

COMP 3610/6361

(recording of 02/08/23)

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- Ed Stein
- Quiz (Voluntary)
- Assessment 1
- Course Representative

successor / next step in a computation :

$$\langle E, s \rangle \rightarrow \langle E', s' \rangle$$

- chaining is possible

$$\langle E, s \rangle \rightarrow \langle E', s' \rangle \rightarrow \langle E'', s'' \rangle \dots$$

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Derivation tree of one single computation step

$$\begin{array}{c} \vdots \\ \frac{\vdots}{a' \rightarrow b'} \qquad \frac{\vdots}{c' \rightarrow d'} \\ \hline a \rightarrow b \end{array} \qquad \text{rules}$$

provides a proof  
that the successor  
is reachable.

$\langle 42, \ell, \phi \rangle$  does not have a successor,  
we say  $\langle 42 +! \ell, \phi \rangle$  is stuck.

$$\frac{\begin{array}{c} \langle !\ell, \phi \rangle \times \\ \text{we can't find a} \\ \text{successor} \end{array}}{\begin{array}{c} \langle 42 +! \ell, \phi \rangle \rightarrow \\ \uparrow \\ \text{the empty function} \\ \text{empty storage/store} \\ \text{no location has been assigned a value.} \end{array}} \text{op}^2$$

Slide 43

$$S + \{l_2 \mapsto 0\} = S$$

$$\frac{\overline{< \ell_2 := 0, s >} \rightarrow < \text{skip}, s + \{ \ell_2 \mapsto 0 \} >}{< E, s > \rightarrow < \text{skip; while } !\ell_1 \geq 1 \text{ do...}, s + \{ \ell_2 \mapsto 0 \} >} \quad \begin{matrix} \text{assign 1} \\ \text{seq 2} \end{matrix}$$

$E = l_2 := 0$ ; while  $!l_1 \geq 1$  do ( $l_2 := !l_2 + !l_1$ ;  $l_1 := !l_1 + -1$ )  
 $s = \{l_1 \mapsto 3, l_2 \mapsto 0\}$

$\langle \text{skip}; \text{while } !l_1 \geq 1 \text{ do } (\dots), s \rangle$

$\rightarrow \langle \text{while } !l_1 \geq 1 \text{ do } \dots, s \rangle$

$\rightarrow \langle \text{if } !l_1 \geq 1 \text{ then "body"; while } (!l_1 \geq 1) \dots \text{ else skip}, s \rangle$

$\rightarrow \langle \text{if } 3 \geq 1 \text{ then "body"; while } \dots \text{ else skip}, s \rangle$

$\rightarrow \dots$

$E = l_2 := 0; \text{while } !l_1 \geq 1 \text{ do } (l_2 := !l_2 + !l_1; l_1 := !l_1 + -1)$   
 $s = \{l_1 \mapsto 3, l_2 \mapsto 0\}$

$(\ell := 1; 0) + (\ell := 2; 0), \{\ell \mapsto 0\} \rangle$

$\rightarrow \langle \text{skip}; 0 + (\ell := 2; 0), \{\ell \mapsto 1\} \rangle$

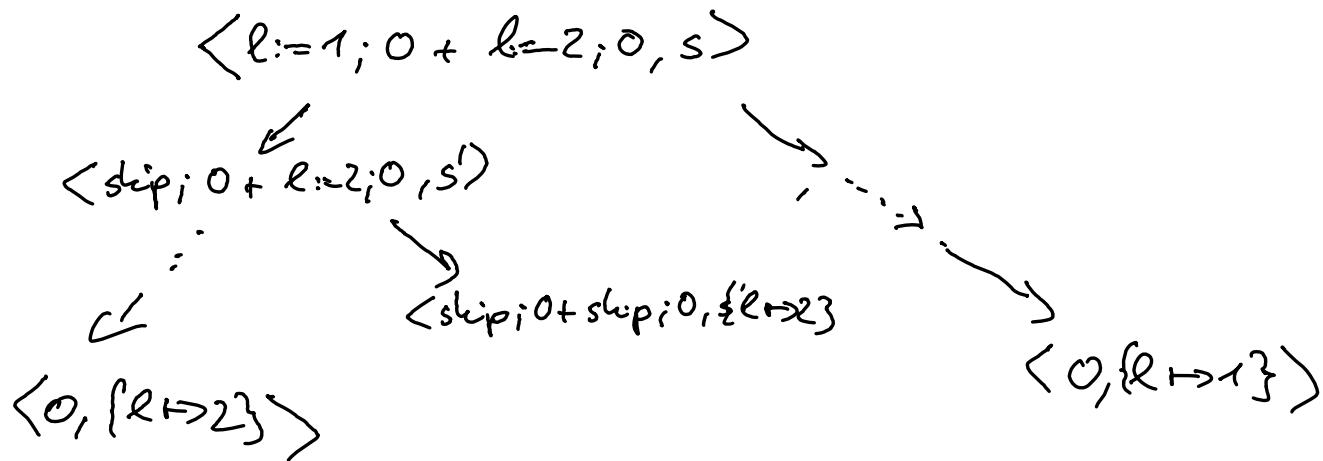
$\rightarrow \langle 0 + (\ell := 2; 0), \{\ell \mapsto 1\} \rangle$

$\rightarrow \langle 0 + (\text{skip}; 0), \{\ell \mapsto 2\} \rangle$

$\rightarrow \langle 0 + 0, \{\ell \mapsto 2\} \rangle$

$\rightarrow \langle 0, \{\ell \mapsto 2\} \rangle$

$| \cap P$  (incl  $OP^1$  and  $OP2$ ) +  $OP1' + OP2'$



Non-deterministic  
language

$\Gamma \vdash E : T$  ← type  
↑  
expression  
type environment / assumptions in ITP

$E$  is of type  $T$  if the assumptions  
in  $\Gamma$  are satisfied.

$$\frac{}{\{3 \vdash \text{false} : \text{bool}\}} \text{bool}$$

$$\frac{\frac{\frac{\text{if}}{\{3 \vdash 2 : \text{int}\}} \text{if}}{\frac{\frac{\text{if}}{\{2 \vdash 3 : \text{int}\}} \text{if}}{\frac{\text{int}}{\{3 \vdash 4 : \text{int}\}}} \text{opt}} \text{if}}{\{3 \vdash 3 + 4 : \text{int}\}} \text{opt}$$

$\{3 \vdash \text{if false then } 2 \text{ else } 3 + 4 : \text{int}\}$  if