Introduction to Concurrency

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References for this chapter

[Ben-Ari06]
M. Ben-Ari

*Principles of Concurrent and Distributed Programming*
Introduction to Concurrency

Forms of concurrency

What is concurrency?

Working definitions:

• Literally ‘concurrent’ means:

  Adj.: Running together in space, as parallel lines; going on side by side, as proceedings; occurring together, as events or circumstances; existing or arising together; conjoint, associated [Oxford's English Dictionary]
Forms of concurrency

What is concurrency?

Working definitions:

- Literally ‘concurrent’ means:
  
  Adj.: Running together in space, as parallel lines; going on side by side, as proceedings; occurring together, as events or circumstances; existing or arising together; conjoint, associated [Oxford's English Dictionary]

- Technically ‘concurrent’ is usually defined negatively as:
  
  If there is no observer who can identify two events as being in strict temporal sequence (i.e. one event has fully terminated before the other one started) then these two events are considered concurrent.
**Forms of concurrency**

**Why do we need/have concurrency?**

- Physics, engineering, electronics, biology, ... 
  
  basically every real world system is concurrent!

- Sequential processing is suggested by most core computer architectures
  
  ... yet (almost) all current processor architectures have concurrent elements
  
  ... and most computer systems are part of a concurrent network.

- Strict sequential processing is suggested by widely used programming languages.

  Sequential programming delivers some fundamental components for concurrent programming

  but we need to add a number of further crucial concepts
Why would a computer scientist consider concurrency?

- ... to be able to connect computer systems with the real world
- ... to be able to employ / design concurrent parts of computer architectures
- ... to construct complex software packages (operating systems, compilers, databases, ...)
- ... to understand when sequential and/or concurrent programming is required
  ... or: to understand when sequential or concurrent programming can be chosen freely
- ... to enhance the reactivity of a system
- ... to enhance the performance of a system
- ... to be able to design embedded systems
- ...
Forms of concurrency

A computer scientist’s view on concurrency

- Overlapped I/O and computation
  - Employ interrupt programming to handle I/O

- Multi-programming
  - Allow multiple independent programs to be executed on one CPU

- Multi-tasking
  - Allow multiple interacting processes to be executed on one CPU

- Multi-processor systems
  - Add physical/real concurrency

- Parallel Machines & distributed operating systems
  - Add (non-deterministic) communication channels

- General network architectures
  - Allow for any form of communicating, distributed entities
Forms of concurrency

A computer scientist’s view on concurrency

Terminology for physically concurrent machines architectures:

- **SISD**
  [single instruction, single data]
  - Sequential processors
- **SIMD**
  [single instruction, multiple data]
  - Vector processors
- **MISD**
  [multiple instruction, single data]
  - Pipelined processors
- **MIMD**
  [multiple instruction, multiple data]
  - Multi-processors or computer networks
Multiple physical, coupled, dynamical systems form the actual environment and/or task at hand.

In order to model and control such a system, its inherent concurrency needs to be considered.

Multiple less powerful processors are often preferred over a single high-performance CPU.

The system design of usually strictly based on the structure of the given physical system.
Forms of concurrency

Does concurrency lead to chaos?

Concurrency often leads to the following features / issues / problems:

- **non-deterministic** phenomena
- **non-observable** system states
- results may depend on more than just the input parameters and states at start time (timing, throughput, load, available resources, signals … throughout the execution)
- **non-reproducible** debugging?
Forms of concurrency

Does concurrency lead to chaos?

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- **non-deterministic** phenomena
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- results may depend on more than just the input parameters and states at start time (timing, throughput, load, available resources, signals ... *throughout* the execution)
- **non-reproducible** debugging?

Meaningful employment of concurrent systems features:

- non-determinism employed where the underlying system is non-deterministic
- non-determinism employed where the actual execution sequence is meaningless
- **synchronization** employed where adequate ... but only there

 ✨ Control & monitor where required (and do it right), but not more ...
Concurrency on different abstraction levels/perspectives

- **Networks**
  - Large scale, high bandwidth interconnected nodes ("supercomputers")
  - Networked computing nodes
  - Standalone computing nodes – including local buses & interfaces sub-systems
  - Operating systems (& distributed operating systems)

- **Implicit concurrency**

- **Explicit concurrent programming (message passing and synchronization)**

- **Assembler level concurrent programming**
  - Individual concurrent units inside one CPU
  - Individual electronic circuits
  - ...
Models and Terminology

The concurrent programming abstraction

1. What appears sequential on a higher abstraction level, is usually concurrent at a lower abstraction level:
   e.g. Concurrent operating system or hardware components, which might not be visible at a higher programming level

2. What appears concurrent on a higher abstraction level, might be sequential at a lower abstraction level:
   e.g. Multi-processing system, which are executed on a single, sequential computing node
Models and Terminology

The concurrent programming abstraction

• ‘concurrent’ is technically defined negatively as:

  If there is no observer who can identify two events as being in
  strict temporal sequence (i.e. one event has fully terminated before the
  other one starts up), then these two events are considered concurrent.

• ‘concurrent’ in the context of programming and logic:

  “Concurrent programming abstraction is the study of
  interleaved execution sequences of the atomic
  instructions of sequential processes.”

  (Ben-Ari)
Models and Terminology

The concurrent programming abstraction

Concurrent program ::= Multiple sequential programs (processes or threads) which are executed concurrently.

P.S. it is generally assumed that concurrent execution means that there is one execution unit (processor) per sequential program

• even though this is usually not technically correct, it is still an often valid, conservative assumption in the context of concurrent programming.
Models and Terminology

The concurrent programming abstraction

No interaction between concurrent system parts means that we can analyze them individually as pure sequential programs [end of course].
No interaction between concurrent system parts means that we can analyze them individually as pure sequential programs.

Interaction occurs in form of:

- **Contention** (implicit interaction):
  Multiple concurrent execution units compete for one shared resource.

- **Communication** (explicit interaction):
  Explicit passing of information and/or explicit synchronization.
Models and Terminology

The concurrent programming abstraction

Time-line or Sequence?

Consider time (durations) explicitly:

Rightarrow Real-time systemsRightarrow join the appropriate courses

Consider the sequence of interaction points only:

Rightarrow Non-real-time systemsRightarrow stay in your seat
Models and Terminology

The concurrent programming abstraction

Correctness of concurrent non-real-time systems

[logical correctness]:

• does not depend on clock speeds / execution times / delays
• does not depend on actual interleaving of concurrent processes

holds true for all possible sequences of interaction points (interleavings)
Models and Terminology

The concurrent programming abstraction

Correctness vs. testing in concurrent systems:

Slight changes in external triggers may (and usually does) result in completely different schedules (interleaving):

- Concurrent programs which depend in any way on external influences cannot be tested without modelling and embedding those influences into the test process.
- Designs which are provably correct with respect to the specification and are **independent** of the actual timing behavior are essential.

P.S. some timing restrictions for the scheduling still persist in non-real-time systems, e.g. ‘fairness’
Models and Terminology

The concurrent programming abstraction

Atomic operations:

Correctness proofs / designs in concurrent systems rely on the assumptions of

‘Atomic operations’ [detailed discussion later]:

• Complex and powerful atomic operations ease the correctness proofs, but may limit flexibility in the design
• Simple atomic operations are theoretically sufficient, but may lead to complex systems which correctness cannot be proven in practice.
Models and Terminology

The concurrent programming abstraction

Standard concepts of correctness:

• **Partial correctness:**

\[(P(I) \land \text{terminates (Program (I, O))}) \Rightarrow Q(I, O)\]

• **Total correctness:**

\[P(I) \Rightarrow (\text{terminates (Program (I, O))} \land Q(I, O))\]

where \(I, O\) are input and output sets,
\(P\) is a property on the input set,
and \(Q\) is a relation between input and output sets

do these concepts apply to and are sufficient for concurrent systems?
Models and Terminology

The concurrent programming abstraction

Extended concepts of correctness in concurrent systems:
¬ Termination is often not intended or even considered a failure

Safety properties:

\((P(I) \land \text{Processes}(I, S)) \Rightarrow \square Q(I, S)\)

where \(\square Q\) means that \(Q\) does always hold

Liveness properties:

\((P(I) \land \text{Processes}(I, S)) \Rightarrow \diamondsuit Q(I, S)\)

where \(\diamondsuit Q\) means that \(Q\) does eventually hold (and will then stay true) and \(S\) is the current state of the concurrent system
Safety properties:

\[(P(I) \wedge \text{Processes}(I,S)) \Rightarrow \Box Q(I,S)\]

where \(\Box Q\) means that \(Q\) does always hold

Examples:

- Mutual exclusion (no resource collisions)
- Absence of deadlocks
  (and other forms of ‘silent death’ and ‘freeze’ conditions)
- Specified responsiveness or free capabilities
  (typical in real-time / embedded systems or server applications)
Models and Terminology

The concurrent programming abstraction

Liveness properties:

$$(P(I) \land \text{Processes}(I, S)) \Rightarrow \Box Q(I, S)$$

where $\Box Q$ means that $Q$ does eventually hold (and will then stay true)

Examples:

- Requests need to complete eventually
- The state of the system needs to be displayed eventually
- No part of the system is to be delayed forever (fairness)

Interesting liveness properties can be very hard to prove
1 CPU per control-flow

Specific configurations only, e.g.:

- Distributed µcontrollers.
- Physical process control systems:
  1 cpu per task, connected via a bus-system.

Process management (scheduling) not required.

Shared memory access need to be coordinated.
Introduction to processes and threads

1 CPU for all control-flows

- OS: emulate one CPU for every control-flow:
  - Multi-tasking operating system

- Support for memory protection essential.
- Process management (scheduling) required.
- Shared memory access need to be coordinated.
Processes

Process ::= Address space + Control flow(s)

Kernel has full knowledge about all processes as well as their states, requirements and currently held resources.
Threads (individual control-flows) can be handled:

- **Inside** the OS:
  - Kernel scheduling.
  - Thread can easily be connected to external events (I/O).

- **Outside** the OS:
  - User-level scheduling.
  - Threads may need to go through their parent process to access I/O.
Symmetric Multiprocessing (SMP)

All CPUs share the same physical address space (and access to resources).

Any process / thread can be executed on any available CPU.
Introduction to processes and threads

Processes ↔ Threads

Also processes can share memory and the specific definition of threads is different in different operating systems and contexts:

- Threads can be regarded as a group of processes, which share some resources (process-hierarchy).
- Due to the overlap in resources, the attributes attached to threads are less than for ‘first-class-citizen-processes’.
- Thread switching and inter-thread communication can be more efficient than switching on process level.
- Scheduling of threads depends on the actual thread implementations:
  - e.g. user-level control-flows, which the kernel has no knowledge about at all.
  - e.g. kernel-level control-flows, which are handled as processes with some restrictions.
Introduction to processes and threads

Process Control Blocks (PCBs)

- **Process Id**
- **Process state:**
  {created, ready, executing, blocked, suspended, bored …}
- **Scheduling attributes:**
  Priorities, deadlines, consumed CPU-time, …
- **CPU state:** Saved/restored information while context switches (incl. the program counter, stack pointer, …)
- **Memory attributes / privileges:**
  Memory base, limits, shared areas, …
- **Allocated resources / privileges:**
  Open and requested devices and files, …

… PCBs (links thereof) are commonly enqueued at a certain state or condition (awaiting access or change in state)
Process states

- **created**: the task is ready to run, but not yet considered by any dispatcher waiting for admission
- **ready**: ready to run waiting for a free CPU
- **running**: holds a CPU and executes
- **blocked**: not ready to run waiting for a resource
Process states

- **created**: the task is ready to run, but not yet considered by any dispatcher
  - waiting for admission
- **ready**: ready to run
  - waiting for a free CPU
- **running**: holds a CPU and executes
- **blocked**: not ready to run
  - waiting for a resource
- **suspended** states: swapped out of main memory
  - (none time critical processes)
  - waiting for main memory space (and other resources)
Process states

- **created**: the task is ready to run, but not yet considered by any dispatcher waiting for admission
- **ready**: ready to run waiting for a free CPU
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- **suspended** states: swapped out of main memory (none time critical processes) waiting for main memory space (and other resources)

.dispatching and suspending can now be independent modules
Introduction to Concurrency

**Process states**

- **Creation**
- **Admitted**
- **Ready**
- **Executing**
- **Blocked**
- **Suspended**

**Directions**

- CPU
- Pre-emption or cycle done
- Termination
- Dispatch
- Unblock
- Unblock
- Suspend (swap-out)
- Suspend (swap-out)
- Swap-in
- Swap-out
- Block or synchronize

**States**

- Ready, suspended
- Blocked, suspended
- Blocked
- Batch

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UNIX processes

In UNIX systems tasks are created by ‘cloning’

 pid = fork ();

resulting in a duplication of the current process

 ... returning ‘0’ to the newly created process (the ‘child’ process)

 ... returning the process id of the child process to the creating process (the ‘parent’ process)

 ... or returning ‘-1’ as C-style indication of a failure (in void of actual exception handling)

Frequent usage:

 if (fork () == 0) {
   ... the child’s task ...
   ... often implemented as: exec (“absolute path to executable file“, “args“);
   exit (0); /* terminate child process */
 } else {
   ... the parent’s task ...
   pid = wait (); /* wait for the termination of one child process */
 }
UNIX processes

Communication between UNIX tasks (‘pipes’)

```c
int data_pipe [2], c, rc;
if (pipe (data_pipe) == -1) {
  perror ("no pipe"); exit (1);
}
if (fork () == 0) {
  close (data_pipe [1]);
  while ((rc = read (data_pipe [0], &c, 1)) > 0) {
    putchar (c);
  }
  if (rc == -1) {
    perror ("pipe broken");
    close (data_pipe [0]);
    exit (1);
  }
  close (data_pipe [0]); exit (0);
} else {
  close (data_pipe [0]);
  while ((c = getchar ()) > 0) {
    if (write(data_pipe[1], &c, 1) == -1) {
      perror ("pipe broken");
      close (data_pipe [1]);
      exit (1);
    };
  }
  close (data_pipe [1]);
  pid = wait ();
}
```
Concurrent programming languages

Requirement

• Concept of tasks, threads or other potentially concurrent entities

Frequently requested essential elements

• Support for management or concurrent entities (create, terminate, …)
• Support for contention management (mutual exclusion, …)
• Support for synchronization (semaphores, monitors, …)
• Support for communication (message passing, shared memory, rpc …)
• Support for protection (tasks, memory, devices, …)
Concurrent programming languages

Language candidates

Explicit concurrency
- Ada, C++, Rust
- Chill
- Erlang
- Go
- Chapel, X10
- Occam, CSP
- All .net languages
- Java, Scala, Clojure
- Algol 68, Modula-2, Modula-3
- ...

Implicit (potential) concurrency
- Lisp, Haskell, Caml, Miranda, and any other functional language
- Smalltalk, Squeak
- Prolog
- Esterel, Lustre, Signal

Wannabe concurrency
- Ruby, Python [mostly broken due to global interpreter locks]

No support:
- Eiffel, Pascal
- C
- Fortran, Cobol, Basic...

Libraries & interfaces (outside language definitions)
- POSIX
- MPI (Message Passing Interface)
- ...

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Languages with implicit concurrency: e.g. functional programming

Implicit concurrency in some programming schemes

Quicksort in a functional language (here: Haskell):

```haskell
qsort [] = []
qsort (x:xs) = qsort [y | y <- xs, y < x] ++ [x] ++ qsort [y | y <- xs, y >= x]
```

Pure functional programming is side-effect free

Parameters can be evaluated independently could run concurrently

Some functional languages allow for lazy evaluation, i.e. sub-expressions are not necessarily evaluated completely:

```haskell
borderline = (n /= 0) && (g (n) > h (n))
```

If n equals zero then the evaluation of g(n) and h(n) can be stopped (or not even be started).

Concurrent program parts should be interruptible in this case.

Short-circuit evaluations in imperative languages assume explicit sequential execution:

```haskell
if Pointer /= nil and then Pointer.next = nil then ...
```
Summary

Concurrency – The Basic Concepts

• Forms of concurrency

• Models and terminology
  • Abstractions and perspectives: computer science, physics & engineering
  • Observations: non-determinism, atomicity, interaction, interleaving
  • Correctness in concurrent systems

• Processes and threads
  • Basic concepts and notions
  • Process states

• Concurrent programming languages:
  • Explicit concurrency: e.g. Ada, Chapel
  • Implicit concurrency: functional programming – e.g. Haskell, Caml