a \{j = 2 * a\}$
- d)  $\{i > j\} j := i + 1; i := j + 1 \{i > j\}$
- e)  $\{i \neq j\} \text{if } i > j \text{ then } m := i - j \text{ else } m := j - i \{m > 0\}$
- f)  $\{i = 3 * j\} \text{if } i > j \text{ then } m := i - j \text{ else } m := j - i \{m - 2 * j = 0\}$
- g)  $\{x = b\} \text{while } x > a \text{ do } x := x - 1 \{b = a\}$

## 2 Hoare Proof Rules I — Assignment Statements

Next up, we will be focusing on the natural deduction rules for precondition strengthening, postcondition weakening and especially the use of the assignment axiom.

- a) Prove  $\{i = 5\} \mathbf{a:=i+2} \{(a = 7) \wedge (i = 5)\}$
- b) Prove  $\{i = 5\} \mathbf{a:=i+2} \{a = 7\}$
- c) Prove  $\{i = 5\} \mathbf{a:=i+2} \{(a = 7) \wedge (i > 0)\}$
- d) Prove  $\{(i = 5) \wedge (a = 3)\} \mathbf{a:=i+2} \{a = 7\}$
- e) Prove  $\{a = 7\} \mathbf{i:=i+2} \{a = 7\}$
- f) Prove  $\{i = a - 1\} \mathbf{i:=i+2} \{i = a + 1\}$
- g) Prove  $\{i = 5\} \mathbf{i:=i+2} \{i > 0\}$
- h) Prove  $\{True\} \mathbf{a:=i+2} \{a = i + 2\}$

## 3 Hoare Proof Rules for Control Structures

Remember the way proofs of larger program fragments are meant to be constructed. If your immediate goal is to prove some property of a conditional statement, or a loop or a sequence of statements, then you use the corresponding Hoare Rule to generate appropriate subgoals.

### 3.1 Sequencing

- a) Prove  $\{a > b\} \mathbf{m:=1; n:=a-b} \{m * n > 0\}$
- b) Prove  $\{s = 2^i\} \mathbf{i:=i+1; s:=s*2} \{s = 2^i\}$

### 3.2 Conditionals

- c) Prove  $\{True\} \mathbf{if\ i<j\ then\ min:=i\ else\ min:=j} \{(min \leq i) \wedge (min \leq j)\}$
- d) Prove  $\{i > 0 \wedge j > 0\} \mathbf{if\ i<j\ then\ min:=i\ else\ min:=j} \{min > 0\}$

### 3.3 Loops

- e) Prove  $\{s = 2^i\} \mathbf{while\ i<n\ do\ [i:=i+1; s:=s*2]} \{s = 2^i\}$

## 4 Want More?

You can get a bit more practice with the warm-up questions. If they were statements that were intuitively true then you should be able to prove them. Otherwise, beat them into shape and prove the better version.

## 5 Appendix: Hoare Logic Rules

- Precondition Strengthening:

$$\frac{\{P_w\} S \{Q\} \quad P_s \rightarrow P_w}{\{P_s\} S \{Q\}}$$

- Postcondition Weakening:

$$\frac{\{P\} S \{Q_s\} \quad Q_s \rightarrow Q_w}{\{P\} S \{Q_w\}}$$

- Assignment:

$$\{Q(e)\} x := e \{Q(x)\}$$

- Sequence:

$$\frac{\{P\} S_1 \{Q\} \quad \{Q\} S_2 \{R\}}{\{P\} S_1; S_2 \{R\}}$$

- Conditional:

$$\frac{\{P \wedge b\} S_1 \{Q\} \quad \{P \wedge \neg b\} S_2 \{Q\}}{\{P\} \text{ if } b \text{ then } S_1 \text{ else } S_2 \{Q\}}$$

- While Loop:

$$\frac{\{P \wedge b\} S \{P\}}{\{P\} \text{ while } b \text{ do } S \{P \wedge \neg b\}}$$