anybody who …

… would like to see immediate real-world involvement in his/her work

… would like to learn how to create predictability and fault-tolerant complex systems

… would like to know more about the usage of 95% of all µprocessors
Lectures:
• 2 lectures à 1.5h per week … all the nice stuff and theory
  Monday & Tuesday, 9:00-10:30; in MCC T4, PSYC G8 resp.

Laboratories:
• 2 hours per week … all the rough stuff and practice
  Monday 11:00-13:00 or Tuesday 13:00-15:00 – in CSIT N114, N113 resp.
  laboratory enrolment: https://cs.anu.edu.au/streams/

Resources:
• introduced in the lectures and collected on the course page:
  … as well as schedules, slides, sources, etc. pp. … keep an eye on this page!

Assessment:
• exam at the end of the course (70%) plus laboratories performance (30%)
Introduction & Real-Time Languages

Features of a Real-Time System?

- Fast context switches?
- Small size?
- Quick responds to external interrupts?
- Multitasking?
- ‘low level’ programming interfaces?
- Interprocess communication tools?
- High processor utilization?

References for this chapter

[Berry99] Gérard Berry
The Esterel v5 Language Primer
(Version 5.21 release 2.0)
Technical report: centre de Mathématiques Appliquées, Ecole des Mines and INRIA

[Bollella01] Greg Bollella, Ben Brosgol, Steve Furr, David Hardin, Peter Dibble, James Gosling, Mark Turnbull & Rudy Belliardi
The Real-Time Specification for Java
http://www.rtj.org

[Burns01] Alan Burns and Andy Wellings
Real-Time Systems and Programming Languages
Addison Wesley, third edition, 2001

[Cohen96] Norman H. Cohen
Ada as a second language
McGraw-Hill series in computer science, 2nd edition

[GI98] eds. GI-working group 4.4.2
PEARL 90, Language report, Version 2.2
Technical report GI

all references and some links are available on the course page
Features of a Real-Time System

The correctness of a real-time systems depends on:

1. the logical correctness of the results as well as
2. the time the result was delivered

according to the specification.

☞ All results are to be delivered just-in-time – not too early, not too late.

Timing constraints are specified in many different ways …
… often as a response to ‘external’ events ⇒ reactive systems

Typical Real-Time Systems

• System sizes vary from traffic light controllers, or heating regulators to big aircrafts (Boeing, Airbus), satellites, or high speed trains (TGV)
• High degree of concurrency
• Close connections to real-world entities (sensors, actuators)
• Often a part of a real-world device (embedded systems)
• Failures may often lead to loss of life, or environmental damages.

☞ Predictability is more important than any other criterion

Typical Real-Time Operating Systems

Often implemented as an integrated run-time environment
i.e. there is ‘no operating system’ (⇒ embedded systems)

⇒ RT-OSs have the smallest possible impact on the timing behaviour

Typical Real-Time Systems Components

Discrete

Analogue

Physical

Phenomena

Typ. non real-time system

Typ. real-time system

Typ. embedded system
what is this course about?

Parts of a Real-Time System

1. A set of physical processes
2. **Physical sensors** transforming all values into analogue voltages
3. Physical actuators delivering forces, lights, temperatures, etc. pp.
4. **D/A and A/D convertors** synchronize and discretize all voltages
5. **Interfaces** and buses feeding the measurements into a computer
6. One or many **controllers** (computers)
7. **Algorithms, Languages & Real-Time Software**

Programming styles

- **Imperative** (sequential)
  - Ada, RT-JAVA, Eiffel, ...
- **Functional** (recursive)
  - Lisp, OCaml, ...
- **Declarative** (logic)
  - Prolog, ...
- **Data-flow machines**
  - Lustre, Signal, ...
- **(hierarchical) Finite state machines**
  - synchronous languages: Esterel, syncEiffel, synERJY, ...

What makes a language suitable for real-time systems?

- **Predictability**
  - no operations which will lead to unforeseeable timing behaviours (e.g. garbage collection)
- **Real-time!!**
- **Concurrence**
- **Distribution** support for message passing or rpc
- **Reliability** detect errors at compile-time or in the run-time environment
- **Large systems** scalable, modular, or object-oriented + separate compilation
Programming styles

Languages considered in this course

- Ada95 (used for assignments → introduced first)
- Esterel
- Pearl
- Real-time JAVA
- POSIX
... others in places

Ada95

Ada95 is a standardized (ISO/IEC 8652:1995(E)) ‘general purpose’ language with core language primitives for

- strong typing, separate compilation (specification and implementation), object-orientation,
- concurrency, monitors, rpcs, timeouts, scheduling, priority ceiling locks
- strong run-time environments

... and standardized language-annexes for

- additional real-time features, distributed programming, system-level programming, numeric, informations systems, safety and security issues.

A crash course

... refreshing:

- specification and implementation (body) parts, basic types
- exceptions
- information hiding in specifications (‘private’)
- generic programming
- class-wide programming (‘tagged types’)
- abstract types and dispatching

Basics

... introducing:

- specification and implementation (body) parts
- constants
- some basic types (integer specifics)
- some type attributes
- parameter specification
package Queue_Pack_Simple is
    constant QueueSize := 10;
    type Element is new Positive range 1_000..40_000;
    type Marker is mod QueueSize;
    type List is array (Marker'Range) of Element;
    type Queue_Type is record
        Top, Free : Marker := Marker'First;
        Elements  : List;
    end record;

procedure Enqueue (Item : in Element; Queue : in out Queue_Type);
procedure Dequeue (Item : out Element; Queue : in out Queue_Type);
end Queue_Pack_Simple;

package body Queue_Pack_Simple is

procedure Enqueue (Item : in Element; Queue : in out Queue_Type) is
end Enqueue;

procedure Dequeue (Item : out Element; Queue : in out Queue_Type) is
    begin      Item := Queue.Elements (Queue.Top);      Queue.Top := Queue.Top - 1;
end Dequeue;

end Queue_Pack_Simple;

procedure Queue_Test_Simple is
    Queue : Queue_Type;
    Item  : Element;
begin
    Enqueue (2000, Queue);
    Dequeue (Item, Queue);
    Dequeue (Item, Queue); -- will produce an unpredictable result!
end Queue_Test_Simple;

… introducing:
• exception handling
• enumeration types
• functional type attributes
package Queue_Pack_Exceptions is

QueueSize : constant Integer := 10;
type Element is (Up, Down, Spin, Turn);
type Marker is mod QueueSize;
type List is array (Marker’Range) of Element;
type Queue_State is (Empty, Filled);
type Queue_Type is record
    Top, Free : Marker := Marker’First;
    State : Queue_State := Empty;
    Elements : List;
end record;

procedure Enqueue (Item: in Element; Queue: in out Queue_Type);
procedure Dequeue (Item: out Element; Queue: in out Queue_Type);

Queueoverflow, Queueunderflow : exception;
end Queue_Pack_Exceptions;

package body Queue_Pack_Exceptions is

procedure Enqueue (Item: in Element; Queue: in out Queue_Type) is
begin
    if Queue.State = Filled and Queue.Top = Queue.Free then
        raise Queueoverflow;
    end if;
    Queue.Elements (Queue.Free) := Item;
    Queue.Free := Marker’Pred (Queue.Free);
    Queue.State := Filled;
end Enqueue;

procedure Dequeue (Item: out Element; Queue: in out Queue_Type) is
begin
    if Queue.State = Empty then
        raise Queueunderflow;
    end if;
    Item := Queue.Elements (Queue.Top);
    Queue.Top := Marker’Pred (Queue.Top);
    if Queue.Top = Queue.Free then
        Queue.State := Empty;
    end if;
end Dequeue;
end Queue_Pack_Exceptions;

with Queue_Pack_Exceptions; use Queue_Pack_Exceptions;
with Ada.Text_IO; use Ada.Text_IO;

procedure Queue_Test_Exceptions is
    Queue : Queue_Type;
    Item : Element;
begin
    Enqueue (Turn, Queue);
    Dequeue (Item, Queue);
    Dequeue (Item, Queue); -- will produce a 'Queue underflow'

exception
    when Queueunderflow => Put ('Queue underflow');
    when Queueoverflow => Put ('Queue overflow');
end Queue_Test_Exceptions;

... introducing:
• private ⇒ assignments and comparisons are allowed
• limited private ⇒ entity cannot be assigned or compared
package Queue_Pack_Private is
   QueueSize : constant Integer := 10;
   type Element is new Positive range 1..1000;
   type Queue_Type is limited private;
   procedure Enqueue (Item: in Element; Queue: in out Queue_Type);
   procedure Dequeue (Item: out Element; Queue: in out Queue_Type);
   Queueoverflow, Queueunderflow : exception;
private
   type Marker is mod QueueSize;
   type List is array (Marker'Range) of Element;
   type Queue_State is (Empty, Filled);
   type Queue_Type is record
      ...
      Top, Free : Marker := Marker'First;
      State     : Queue_State := Empty;
      Elements  : List;
   end record;
end Queue_Pack_Private;

package body Queue_Pack_Private is
   procedure Enqueue (Item: in Element; Queue: in out Queue_Type) is
      begin
         if Queue.State = Filled and Queue.Top = Queue.Free then raise Queueoverflow; end if;
         Queue.Elements (Queue.Free) := Item;
         Queue.Free := Queue.Free - 1;
         Queue.State := Filled;
      end Enqueue;
   procedure Dequeue (Item: out Element; Queue: in out Queue_Type) is
      begin
         if Queue.State = Empty then raise Queueunderflow; end if;
         Item := Queue.Elements (Queue.Top);
         Queue.Top := Queue.Top - 1;
         if Queue.Top = Queue.Free then Queue.State := Empty; end if;
         end Dequeue;
      end Dequeue;
end Queue_Pack_Private;

with Queue_Pack_Private; use Queue_Pack_Private;
with Ada.Text_IO; use Ada.Text_IO;
procedure Queue_Test_Private is
   Queue, Queue_Copy : Queue_Type;
   Item              : Element;
begin
   Queue_Copy := Queue;
   -- compiler-error: left hand of assignment must not be limited type
   Enqueue (Item => 1, Queue => Queue);
   Dequeue (Item, Queue);
   Dequeue (Item, Queue); -- will produce a 'Queue underflow'
   exception
      when Queueunderflow => Put ("Queue underflow");
      when Queueoverflow => Put ("Queue overflow");
   end Queue_Test_Private;

... introducing:
• specification of generic packages
• instantiation of generic packages
A generic queue **specification**

```ada
package Queue_Pack_Generic is
  QueueSize: constant Integer := 10;
  type Queue_Type is limited private;
  procedure Enqueue (Item: in  Element; Queue: in out Queue_Type);
  procedure Dequeue (Item: out Element; Queue: in out Queue_Type);
  Queueoverflow, Queueunderflow : exception;
private
  type Marker is mod QueueSize;
  type List is array (Marker'Range) of Element;
  type Queue_State is (Empty, Filled);
  type Queue_Type is record
    Top, Free : Marker := Marker'First;
    State : Queue_State := Empty;
    Elements : List;
  end record;
end Queue_Pack_Generic;
```

A generic queue **implementation**

```ada
package body Queue_Pack_Generic is
  procedure Enqueue (Item: in Element; Queue: in out Queue_Type) is
    begin
      if Queue.State = Filled and Queue.Top = Queue.Free then
        raise Queueoverflow;
      end if;
      Queue.Elements (Queue.Free) := Item;
      Queue.Free  := Queue.Free - 1;
      Queue.State := Filled;
    end Enqueue;

  procedure Dequeue (Item: out Element; Queue: in out Queue_Type) is
    begin
      if Queue.State = Empty then
        raise Queueunderflow;
      end if;
      Item      := Queue.Elements (Queue.Top);
      Queue.Top := Queue.Top - 1;
      if Queue.Top = Queue.Free then Queue.State := Empty; end if;
      end Dequeue;
    end Dequeue;
end Queue_Pack_Generic;
```

A generic queue **test program**

```ada
with Queue_Pack_Generic;
with Ada.Text_IO;   use Ada.Text_IO;

procedure Queue_Test_Generic is
  use Queue_Pack_Positive;
  Queue : Queue_Type;
  Item  : Positive;
begin
  Enqueue (Item => 1, Queue => Queue);   Dequeue (Item, Queue);   Dequeue (Item, Queue); -- will produce a 'Queue underflow'
  exception
    when Queueunderflow => Put ("Queue underflow");
    when Queueoverflow  => Put ("Queue overflow");
end Queue_Test_Generic;
```

**Ada95**

Object oriented programming 1

... introducing:

- tagged types = the Ada-way to say that this type can be extended
- derivation of tagged types
- method overwriting
- usage of parent entities
An open queue base class specification

package Queue_Pack_Object_Base is
  QueueSize : constant Integer := 10;
  type Element is new Positive range 1..1000;
  type Marker is mod QueueSize;
  type List is array (Marker'Range) of Element;
  type Queue_State is (Empty, Filled);
  type Queue_Type is
      tagged record
        Top, Free : Marker := Marker'First;
        State     : Queue_State := Empty;
        Elements  : List;
      end record;

  procedure Enqueue (Item: in Element; Queue: in out Queue_Type);
  procedure Dequeue (Item: out Element; Queue: in out Queue_Type);

  Queueoverflow, Queueunderflow : exception;
end Queue_Pack_Object_Base;

A derived open queue class specification

with Queue_Pack_Object_Base; use Queue_Pack_Object_Base;

package Queue_Pack_Object is
  type Ext_Queue_Type is new Queue_Type with record
    Reader       : Marker := Marker'First;
    Reader_State : Queue_State := Empty;
  end record;
  procedure Enqueue (Item: in Element; Queue: in out Ext_Queue_Type);
  procedure Read_Queue (Item: out Element; Queue: in out Ext_Queue_Type);
end Queue_Pack_Object;

An open queue base class implementation

package body Queue_Pack_Object_Base is
  procedure Enqueue (Item: in Element; Queue: in out Queue_Type) is
    begin
      if Queue.State = Filled and Queue.Top = Queue.Free then
        raise Queueoverflow;
      end if;
      Queue.Elements (Queue.Free) := Item;
      Queue.Free := Queue.Free - 1;
      Queue.State := Filled;
      Enqueue;
    end Enqueue;

  procedure Dequeue (Item: out Element; Queue: in out Queue_Type) is
    begin
      if Queue.State = Empty then
        raise Queueunderflow;
      end if;
      Item := Queue.Elements (Queue.Top);
      Queue.To := Queue.Top - 1;
      if Queue.To = Queue.Free then Queue.State := Empty; end if;
      Dequeue;
    end Dequeue;
end Queue_Pack_Object_Base;

A derived open queue class implementation

package body Queue_Pack_Object is
  procedure Enqueue (Item: in Element; Queue: in out Ext_Queue_Type) is
    begin
      Enqueue (Item, Queue_Type (Queue));
      Queue.Reader_State := Filled;
      Enqueue;
    end Enqueue;

  procedure Read_Queue (Item: out Element; Queue: in out Ext_Queue_Type) is
    begin
      if Queue.Reader_State = Empty then
        raise Queueunderflow;
      end if;
      Item := Queue.Elements (Queue.Reader);
      Queue.Reader := Queue.Reader - 1;
      if Queue.Reader = Queue.Free then Queue.Reader_State := Empty; end if;
      Read_Queue;
    end Read_Queue;
end Queue_Pack_Object;
An open class test program

```ada
with Queue_Pack_Object_Base; use Queue_Pack_Object_Base;
with Queue_Pack_Object; use Queue_Pack_Object;
with Ada.Text_IO; use Ada.Text_IO;

procedure Queue_Test_Object is
  Queue : Ext.Queue_Type;
  Item  : Element;
begin
  Enqueue (Item => 1, Queue => Queue);
  Read_Queue (Item, Queue);
  Enqueue (Item => 5, Queue => Queue);
  Dequeue (Item, Queue);
  Dequeue (Item, Queue);
  Dequeue (Item, Queue); -- will produce a 'Queue underflow'
  exception
  when Queueunderflow => Put ("Queue underflow");
  when Queueoverflow    => Put ("Queue overflow");
end Queue_Test_Object;
```

Object oriented programming II

- private tagged types
- objects which are protected against their children also

An encapsulated queue base class specification

```ada
package Queue_Pack_Object_Base_Private is
  QueueSize : constant Integer := 10;
  type Element is new Positive range 1..1000;
  type Queue_Type is
tagged limited private
    record
      Top, Free : Marker      := Marker'First;
      State     : Queue_State := Empty;
      Elements  : List;    end record;
  end Queue_Type;
  procedure Enqueue (Item: in  Element; Queue: in out Queue_Type);
  procedure Dequeue (Item: out Element; Queue: in out Queue_Type);
  Queueoverflow, Queueunderflow : exception;
private
  type Marker is mod QueueSize;
  type List is array (Marker'Range) of Element;
  type Queue_State is (Empty, Filled);
  procedure Enqueue (Item: in Element; Queue: in out Queue_Type);
  procedure Dequeue (Item: out Element; Queue: in out Queue_Type);
end Queue_Pack_Object_Base_Private;
```

An encapsulated queue base class implementation

```ada
package body Queue_Pack_Object_Base_Private is
  procedure Enqueue (Item: in  Element; Queue: in out Queue_Type) is
  begin
    if Queue.State = Filled and Queue.Top = Queue.Free then
      raise Queueoverflow;
    end if;
    Queue.Elements (Queue.Free) := Item;
    Queue.Free  := Queue.Free - 1;
    if Queue.Top = Queue.Free then Queue.State := Empty; end if;
  end Enqueue;
  procedure Dequeue (Item: out Element; Queue: in out Queue_Type) is
  begin
    if Queue.State = Empty then
      raise Queueunderflow;
    end if;
    Item   := Queue.Elements (Queue.Top);
    Queue.To := Queue.Top - 1;
    if Queue.To > Queue.Free then Queue.State := Empty; end if;
    end Dequeue;
end Queue_Pack_Object_Base_Private;
```
Real-Time & Embedded Systems

A derived encapsulated queue class specification

with Queue_Pack_Object_Base_Private; use Queue_Pack_Object_Base_Private;
package Queue_Pack_Object_Private is
  type Ext_Queue_Type is new Queue_Type with private;
  subtype Depth_Type is Positive range 1..QueueSize;
  procedure Look_Ahead (Item: out Element;
                        Depth: in Depth_Type; Queue: in out Ext_Queue_Type);
private
  type Ext_Queue_Type is new Queue_Type with null record;
end Queue_Pack_Object_Private;

package body Queue_Pack_Object_Private is
  procedure Look_Ahead (Item: out Element;
                        Depth: in Depth_Type; Queue: in out Ext_Queue_Type) is
    Storage     : Queue_Type;
    ShuffleItem : Element;
  begin
    for I in 1..Depth - 1 loop
      Dequeue (ShuffleItem, Queue);
      Enqueue (ShuffleItem, Storage);
    end loop;
    Dequeue (Item, Queue);
    Enqueue (Item, Storage);
  (...) Read_The_Rest:
    begin
      for I in 1..QueueSize - Depth loop
        Dequeue (ShuffleItem, Queue);
        Enqueue (ShuffleItem, Storage);
      end loop;
      exception
        when Queueunderflow => null; -- restore is done
    end Restore_The_Queue;
end Look_Ahead;
end Queue_Pack_Object_Private;

An encapsulated class test program

with Queue_Pack_Object_Base_Private; use Queue_Pack_Object_Base_Private;
with Queue_Pack_Object_Private; use Queue_Pack_Object_Private;
with Ada.Text_IO; use Ada.Text_IO;
procedure Queue_Test_Object_Private is
  Queue : Ext_Queue_Type;
  Item  : Element;
begin
  Enqueue (Item => 1, Queue => Queue);
  Enqueue (Item => 1, Queue => Queue);
  Look_Ahead (Item => Item, Depth => 2, Queue => Queue);
  Enqueue (Item => 5, Queue => Queue);
  Dequeue (Item, Queue);
  Dequeue (Item, Queue);
  Dequeue (Item, Queue);
  Dequeue (Item, Queue); -- will produce a 'Queue underflow'
  exception
    when Queueunderflow => Put ("Queue underflow");
    when Queueoverflow  => Put ("Queue overflow");
end Queue_Test_Object_Private;
... introducing:

- protected types
- tasks (definition, instantiation and termination)
- task synchronisation
- entry guards
- entry calls
- accept and selected accept statements

```
package Queue_Pack_Protected is
  QueueSize : constant Integer := 10;
  subtype Element is Character;
  type Queue_Type is limited private;
  Protected type Protected_Queue is
    entry Enqueue (Item: in Element);
    entry Dequeue (Item: out Element);
  private
    type Marker is mod QueueSize;
    type List is array (Marker'Range) of Element;
    type Queue_State is (Empty, Filled);
  protected body Protected_Queue is
    entry Enqueue (Item: in Element) when Queue.State = Empty or Queue.Top /= Queue.Free is
      begin
        Queue.Elements (Queue.Free) := Item;
        Queue.Free := Queue.Free - 1;
        Queue.State := Filled;
      end Enqueue;
    entry Dequeue (Item: out Element) when Queue.State = Filled is
      begin
        Item := Queue.Elements (Queue.Top);
        Queue.Top := Queue.Top - 1;
        if Queue.Top = Queue.Free then Queue.State := Empty; end if;
      end Dequeue;
    end Protected_Queue;
end Queue_Pack_Protected;
```

```
with Queue_Pack_Protected; use Queue_Pack_Protected;
with Ada.Text_IO; use Ada.Text_IO;
procedure Queue_Test_Protected is
  Queue : Protected_Queue;
  task Producer is
    entry shutdown;
  task Consumer is
  task body Producer is
    Item : Element;
    Got_It : Boolean;
    begin
      loop
        select
          accept shutdown; exit; -- main task loop
        else
          Get_Immediate (Item, Got_It);
          if Got_It then
            Queue.Enqueue (Item); -- task might be blocked here!
          else
            delay 0.1; -- sec.
          end if;
        end select;
      end loop;
    end Producer;
```
A multitasking protected queue test program (cont.)

```ada
(task body Consumer is
  Item : Element;
  begin
    loop
      Queue.Dequeue (Item); -- task might be blocked here!
      Put ("Received: "); Put (Item); Put_Line ("!");
      if Item = 'q' then
        Put_Line ("Shutting down producer"); Producer.Shutdown;
        Put_Line ("Shutting down consumer"); exit; -- main task loop
      end if;
    end loop;
  end Consumer;

begin
  null;
end Queue_Test_Protected;
```

... introducing:
- abstract tagged types
- abstract subroutines
- concrete implementation of abstract types
- dispatching to different packages, tasks, and partitions according to concrete types

### Abstract types & dispatching

... introducing:

- abstract tagged types
- abstract subroutines
- concrete implementation of abstract types
- dispatching to different packages, tasks, and partitions according to concrete types

#### Ada95

... introducing:

- abstract tagged types
- abstract subroutines
- concrete implementation of abstract types
- dispatching to different packages, tasks, and partitions according to concrete types

### Abstract types & dispatching

#### An abstract queue specification

```ada
package Queue_Pack_Abstract is
  subtype Element is Character;
  type Queue_Type is abstract tagged limited private;
  procedure Enqueue (Item: in  Element; Queue: in out Queue_Type) is
    abstract;
  procedure Dequeue (Item: out Element; Queue: in out Queue_Type) is
    abstract;

private
  type Queue_Type is abstract tagged limited null record;
end Queue_Pack_Abstract;
```

#### A concrete queue specification

```ada
package Queue_Pack_Concrete is
  QueueSize : constant Integer := 10;
  type Real_Queue is new Queue_Type with
    private
      type Marker is mod QueueSize;
      type List is array (Marker'Range) of Element;
      type Queue_State is (Empty, Filled);
      type Real_Queue is new Queue_Type with record
        Top, Free : Marker := Marker'First;
        State     : Queue_State := Empty;
        Elements  : List;
        end record;
end Queue_Pack_Concrete;
```
A concrete queue implementation

package body Queue_Pack_Concrete is
  procedure Enqueue (Item: in Element; Queue: in out Real.Queue) is
    begin      if Queue.State = Filled and Queue.Top = Queue.Free then         raise Queueoverflow;      end if;
      Queue.Elements (Queue.Free) := Item;
      Queue.Free  := Queue.Free - 1;
      Queue.State := Filled;   end Enqueue;

  procedure Dequeue (Item: out Element; Queue: in out Real.Queue) is
    begin      if Queue.State = Empty then         raise Queueunderflow;      end if;
      Item      := Queue.Elements (Queue.Top);
      Queue.Top := Queue.Top - 1;
      if Queue.Top = Queue.Free then Queue.State := Empty; end if;
      end Dequeue;
end Queue_Pack_Concrete;

A multitasking dispatching test program

with Queue_Pack_Abstract; use Queue_Pack_Abstract;
with Queue_Pack_Concrete; use Queue_Pack_Concrete;
procedure Queue_Test_Dispatching is
  type Queue_Class is access all Queue_Type'class;

  task Queue_Holder is
    -- could be on an individual partition
    entry Queue_Filled;
  end Queue_Holder;

  task Queue_User is
    -- could be on an individual partition
    entry Send_Queue (Remote_Queue: in Queue_Class);
  end Queue_User;

  task body Queue_Holder is
    Local_Queue : Queue_Class;
    Item        : Element;
    begin      Local_Queue := new Real.Queue; -- could be a different implementation!
                  Queue_User.Send_Queue (Local_Queue);
                  accept Queue_Filled do
                    Dequeue (Item, Local_Queue.all); -- Item will be 'r'
                  end Queue_Filled;
    end Queue_Holder;

  task body Queue_User is
    Local_Queue : Queue_Class;
    Item        : Element;
    begin      Local_Queue := new Real.Queue; -- could be a different implementation!
                accept Send_Queue (Remote_Queue: in Queue_Class) do
                    Enqueue ('r', Remote_Queue.all); -- potentially a rpc!
                    Enqueue ('l', Local_Queue.all);
                end Send_Queue;
                Queue_Holder.Queue_Filled;
                Dequeue (Item, Local_Queue.all); -- Item will be 'l'
                end Queue_User;
    begin null; end Queue_Test_Dispatching;

Real-time & Embedded Systems

Ada95

Ada95 language status

• Established language standard with free and commercial compilers available for all major OSs.
• Stand-alone runtime environments for embedded systems (some are only available commercially).
• Special (yet non-standard) extensions (i.e. language reductions and proof systems) for extreme small footprint embedded systems or high integrity real-time environments available ≫ Ravenscar profile systems.
• has been used and is in use in numberless large scale projects (e.g. in the international space station, and in some spectacular crashes: e.g. Ariane 5)
Transformational ↔ Interactive ↔ Reactive

- Transformational (functional) systems:
  ... generating outputs based on input and stop, utilizing no or a small number of internal states.

- Interactive systems:
  ... servers and other systems in longer long-term operation, requesting occasional inputs, and accepting service-calls, when there are resources to do so.

- Reactive (reflex) systems:
  ... systems, which are reacting to external stimuli only (by generating other stimuli). Can be viewed as a predictable, functional system, which is listening to inputs continuously, while holding enough resources to ensure reasonable reaction times.

Strong synchrony or ‘zero delay’ assumption:

☞ theoretical perspective: all operations are ‘instantaneous’

... or a little bit more realistic:

☞ there is no observable delay, i.e. all operations are finished before the next signal occurs.

Control- ↔ Data-handling

- Data-handling:
  ... continuous data-streams, functional processing (DSPs, ‘number crunching’), ⇒ high bandwidth

- Control-handling:
  ... discrete signals, controlling data-streams and processes, ⇒ low bandwidth

Control-dominated reactive systems

- Real-Time process control:
  ... reaction to (sparse) stimuli in predefined time-spans

- Embedded systems / device control:
  ... local, discrete control

- Complex systems control:
  ... supervision of complex data-streams

- Communication protocols:
  ... control part of communication systems

- Human-machine interface:
  ... switching modes, event handling

- Control logic (hardware):
  ... glue logic, interfaces, pipe control
Determinism

- Control-dominated reactive systems are mostly deterministic
  ... e.g. real-time & embedded systems
- Interactive systems are usually non-deterministic
  ... e.g. operating systems, IP-servers

Is determinism lost automatically in concurrent systems?

---

Module in Esterel:

```esterel
module A_and_B_gives_O;
  input A, B, R;
  output O;
  loop
    [await A || await B];
    emit O;
  end loop;
end module;
```

---

**a simple reactive integrator example**

**Specification:** A module should count the number of metres per second and emit this number as ‘speed’ once per second.

```esterel
Module Speed;
  input  Metre, Second;  relation Metre # Second;
  output Speed: integer;
  loop
    var Distance := 0: integer in
      abort
      every Metre do
        Distance := Distance + 1;
      end every;
    when Second do
      emit Speed (Distance);
    end abort;
  end loop;
end module;
```

---

**Immediate reactions:**

- by default all synchronization points:
  
  ```esterel
  await <signal>;
  abort ... when <signal>;
  every <signal> do ... end every;
  loop ... each <signal>;
  ```

  ... wait for the next signal occurrence (‘rising edge trigger’).

- but with an additional ‘immediate’:

  ```esterel
  await immediate <signal>;
  abort ... when immediate <signal>;
  every immediate <signal> do ... end every;
  loop ... each immediate <signal>;
  ```

  ... a currently active signal will trigger these statements (‘level trigger’).
by default a code block is aborted immediately, when a signal occurs:

\[
\text{abort} \left[ \langle \text{statement} \rangle ; \right] + \\
\text{when } \langle \text{signal} \rangle;
\]

Sometimes the semantic like

‘activate the code block for one last time, when \( \langle \text{signal} \rangle \) occurs’

is more useful and expressed in Esterel as:

\[
\text{weak abort} \left[ \langle \text{statement} \rangle ; \right] + \\
\text{when } \langle \text{signal} \rangle;
\]

the code block is now activated for a ‘final wish’, when \( \langle \text{signal} \rangle \) occurs.

assuming ‘Metre’ and ‘Second’ occur simultaneously \( \Rightarrow \) this code is wrong!

\[
\text{Module Speed;}
\text{input Metre, Second;}
\text{output Speed: integer;}
\text{loop}
\text{var Distance := 0 : integer in}
\text{weak abort}
\text{every Metre do}
\text{Distance := Distance + 1;}
\text{end every;}
\text{when Second do}
\text{emit Speed (Distance);}
\text{end abort;}
\text{end var;}
\text{end loop;}
\text{end module;}
\]

these are no longer exclusive

\( \text{1. above block is immediately aborted} \)
\( \text{2. ‘every Metre’ waits for the next ‘Metre’ ⇒ 1 metre is lost} \)

Parallel signals

‘Metre’ and ‘Second’ occur simultaneously – using ‘weak abort’:

\[
\text{Module Speed;}
\text{input Metre, Second;}
\text{output Speed: integer;}
\text{loop}
\text{var Distance := 0 : integer in}
\text{weak abort}
\text{every Metre do}
\text{Distance := Distance + 1;}
\text{end every;}
\text{when Second do}
\text{emit Speed (Distance);}
\text{end abort;}
\text{end var;}
\text{end loop;}
\text{end module;}
\]

these are no longer exclusive

\( \text{1. above block is immediately aborted} \)
\( \text{2. ‘every Metre’ takes the current ‘Metre’ into account} \)

Parallel signals

‘Metre’ and ‘Second’ occur simultaneously – using ‘every immediate’:

\[
\text{Module Speed;}
\text{input Metre, Second;}
\text{output Speed: integer;}
\text{loop}
\text{var Distance := 0 : integer in}
\text{every immediate Metre do}
\text{Distance := Distance + 1;}
\text{end every;}
\text{when Second do}
\text{emit Speed (Distance);}
\text{end abort;}
\text{end var;}
\text{end loop;}
\text{end module;}
\]

these are no longer exclusive

\( \text{1. above block is immediately aborted} \)
\( \text{2. ‘every immediate’ takes the current ‘Metre’ into account} \)
Causality and synchronous languages

Causality in general terms: ‘the future should not influence the past’

Technically:

**Causal synchronous programs need to be**

☞ **Reactive**
  provides a well-defined output for each signal sequence

☞ **Deterministic**
  provides exactly one output for each signal sequence

---

Causality: counter examples

☞ **non-reactive programming:**

```plaintext
module non-reactive;
  output O;
  present O else emit O end;
end module;
```

☞ **non-deterministic programming:**

```plaintext
module non-deterministic;
  output O;
  present O then emit O end;
end module;
```

☞ **cyclic dependencies with multiple signals:**

```plaintext
module cyclic_dependency;
  output A, B;
  [ present A then emit B end || present B else emit A end ]
end module;
```

---

More on ‘strong synchrony’ and ‘zero delay’ assumptions:

Assuming a system operates in three phases:

1. collect current input signals
2. calculate responses
3. emit new output signals and goto step 1.

☞ The system is assumed to be ‘synchronous’ or ‘instantaneous’, iff the total worst case computation time is smaller than the minimal time between two observable changes in the environment.

☞ Synchronous systems assume a logical rather than a continuous time.
More on ‘strong synchrony’ and ‘zero delay’ assumptions:

In many applications these assumptions are justified and enable:

- a strong analysis and simplification theory (Boolean calculus and automata theory)
- a significantly easier program verification
- an easier hardware implementation

Esterel language status

- Created by the Esterel group at INRIA, France.
- Available freely for Linux, Solaris, and Windows NT.
- Produces C/C++ code, which need to be cross-compiled for an actual target system.
- Currently maintained by Esterel Technologies (a spin-off company).
- No standards.
- Employed in telecommunication, automotive, energy, aerospace, and defence projects by some major companies.

PEARL – Example

```plaintext
MODULE;
SYSTEM;
Alert: Hard_Int(7);
PROBLEM;
SPECIFY Alert INTERRUPT;
   SPECIFY Help TASK GLOBAL;
   SPECIFY Pushed BIT GLOBAL;
   DECLARE Switch BIT INITIAL 0;
Init : TASK MAIN;
   WHEN Alert ACTIVATE Recovery;
   ENABLE Alert;
   Switch := Pushed;
END;
Recovery: TASK PRIORITY 9;
   DISABLE Alert;
   IF Switch = 1 THEN ACTIVATE Help; FIN;
   AFTER 30 MIN ALL 5 MIN DURING 1 HRS ACTIVATE Help;
END;
MODEND;
```

Is a settled standard:

- DIN 66253-1: Basic PEARL 1981
- DIN 66253-2: Full PEARL 1982
- DIN 66253-3: PEARL for distributed systems 1989
**PEARL**

**PEARL – language status**

- Established standard.
- Compilers available for all major OSs (and some RT-OSs) as well as for a number of single-board systems (one free compiler for academic users).
- Used for educational purposes mainly.
- Currently maintained by a German special-interest community and a small company (IEP).

**Real-Time Java**

**Specific Java engines and classes enhance:**

- **Threads**: Priorities, scheduling, and dispatching
- **Memory**: Controlled garbage collection and physical memory access
- **Synchronization**: Ordered queues, and priority ceiling protocols
- **Asynchronism**: Generalized asynchronous event handling, asynchronous transfer of control, timers, and an operational implementation of thread termination

☞ All current real-time Java extensions keep the underlying, consequent object orientation.

☞ Some restrict the language standard, some extend it.

**Real-Time Specification for Java 1.0**

*(final 11/01 – currently no release date for 1.0.1)*

- Enhanced thread model (memory attributes, more precise specs)
- Enabling powerful and highly adaptive scheduling policies
- Introducing scoped, immortal, and physical memory to Java
- Introducing timers, interrupts, and more exceptions
- Higher resolution time model
- Optional support for POSIX signals

☞ **Backward compatible**: offers the full standard Java specification, no syntactical extensions

☞ **Allows for different** Java-engine implementations:
  - In terms of completeness: e.g. scheduling is not mandatory
  - In terms of interpretations: e.g. ‘instantiations per time-span’ (RationalTime) strongly suggests but does not enforce equal distance internals

Currently one reference implementation available (for TimeSys Linux)
Real-Time Java

RT-Java – Language status

- RT-Java is still a consequently object-oriented language.
- Garbage collection can be restricted or even fully suppressed (mandatory requirement for predictable systems).
- How do you program in a clean object oriented manner without garbage collection?
- Using it in hard-realtime environments implies 'program badly' (in terms of strong OOP).

Many potential applications in soft- or mixed-realtime environments.

POSIX

Portable Operating System Interface for Computing Environments

- IEEE/ANSI Std 1003.1 and following
- Program Interface (API) [C Language]
- more than 30 different POSIX standards (a system is 'POSIX compliant', if it implements parts of just one of them!)

POSIX – some of the real-time relevant standards

- IEEE Std 1003.1
  - OS Definition
    - single process, multi process, job control, signals, user groups, file system, file attributes, file device management, file locking, device I/O, device-specific control, system database, pipes, FIFO, ...

- IEEE Std 1003.1b
  - Real-time Extensions
    - real-time signals, priority scheduling, timers, asynchronous I/O, prioritized I/O, synchronized I/O, file sync, mapped files, memory locking, memory protection, message passing, semaphore, ...

- IEEE Std 1003.1c
  - Threads
    - multiple threads within a process; includes support for: thread control, thread attributes, priority scheduling, mutexes, mutex priority inheritance, mutex priority ceiling, and condition variables

- IEEE Std 1003.1d
  - Additional Real-time Extensions
    - new process create semantics (spawn), sporadic server scheduling, execution time monitoring of processes and threads, I/O advisory information, timeouts on blocking functions, device control, and interrupt control

- IEEE Std 1003.1j
  - Advanced Real-time Extensions
    - typed memory, nanosleep improvements, barrier synchronization, reader/writer locks, spin locks, and persistent notification for message queues

- IEEE Std 1003.21
  - Distributed Real-time
    - buffer management, send control blocks, asynchronous and synchronous operations, bounded blocking, message priorities, message labels, and implementation protocols

Frequently found POSIX RT-features include:

- **Timers**: delivery is accomplished using POSIX signals
- **Priority scheduling**: fixed priority, 32 priority levels
- **Real-time signals**: signals with multiple levels of priority
- **Semaphore**: named semaphore
- **Memory queues**: message passing using named queues
- **Shared memory**: memory regions shared between multiple processes
- **Memory locking**: no virtual memory swapping of physical memory pages
**Real-Time & Embedded Systems**

### POSIX – support in some OSs

<table>
<thead>
<tr>
<th>POSIX 1003.1 (Base POSIX)</th>
<th>POSIX 1003.1b (Real-time extensions)</th>
<th>POSIX 1003.1c (Threads)</th>
</tr>
</thead>
<tbody>
<tr>
<td>Solaris</td>
<td>Full support</td>
<td>Full support</td>
</tr>
<tr>
<td>IRIX</td>
<td>Conformant</td>
<td>Full support</td>
</tr>
<tr>
<td>LynxOS</td>
<td>Conformant</td>
<td>Conformant (Version 3.1)</td>
</tr>
<tr>
<td>QNX Neutrino</td>
<td>Full support</td>
<td>Partial support (no memory locking)</td>
</tr>
<tr>
<td>Linux</td>
<td>Full support</td>
<td>Partial support (no timers, no message queues)</td>
</tr>
<tr>
<td>VxWorks</td>
<td>Partial support (different process model)</td>
<td>Supported through third party product</td>
</tr>
</tbody>
</table>

**POSIX – other languages**

POSIX is a ‘C’ standard …

… but bindings to other languages are also (suggested) POSIX standards:

- Ada: 1003.5*, 1003.24 (some PAR approved only, some withdrawn)
- Fortran: 1003.9 (6/92)
- Fortran90: 1003.19 (withdrawn)

… and there are POSIX standards for task-specific POSIX profiles, e.g.:

- Super computing: 1003.10 (6/95)
- Realtime: 1003.13, 1003.13b (3/98)
  - profiles 51-54: combinations of the above RT-relevant POSIX standards ↔ RT-Linux
- Embedded Systems: 1003.13a (PAR approved only)

**Real-Time & Embedded Systems**

### POSIX – example: setting a timer

```c
void timer_create(int num_secs, int num_nsecs)
{
    struct sigaction sa;
    struct sigevent sig_spec;
    sigset_t allsigs;
    struct itimerspec tmr_setting;
    timer_t timer_h;

    /* setup signal to respond to timer */
    sigemptyset(&sa.sa_mask);
    sa.sa_flags = SA_SIGINFO;
    sa.sa_sigaction = timer_intr;
    if (sigaction(SIGRTMIN, &sa, NULL) < 0)
        perror('sigaction');

    sig_spec.sigev_notify = SIGEV_SIGNAL;
    sig_spec.sigev_signo = SIGRTMIN;

    /* create timer, which uses the REALTIME clock */
    if (timer_create(CLOCK_REALTIME, &sig_spec, &timer_h) < 0)
        perror('timer create');

    /* set the initial expiration and frequency of timer */
    tmr_setting.it_value.tv_sec = 1;
    tmr_setting.it_value.tv_nsec = 0;
    tmr_setting.it_interval.tv_sec = num_secs;
    tmr_setting.it_interval.tv_nsec = num_nsecs;
    if (timer_settime(timer_h, 0, &tmr_setting, NULL) < 0)
        perror('settimer');

    /* wait for signals */
    sigemptyset(&allsigs);
    while (1)
    {
        sigsuspend(&allsigs);
    }
}
```

**POSIX – example: setting a timer (cont.)**

```c
/* routine that is called when timer expires */
void timer_intr(int sig, siginfo_t *extra, void *cruft)
{
    /* perform periodic processing and then exit */
}
```
/* create timer, which uses the REALTIME clock */
if (timer_create(CLOCK_REALTIME, &sig_spec, &timer_h) < 0)
    perror("timer create");
/* set the initial expiration and frequency of timer */
tmr_setting.it_value.tv_sec = 1;
tmr_setting.it_value.tv_nsec = 0;
tmr_setting.it_interval.tv_sec = num_secs;
tmr_setting.it_interval.tv_nsec = num_nsecs;
if (timer_settime(timer_h, 0, &tmr_setting,NULL) < 0)
    perror("settimer");
/* wait for signals */
sigemptyset(&allsigs);
while (1) {
    sigsuspend(&allsigs);
}
/* routine that is called when timer expires */
void timer_intr(int sig, siginfo_t *extra, void *cruft) {
    /* perform periodic processing and then exit */
}

Remember the Pearl timers?
After 30 min all 5 min during 1 hr activate help;
References for this chapter

[Burns01]
Alan Burns and Andy Wellings
*Real-Time Systems and Programming Languages*
Addison Wesley, third edition, 2001

[Edler01]
F. Edler, M. Kühne, E. Tegeler
*Noise temperature measurements for the determination of the thermodynamic temperature of the melting point of Palladium*  

[Peacock97]
G.R. Peacock
*Standards for temperature sensors*
web guide: www.temperatures.com/temps.html

All references and some links are available on the course page.

---

Physical coupling

First step to embed a system into the real world:

*Transform all kinds of physical phenomena into analogue voltages*

- e.g. speed, pressure, brightness, loudness, colour, force, humidity, distance, salinity, density, radioactivity, spectrograms, reflectivity, acceleration, deformation, ..., ..., ..., or: temperature

Measuring temperature

Some observable effects of temperature changes:
- Mean square noise voltage changes
- Volume changes (gas, liquids, metals)
- Thermovoltage
- Changes in conductors and semiconductors
- State changes: into solid, liquid, or gaseous
Thermoelements

\[ E_{th} = K \cdot \text{grad}(T) \] with \( K \): Seebeck coefficient (depends on material)

\[ U_{th} = \int_0^L E_{th} \, dl = \int_0^L K \cdot \text{grad}(T) \, dl = K(T_0 - T_L) \]

Temperature measurement

- Take two wires (A & B) with different Seebeck coefficients \( K_A \) & \( K_B \)
- Connect them on one side
- Place the connected side in one temperature zone, and the open ends in another
- Measure the voltage difference \( U_{th} \)

\[ U_{th} = K_A(T_1 - T_2) + K_B(T_2 - T_1) = (K_A - K_B)(T_1 - T_2) \] with \( T_1 \) known ⇒ \( T_2 \)

Some standard combinations: (typical shape: ⬇️)

<table>
<thead>
<tr>
<th>Short Name</th>
<th>Material</th>
<th>( T_{max} )</th>
<th>( U_{th} ) with 0° to ( T_{max} )</th>
<th>( K_A - K_B )</th>
</tr>
</thead>
<tbody>
<tr>
<td>(T)</td>
<td>Cu-Constantan</td>
<td>400°C</td>
<td>21.000 mV</td>
<td>42.5×10⁻⁶</td>
</tr>
<tr>
<td>(J)</td>
<td>Fe-Constantan</td>
<td>700°C</td>
<td>39.720 mV</td>
<td>53.7×10⁻⁶</td>
</tr>
<tr>
<td>(K)</td>
<td>NiCr-Ni</td>
<td>1000°C</td>
<td>41.310 mV</td>
<td>41.1×10⁻⁶</td>
</tr>
<tr>
<td>(S)</td>
<td>PtRh-Pt</td>
<td>1300°C</td>
<td>13.138 mV</td>
<td>6.43×10⁻⁶</td>
</tr>
</tbody>
</table>
Physical coupling: Measuring temperature (thermoelements)

Pro:
• accepts high temperatures
• small
• relatively cheap

Contra:
• requires stable amplifier
• temperature differences only
• cables to the sensor need to have the same Seebeck coefficient

Applications of standard thermocouples

• (TYPE N) Nicrosil-Nisil thermocouples are suitable for use in oxidizing inert or dry reducing atmospheres. Must be protected from sulphurous atmospheres. Very accurate at high temperatures. Virtually the same emf (electromotive force) and range as Type K.

• (TYPE J) Iron-Constantan thermocouples are suitable for use in vacuum, oxidizing, reducing or inert atmospheres. Suitable for measuring temperatures up to 760°C for largest wire size.

• (TYPE K) Chromel-Alumel thermocouples are suitable for continuous use in oxidizing or inert atmospheres up to 1260°C for largest wire size. Because their oxidation resistance characteristics are better than those of other base metal thermocouples, they find widest use at temperatures above 538°C.

• (TYPE E) Chromel-Constantan thermocouples are suitable for use up to 781°C in oxidizing or inert atmospheres for largest gauge wires. Type E thermocouples develop the highest emf per degree of all commonly used thermocouples.

Applications of standard thermocouples (cont.)

• (TYPE T) Copper-Constantan thermocouples are suitable for subzero temperatures with an upper temperature limit of 371°C and can be used in vacuum, oxidizing, reducing or inert atmospheres.

• (TYPE R) Platinum 13% Rhodium - Platinum thermocouples are suitable for continuous use in oxidizing or inert atmospheres at temperatures up to 1482°C.

• (TYPE S) Platinum 10% Rhodium - Platinum thermocouples are suitable for continuous use in oxidizing or inert atmospheres at temperatures up to 1482°C.

• (TYPE B) Platinum 30% Rhodium - Platinum 6% Rhodium thermocouples are suitable for continuous use in oxidizing or inert atmospheres and short-term use in vacuum atmospheres at temperatures up to 1705°C.

• (TYPE W) Tungsten - Rhenium Alloy thermocouples are used to measure temperatures up to 2760°C. These thermocouples have inherently poor oxidation resistance and should be used in vacuum, hydrogen or inert atmospheres.

Linearity of standard thermocouples
Physical coupling: Measuring temperature (thermoelements)

Three standard forms

- grounded (G) junction
- ungrounded (U) junction
- exposed (E) junction

source: Alltemp Sensors

Physical coupling: Measuring temperature (resistors)

Thermoresistors: Pt$_{100}$ sensor

- e.g. in thinfilm technology:

  - Close-ceramic passivation layer
  - Connection pads
  - Pt/Al thin film
  - Glass or ceramics tube

  source: sensor-nite

- also common as a Platinum wire around a glass or ceramics tube

Some characteristics of the Pt$_{100}$ sensor:

- general: $R_T = R_0[1 + A(T - T_0) + B(T - T_0)^2 + C(T - T_0)^3 + ...]$  
- Platinum: $A = 3.27 \times 10^{-3} \frac{1}{K}$ (550°C) .. $4.2 \times 10^{-3} \frac{1}{K}$ (-150°C)  
- $Pt_{100}$: $0°C \leftrightarrow R_T = 100\, \Omega \pm 0.1\, \Omega \leftrightarrow \pm 0.26\, K$
- correction required to compensate for non-linearity  
- range approx. -200°C .. +650°C
- speed: 0.1s in flowing water .. multiple seconds in still air

Application of the Pt$_{100}$ sensor:

The problem of heating the sensor itself

- Pt$_{100}$ in a TO18 enclosure:

  $R_{th} = \frac{T_{Pt_{100}} - T_E}{P_V} = 480°C/W$ in non-moving air

- limiting the sensor error to 0.5°C around 0°C ($R_T = 100\, \Omega$)

  $R_{th} = \frac{T_{Pt_{100}} - T_E}{P_V} \rightarrow P_V = \frac{\Delta T}{R_{th}} \rightarrow U_{max} = \frac{\Delta T \cdot R_T}{R_{th}} = 0.323 \, V$
More things to consider for the $Pt_{100}$ sensor:

- compensate for the non-linearity of sensor itself
- choose an adequate (bridge) circuit, e.g. a ‘four wire’ setup:
- limit the cable length
- keep the cables on the same temperature
- limit the influence of the sensor to the environment

In relation to thermocouples:

Pro:

- higher accuracy 
  (if all side-effects have been taken care of)
- non-linearity less worse than with thermocouples
- long term stability
- measures absolute temperatures

Contra:

- smaller temperature range  
  -200°C .. +650°C  
  (some thermocouples exceed +2000°C)
- more expensive
- less robust (can be compensated by probe construction)
- usually bigger

Temperature sensitive semi-conductors (Thermistors):

Pro:

- very high accuracy  
  (with special models, if extensive compensation is applied)
- usually cheaper than $Pt_{100}$
- high long term stability  
  (with special models)
- large changes with temperature

Contra:

- even smaller temperature range  
  -40°C .. +350°C
- terrible non-linearities
- standard components are very inaccurate and instable
- larger than thermocouples and $Pt_{100}$

In relation to thermocouples and $Pt_{100}$:

Pro:

- higher accuracy 
  (with special models, if extensive compensation is applied)
- usually cheaper than $Pt_{100}$
- high long term stability  
  (with special models)
- large changes with temperature

Contra:

- even smaller temperature range  
  -40°C .. +350°C
- terrible non-linearities
- standard components are very inaccurate and instable
- larger than thermocouples and $Pt_{100}$
Physical coupling: Measuring temperature (noise)

Noise temperature measurement

Based on Nyquist formula: $\overline{U^2} = 4kTR\Delta f$

with $k$: Boltzman constant, $T$: thermodynamic temperature, $R$: electric resistance, and $\Delta f$: the measurement bandwidth

Pro:
- wide range
- high accuracy
- long term stability

Contra:
- expensive
- sophisticated amplifier setups

An actual device: range: $1..2500^\circ K$, accuracy: ±0.1% (over the full range)

More ways to measure temperature:
- Spreading resistors
- Piezos and other temperature sensitive quartz elements
- Temperature controlled current sources (e.g. AD590)
- Mercury filled thermometers
- …

Basic conclusions

☞ we just scratched the surface of conversion methods for ONE physical value (temperature).
☞ converting physical phenomena into analogue voltages seems to be a complex matter
  … in fact a whole industry is dedicated to this field exclusively.
☞ always ask for the full sensor specifications (and read them).
☞ never assume that the output is a linear translation of a physical value.

Physical coupling is not the only loss afflicted stage of conversion, but it is usually the most complex one
Range and relative speed measurements

- **Triangulation** (optical)
- **Time of flight** (optical, acoustical, electro-magnetic)
  - Phase correlation (optical, acoustical, electro-magnetic)
- **Intensity** (optical, acoustical)
- **Doppler methods** (acoustical, electro-magnetic)
- **Interferometry** (optical, electro-magnetic)

**Range measurements – Triangulation**

\[ d = b \cdot \tan \alpha \]

- non-linear, very poor resolution, if \( d \approx b \)
- highly focused light-beam required \( \rightarrow \) laser \( \rightarrow \) safety considerations
- distances are measured along the optical axis only.

**Range measurements – Time of flight - Phase correlation**

Method: measure the time of flight between the outgoing signal and the received reflected signal.

\( \rightarrow \) in case of light, this method requires high resolution timers (> 1GHz)

- Method is perfectly linear.
- The achieved resolution depends on the precision of the signal’s rising edge and the resolution of employed timers.
- Signals can be formed and volume measurements are possible.

\( \rightarrow \) in order to increase the resolution, the outgoing signals are often modulated and the phase shifts between outgoing and reflected signals are detected.
in the special case of one ultrasound transducer for sending and receiving (like above), there is a short time delay before the transducer is ready to receive a signal (oscillations need to die away first).

---

### Range measurements: Laser

a common laser range finder (SICK):
- range: max. 80 m
- angular resolution: 0.25°
- response time: max. 53 ms
- resolution: 10 mm
- accuracy: typ. 5-10 mm

---

### Speed measurements: Doppler current profilers

Physical effect: Doppler shift frequency \( f_d = -2 f_s \frac{V}{c} \)

with source frequency \( f_s \), relative velocity \( v \) and signal speed \( c \).

a common current profiler (SonTek):
- ping signal: 250 kHz - 3 MHz
- range: up to 160 m; velocity: ±10 m/s
- resolution: 0.15 - 2 m; 1 mm/s
- accuracy: ±1%
- blanking zone: 0.2 - 2.0 m
Summary

Physical coupling

- Physical phenomena
- Measuring temperature
  - thermoelements, thermocouples
  - thermoresistors
  - thermistors
  - noise temperature measurement and others
- Measuring range and relative speed
  - triangulation
  - time of flight
  - intensity
  - Doppler methods
  - interferometry
- Examples: time-of-flight ultrasound & laser, Doppler current profiler

References for this chapter

[Burns01]
Alan Burns and Andy Wellings
Real-Time Systems and Programming Languages
Addison Wesley, third edition, 2001

[Kester93]
Walt Kester, James Bryant, Joe Buxton
High Resolution Signal Conditioning ADCs
System Applications Guide, Analog Devices, 1993, Section 3

[M_MPC565/566_00]
Motorola
MPC565/MPC66 microcontroller
Reference manual, 2000

[NS_ADC08200_02]
National Semiconductor
ADCO8200 – 8-Bit, 20 MSPS to 200 MSPS, 1.05 mW/MSPS A/D Converter
Datasheet, 2002

[NS_LM12L458_99]
National Semiconductor
LM12L458 – 12-Bit + Sign Data Acquisition System with Self-Calibration
Datasheet, 1999

all references and links are available on the course page
Sample data with a frequency $f_s$.

Interpolation suggests an original signal like

Sample data with a higher frequency $f_s$.

suggests an original signal like

the phenomenon of the observed signal at a lower $f_s$ is called “aliasing”.
An analog signal with a bandwidth of \( f_a \) must be sampled at 
\[ f_s > 2f_a \]
in order to avoid loss of information.

Analog signal with a bandwidth of \( f_a \) must be sampled at 
\[ f_s > 2f_a \]
⇒ theoretical sufficient: \[ \square \] | practical case: \[ \square \] (oversampling)

A resolution of \( N \) bit gives \( 2^N \) possible discrete output levels:
- smallest distinguishable value (least significant bit or LSB):
  \[ q = \frac{1}{2^N} \]
- ratio \( 1/q \) expressed in dB:
  \[ 10\log(2^{2N}) = N \cdot 20\log2 \approx N \cdot 6.02\,\text{db} \]

The mean square error over one step and the rms noise voltage is:

\[ E^2 = \frac{1}{q} \int_{-q/2}^{q/2} E^2\,dE = \frac{q^2}{12} \]

rms noise voltage:

\[ \frac{q}{\sqrt{12}} \]
The signal $S$ with respect to the rms error is thus:

$$S = \frac{q}{\sqrt{12}} = S \cdot 2^N \sqrt{12}$$

or as the “Signal to Noise Ratio”:

$$SNR[db] = 10\log \left( \frac{S^2}{q^2/12} \right)$$

Assuming an ideal input signal:

$$F(t) = A\sin \omega t ; q = \frac{2A}{2^N}$$

then

$$SNR[db] = 10\log \left( \frac{A^2}{2} \cdot \frac{q^2}{12} \right)$$

$$SNR[db] = 10\log \left( \frac{A^2}{2} \cdot \frac{A^2}{3 \cdot 2^{2N}} \right)$$

$$SNR[db] = 10\log \left( \frac{3 \cdot 2^{2N}}{2} \right)$$

$$\approx SNR[db] = N \cdot 6.02 + 1.76$$
Determining the effective number of bits (ENOB):

\[ SNR_{\text{ideal (dB)}} = 20N\log 2 + 10\log \frac{3}{2} \]

\[ ENOB = \frac{SNR_{\text{actual}} - 10\log (3/2)}{20\log 2} \]

\[ ENOB = \frac{SNR_{\text{actual}} - 1.76}{6.02} \]

(ENOB = N for an ideal A/D converter)

Actual A/D converters also have:

- **Missing codes**
  (reduce SNR by \(20\log 2\) or 6.02\(\text{db}\) for each missing code)

- **Differential non-linearity (DNL):**
  differences between successive code-widths

- **Integral non-linearity (INL):**
  maximal difference between code-centres and the ideal line.

Some central criteria for A/D converters:

- Throughput (max. sampling frequency)
- Resolution (ENOB)
- Latency
- Accuracy

Trade-off between: Throughput and Resolution | Accuracy

A/D converter types

- **Flash converter**
  - \(2^N - 1\) comparators are directly connected to the input signal \(A_{IN}\)
  - Accuracy depends on the accuracy of the reference voltages only
  - Conversion in **one cycle**

but

→ a straight 16-bit ADC flash converter would require: 65535 comparators and references
Pipelined flash converter

- divides the conversion in multiple stages with an
  \( m \) bit ADC and DAC each.
- digital conversion results are accumulated \((D)\)
- Analog residues are forwarded \((A)\)

\[ \text{with } p \text{ } m\text{-bit pipeline stages the converter achieves a resolution of:} \]
\[ n = pm \text{ bits} \]

Successive approximation shift register (SAR)

- Compare \( A \) successively with \((\text{MSB to LSB}): \frac{1}{2}, \frac{1}{4}, \frac{1}{8}, \ldots, \frac{1}{2^N}\) of full scale
- Reduce the input in each step according to already converted bits (and shift the register).
- Conversion takes \( N \) steps
  - Only one comparator required

\[ \text{2nd order } \Sigma - \Delta \text{ A/D Converters} \]

- The output of a 1-bit DAC \((\pm U_{ref})\) is subtracted from the input and integrated.
- The output of the integrator is compared against signal ground.
- According to the comparator the output of the 1-bit DAC changes.
  - This is done at a very high rate!
  - The density of ‘1’s in the bit-stream represents the signal level of \( A \)

- Converter is inherently linear.
- Due to the high ‘oversampling’ rates, relatively slow.
- High resolutions (typically 16-24 bits).
- Higher order S-D converters achieve an even better signal to noise ratio
Relation between resolution (ENOB, SNR) and oversampling in higher order Σ-Δ Converters (qualities):

<table>
<thead>
<tr>
<th>Oversampling</th>
<th>4</th>
<th>8</th>
<th>16</th>
<th>32</th>
<th>64</th>
<th>128</th>
<th>256</th>
</tr>
</thead>
<tbody>
<tr>
<td>1st order</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>2nd order</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>3rd order</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

Comparison between these four A/D converters

<table>
<thead>
<tr>
<th></th>
<th>Flash</th>
<th>Pipelined Flash</th>
<th>SAR</th>
<th>Σ-Δ</th>
</tr>
</thead>
<tbody>
<tr>
<td>Throughput</td>
<td>1</td>
<td>2</td>
<td>3</td>
<td>4</td>
</tr>
<tr>
<td>Resolution</td>
<td>4</td>
<td>3</td>
<td>2</td>
<td>1</td>
</tr>
<tr>
<td>Latency</td>
<td>1</td>
<td>3</td>
<td>2</td>
<td>4</td>
</tr>
<tr>
<td>Accuracy</td>
<td>2</td>
<td>2</td>
<td>2</td>
<td>1</td>
</tr>
</tbody>
</table>

1 = best

ADC08200 – basic characteristics

- Resolution: 8 bit
- Sampling frequency: 10 - 200 (230) MHz
- Differential non-linearity (DNL): ±0.4 LSB (typical), ±0.95 LSB (max.)
- ENOB: 7.5 (at 4 MHz), 7.3 (at 50 MHz), 7.0 (at 100 MHz)
- No missing codes
- Power consumption: 1.05 mW/MSPS, 1 mW (power down)
- Latency: 6 cycles (pipeline delay)
- Aperture (sampling delay): 2.6 ns, with 2 ps rms jitter

Pipelined Flash converter with 2 stages and 2 pipelines
ADC08200 – timing

![Timing Diagram]

ADC08200 – integral non-linearity

![INL Graph]

LM12L458

12-Bit + sign, 8 channel, A/D converter, controller and interface

Controller features:
- Programmable acquisition times and conversion rates
- 32-word conversion FIFO
- Self-calibration and diagnostic mode
- 8- or 16-bit wide data bus microprocessor or DSP

Typ. applications:
- Data Logging
- Process Control
### LM12L458 – accessible registers

<table>
<thead>
<tr>
<th>A4 A3 A2 A1</th>
<th>Purpose</th>
<th>Type</th>
<th>D15 D14 D13 D12 D11 D10 D9 D8 D7 D6 D5 D4 D3 D2 D1 D0</th>
</tr>
</thead>
<tbody>
<tr>
<td>0 0 0 0</td>
<td>Instruction RAM</td>
<td>R/W</td>
<td>Acquisition Time</td>
</tr>
<tr>
<td>0 to 1</td>
<td></td>
<td></td>
<td>Watchdog</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>&amp;×2 Timer Sync</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>VA, VA</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>Paused Loop</td>
</tr>
<tr>
<td>0 0 0 0</td>
<td>Instruction RAM</td>
<td>R/W</td>
<td>Acquisition Time</td>
</tr>
<tr>
<td>0 to 1</td>
<td></td>
<td></td>
<td>Don't Care</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>&gt;× Sign</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>Limit #1</td>
</tr>
<tr>
<td>0 0 0 0</td>
<td>Instruction RAM</td>
<td>R/W</td>
<td>Acquisition Time</td>
</tr>
<tr>
<td>0 to 1</td>
<td></td>
<td></td>
<td>Don't Care</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>&gt;× Sign</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>Limit #2</td>
</tr>
<tr>
<td>0 0 0 0</td>
<td>Configuration Register</td>
<td>R/W</td>
<td>Don't Care</td>
</tr>
<tr>
<td>0 to 1</td>
<td></td>
<td></td>
<td>D0</td>
</tr>
<tr>
<td>0 0 0 0</td>
<td>Interrupt Enable Register</td>
<td>R/W</td>
<td>Number of Conversions</td>
</tr>
<tr>
<td>0 to 1</td>
<td></td>
<td></td>
<td>in Conversion FIFO</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>to Generate INT2</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>INT7 INT4 INT5 INT2 INT1</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>INT0</td>
</tr>
<tr>
<td>0 0 0 0</td>
<td>Interrupt Status Register</td>
<td>R</td>
<td>Actual Number of</td>
</tr>
<tr>
<td>0 to 1</td>
<td></td>
<td></td>
<td>Conversion Results</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>in Conversion FIFO</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>Address of</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>INST7 INST5 INST4 INST3</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>INST2 INST1 INST0</td>
</tr>
<tr>
<td>0 0 0 0</td>
<td>Timer Register</td>
<td>R/W</td>
<td>Timer Preset High Byte</td>
</tr>
<tr>
<td>0 to 1</td>
<td></td>
<td></td>
<td>Timer Preset Low Byte</td>
</tr>
<tr>
<td>0 0 0 0</td>
<td>Conversion Register</td>
<td>R</td>
<td>Address</td>
</tr>
<tr>
<td>0 to 1</td>
<td></td>
<td></td>
<td>Sign</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>Conversion</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>Conversion Data</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>LSBS</td>
</tr>
<tr>
<td>0 0 0 0</td>
<td>Limit Status Register</td>
<td>R</td>
<td>Limit #2: Status</td>
</tr>
<tr>
<td>0 to 1</td>
<td></td>
<td></td>
<td>Limit #1: Status</td>
</tr>
</tbody>
</table>

### LM12L458 – instruction RAM

<table>
<thead>
<tr>
<th>A4 A3 A2 A1</th>
<th>Purpose</th>
<th>Type</th>
<th>D15 D14 D13 D12 D11 D10 D9 D8 D7 D6 D5 D4 D3 D2 D1 D0</th>
</tr>
</thead>
<tbody>
<tr>
<td>0 0 0 0</td>
<td>Instruction RAM</td>
<td>R/W</td>
<td>Acquisition Time</td>
</tr>
<tr>
<td>0 to 1</td>
<td></td>
<td></td>
<td>Watchdog</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>&amp;×2 Timer Sync</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>VA, VA</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>Paused Loop</td>
</tr>
<tr>
<td>0 0 0 0</td>
<td>Instruction RAM</td>
<td>R/W</td>
<td>Acquisition Time</td>
</tr>
<tr>
<td>0 to 1</td>
<td></td>
<td></td>
<td>Don't Care</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>&gt;× Sign</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>Limit #1</td>
</tr>
<tr>
<td>0 0 0 0</td>
<td>Instruction RAM</td>
<td>R/W</td>
<td>Acquisition Time</td>
</tr>
<tr>
<td>0 to 1</td>
<td></td>
<td></td>
<td>Don't Care</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>&gt;× Sign</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>Limit #2</td>
</tr>
<tr>
<td>0 0 0 0</td>
<td>Configuration Register</td>
<td>R/W</td>
<td>Don't Care</td>
</tr>
<tr>
<td>0 to 1</td>
<td></td>
<td></td>
<td>D0</td>
</tr>
<tr>
<td>0 0 0 0</td>
<td>Instruction RAM</td>
<td>R/W</td>
<td>Number of Conversions</td>
</tr>
<tr>
<td>0 to 1</td>
<td></td>
<td></td>
<td>in Conversion FIFO</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>to Generate INT2</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>INT7 INT4 INT5 INT2 INT1</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>INT0</td>
</tr>
<tr>
<td>0 0 0 0</td>
<td>Interrupt Enable Register</td>
<td>R/W</td>
<td>Number of Conversions</td>
</tr>
<tr>
<td>0 to 1</td>
<td></td>
<td></td>
<td>in Conversion FIFO</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>to Generate INT2</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>INT7 INT4 INT5 INT2 INT1</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>INT0</td>
</tr>
<tr>
<td>0 0 0 0</td>
<td>Interrupt Status Register</td>
<td>R</td>
<td>Actual Number of</td>
</tr>
<tr>
<td>0 to 1</td>
<td></td>
<td></td>
<td>Conversion Results</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>in Conversion FIFO</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>Address of</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>INST7 INST5 INST4 INST3</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>INST2 INST1 INST0</td>
</tr>
<tr>
<td>0 0 0 0</td>
<td>Timer Register</td>
<td>R/W</td>
<td>Timer Preset High Byte</td>
</tr>
<tr>
<td>0 to 1</td>
<td></td>
<td></td>
<td>Timer Preset Low Byte</td>
</tr>
<tr>
<td>0 0 0 0</td>
<td>Conversion Register</td>
<td>R</td>
<td>Address</td>
</tr>
<tr>
<td>0 to 1</td>
<td></td>
<td></td>
<td>Sign</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>Conversion</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>Conversion Data</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>LSBS</td>
</tr>
<tr>
<td>0 0 0 0</td>
<td>Limit Status Register</td>
<td>R</td>
<td>Limit #2: Status</td>
</tr>
<tr>
<td>0 to 1</td>
<td></td>
<td></td>
<td>Limit #1: Status</td>
</tr>
</tbody>
</table>

Every entry in the instruction RAM consists of:

- **Loop** (1bit): indicates the last instruction and branches to the first one.
- **Pause** (1bit): halts the sequencer before this instruction.
- **VIN+**, **VIN-** (2*3bit): select the input channels (000 selects ground in VIN.)
- **Sync** (1bit): wait for an external sync. signal before this instruction.
- **Timer** (1bit): wait for a preset 16-bit counter delay before this instruction.

**Type ChannelPlus** is (Ch0, Ch1, Ch2, Ch3, Ch4, Ch5, Ch6, Ch7);
**Type ChannelMinus** is (Gnd, Ch1, Ch2, Ch3, Ch4, Ch5, Ch6, Ch7);
**Type Resolution** is (12bit, 8bit);
**Type Acquisition** is (15; 9+2D, 2+2D (8bit))

```
--9+2D (12bit), 2+2D (8bit)
```
Units_Per_Word : constant Integer := Word_Size / Storage_Unit;
for Instruction use record
  EndOfLoop at 0*Units_Per_Word range 0..0;
  Pause at 0*Units_Per_Word range 1..1;
  Vplus at 0*Units_Per_Word range 2..4;
  Vminus at 0*Units_Per_Word range 5..7;
  Sync at 0*Units_Per_Word range 8..8;
  Timer at 0*Units_Per_Word range 9..9;
  Resolution at 0*Units_Per_Word range 10..10;
  Watchdog at 0*Units_Per_Word range 11..11;
  AcquisitionTime at 0*Units_Per_Word range 12..15;
end record;

ADC_Instructions (0) := (EndOfLoop => False,
                          Pause => False,
                          Vplus => Ch0,
                          Vminus => Gnd,
                          ... => EightBit,
                          Watchdog => False,
                          AcquisitionTime => 10);
ADC_Instructions (1) := (EndOfLoop => True, -- last instruction
                          Pause => False,
                          Vplus => Ch1,
                          Vminus => Ch2,
                          Sync => True,
                          Timer => False,
                          Resolution => EightBit,
                          Watchdog => False,
                          AcquisitionTime => 10);

Data structures in ‘C’:
enum ChannelPlus {Ch0=0, Ch1, Ch2, Ch3, Ch4, Ch5, Ch6, Ch7};
enum ChannelMinus {Gnd=0, Ch1, Ch2, Ch3, Ch4, Ch5, Ch6, Ch7};
enum Resolutions {TwelveBit=0, EightBit};
struct {
  unsigned int EndOfLoop  : 1;
  unsigned int Pause       : 1;
  ChannelPlus    Vplus     : 3;
  ChannelMinus   Vminus    : 3;
  unsigned int Sync     : 1;
  unsigned int Timer    : 1;
  Resolutions   Resolution : 1;
  unsigned int Watchdog : 1;
  unsigned int AcquisitionTime : 4;
} Instruction;
Data structures in ‘C’:

```c
struct {
  unsigned int EndOfLoop : 1;
  unsigned int Pause : 1;
  ChannelPlus Vplus : 3;
  ChannelMinus Vminus : 3;
  unsigned int Sync : 1;
  unsigned int Timer : 1;
  Resolutions Resolution : 1;
  unsigned int Watchdog : 1;
  unsigned int AcquisitionTime : 4;
} Instruction;
```

```c
Instruction InstructionsA[8];
InstructionsA *Instructions = 0x0000132D;
```

If this works, you were lucky two times:
- The compiler implemented the struct-fields in the intended places and order.
- The bit ordering in your device is the way the compiler assumed it.

Macro-Assembler style programming:

In order to produce portable code in ‘C’, it is necessary to set bits manually:

```c
unsigned int setbits (unsigned int *r, unsigned int n, unsigned int p, unsigned int x)
{   unsigned int mask;
    mask  = ~(~0 << n);
    *r   &= ~(mask << p);   *r   |= (x & mask) << p;   return (*r);
}
```

The configuration register consists of (1bit each):

- **Start**: starts the sequencer.
- **Reset**: sets the instruction pointer to ‘000’.
- **Auto-Zero**: triggers a ‘short’ calibration (76 cycles – 1 offset sample).
- **Full-Cal**: initiates a full calibration (4944 cycles, 8 samples) interrupt.
- **Stand-by**: disconnects the external clock, preserves the registers.
  After powering up again (≈ 10 ms): a specific interrupt is issued.
- **Chan-Mask**: format selection for the FIFO output registers.
- **Auto-Zero<sub>ec</sub>**: auto-zeros the ADC automatically in every conversion.
The configuration register consists of (1 bit each) (cont.):

- **I/O Sel**: sets the Sync pin to input or output mode.
- **RAM Pointer**: selects the current (16-bit) part in each 48-bit instruction.
- **Test=0**: production testing mode: leave this bit at ‘0’.
- **DIAG**: connects $V_{IN^+}$ and $V_{IN^-}$ to $V_{REF^+}$ and $V_{REF^-}$ for testing purposes.

The conversion FIFO is accessible via one address only:

- Every read on this address will delete this result from the internal memory and shift the next one into the visible conversion FIFO register.
- The FIFO holds 32 conversion results max.
- **Data will be lost**, if the results are not read fast enough to prevent a buffer overrun.

The controller can issue specific interrupts or initiate a DMA transfer, when a given number of results are accumulated or a certain instruction is completed.

☞ on the CPU side: **Polling**, **interrupt** or DMA scheme
**μControllers**

**MC68HC05**

- **Clock:** max. 2.1MHz internal (4.2MHz external)
- **RAM:** 176 bytes
- **ROM:** 5936 bytes
- **EEPROM:** 256 bytes
- **Power saving modes** (stop, wait, slow)
- **Serial:** 46-76800 baud (at 2.4576MHz)
  79-131072 baud (at 4.194394MHz)
- **Parallel I/O:** 3*8bit; Parallel in: 1*8bit
- **Timers:** 1*16bit
- **A/D:** 8 channels, 8bit
- **PWM:** 2 generators

**µControllers**

**MPC565**

- **-40° - +125°C,** power dissipation: 0.8 - 1.12W
- **CPU:** PowerPC core (incl. FPU & BBC), 40/56MHz
- **RAM:** flash: 1M, static: 36K
- **Time processing units:** 3 (via dual-ported RAM)
- **Timers:** 22 channels (PWM & RTC supported)
- **A/D convertors:** 40 channels, 10bit, 250kHZ
- **Can-bus:** 3 TOUCAN modules
- **Serial:** 2 interfaces
- **Data link controller:**
  SAE J1850 class B communications module
- **Real-time embedded application development interface:** NEXUS debug port (IEEE-ISTO 5001-1999)
- **Packaging:** 352/388 ball PBGA
a special-purpose µcontroller:
- Independent µengine.
- 16 digital I/O channels with independent match and capture capabilities.
- Meant to operate these I/O channels for timing control purposes.
- Predefined µengine command set (ROM functions in control store).
- 2 16-bit time bases

The period/pulse-width accumulator (PPWA) algorithm accumulates a 16-bit or 24-bit sum of either the period or the pulse width of an input signal over a programmable number of periods or pulses (from 1 to 255).

The stepper motor (SM) control algorithm provides for linear acceleration and deceleration control of a stepper motor with a programmable number of step rates of up to 14.

The PSP function generates pulses of variable length at specified "angles." Angle clock period is measured (in TCR1 clocks) using the PMA/PMM function on another channel.

This function measures the period (in TCR1 clocks) between regularly occurring input transitions and makes this period available for use by other functions or by the CPU (optional detection of misses and additional transitions).

The TPU can generate a pulse-width modulation (PWM) waveform with any duty cycle from zero to 100% (within the resolution and latency capability of the TPU).

Three different operating modes allow the function to maintain complex timing relationships between channels without CPU intervention.

QDEC uses two channels to decode a pair of out-of-phase signals in order to present the CPU with directional information and a position value.

Refer the control store of the µengine to the dual-ported RAM instead of the integrated ROM area and supply:
- up to 16 µengine commands (functions)
- in 2-8Kbyte of long-word (32-bit) organized memory
- programmed in a 32-bit µinstruction format (explained next)

the dual-ported RAM is now cut off from the CPU (the TPU parameter RAM is not affected)
µControllers MPC565 – TPU: the µinstruction formats:

### 1: Execution unit and RAM

- **RW**: Read/Write Control
- **T1ABS**: A-Bus Source Control
- **T3ABD**: A-Bus Destination Control
- **SHF**: Shifter Control
- **SRC**: Shift Register Control
- **CCL**: Condition Code Latch Control
- **T1BBS**: B-Bus Source Control
- **CIN**: B-Bus Carry Control
- **BINV**: B-Bus Invert Control
- **IOM**: Input/Output Mode Control
- **AIS**: Address
- **DEC/END**: Decrementor / End Control

### Operation groups:

- **Execution unit**
- **Channel control**
- **RAM**
- **Sequencer**

### Operation groups:

- **Execution unit**
- **Channel control**
- **RAM**
- **Sequencer**

---

### 2: Execution unit, flag, and channel control

- **ERW**: Event Register Write Control
- **T1ABS**: A-Bus Source Control
- **T3ABD**: A-Bus Destination Control
- **SHF**: Shifter Control
- **SRC**: Shift Register Control
- **CCL**: Condition Code Latch Control
- **T1BBS**: B-Bus Source Control
- **CIN**: B-Bus Carry Control
- **BINV**: B-Bus Invert Control
- **PAC**: Pin Action Control
- **LSL**: Link Service Latch Negation Control
- **PSC**: Pin State Control
- **FLC**: Flag Control
- **CIR**: Channel Interrupt Request
- **DEC/END**: Decrementor / End Control

### Operation groups:

- **Execution unit**
- **Channel control**
- **RAM**
- **Sequencer**

---

### 3: Conditional branch, flag, and channel control

- **BCC**: Branch Condition Code Field
- **FLS**: µPC Flush Control
- **BAF**: Branch Address Field
- **TBS**: Time Base Select Control
- **PAC**: Pin Action Control
- **BCF**: Branch Condition Control
- **PSC**: Pin State Control
- **FLC**: Flag Control
- **CIR**: Channel Interrupt Request
- **MTSR**: Match/Transition Detect Service Request Inhibit Control

### Operation groups:

- **Execution unit**
- **Channel control**
- **RAM**
- **Sequencer**

---

### 4: Jump, flag, and RAM

- **RW**: Read/Write Control
- **NMA**: Next µPC Address Mode Control
- **FLS**: µPC Flush Control
- **BAF**: Branch Address Field
- **FLC**: Flag Control
- **LSL**: Link Service Latch Negation Control
- **IOM**: Input/Output Mode Control
- **AIS**: Address
- **DEC/END**: Decrementor / End Control

### Operation groups:

- **Execution unit**
- **Channel control**
- **RAM**
- **Sequencer**
µControllers MPC565 – TPU: the µinstruction formats:

- T1ABS: T1 A-Bus Source Control
- T3ABD: T3 A-Bus Destination Control
- SHF: AU Shifter Control
- SRC: Shift Register Control
- CCL: AU Condition Code Latch Control
- T1BBI: T1 B-Bus Immediate Data
- LSL: CC Link Service Latch Negation Control
- EQ/GE: CC Match Compare Register Control
- FLC: CC Flag Control
- CIR: CC Channel Interrupt Request
- DEC/END: SEQ Decrementor / End Control

Entities to consider:

- **States**: non-interruptible µcode-blocks
- **Functions**: constructed of one or multiple states
- **Channels**: digital I/O lines with 16-bit match and capture
- **Priorities**: of channels
- **Timers**: 2 16-bit time-bases

☞ Associate functions, time-bases, channels and priorities!

Function are composed of one or multiple “states”
☞ States are atomic instructions blocks!

Ready to go …

… and let it run …

Time slots for states on three priorities

☞ no state will be suspended permanently
Free time slots are transferred:

Free slots are made available in priority order

Channels on the same level are scheduled by their channel-number

States have variable lengths

a set of influences need to be considered to calculate latencies

Latencies

... for latencies of capture and match at each channel mind:
- only the time-base resolution (all channels are evaluated independently and in parallel)

... for the functions associated with individual channels mind the:
- number of active channels (max. 16)
- number of channels on each priority level (add max. 2 µcycles for each “state-switch”)
- number of available time slots on each priority level per full scheduler-cycle (4, 2, 1 slot(s))
- number of µcycles to execute individual states of a function (2 µcycles per instruction)
- number of RAM accesses during the execution of a state (each access may stall for 2 CPU cycles)
- TPU clock cycle frequency

Determining latencies

assuming that all states may execute at all times!

calculate the number of µcycles per state
Determining latencies

- assuming that all states may execute at all times!
- calculate a schedule based on priorities

Add potential state switches (2 cycles)

Determine maximal latencies!

- assuming that all states may execute at all times!
- Add potential stall times for each RAM-access (2 µcycles)
Determining latencies

Assuming that all states on the same level may execute at all times!
On all other levels: the longest state is always executed!

• Handles up to 48 different sources
  (32 from internal modules, 8 from timers and clocks, and 8 external vectorized sources) and supply each of them with a unique interrupt-vector

• 8 interrupt levels are distinguished by the interrupt controller
  (32 interrupt levels are supplied by the internal modules, prioritized and vectorized interrupts are supplied by external sources)

• Latency: 20 clock cycles
  + bus collisions + CPU state saving + tasking system overhead!

Time-base & Real-time clock

• 64 bit time base
  driven by an external clock: e.g. 20MHz ↔ resolution: 50ns; range: ≈30000 years

• Free running
  (not influenced by any CPU action or resets)

• 2 reference registers
  (used for compares and interrupt generation)

• Real-Time clock supplies full seconds (32bit ↔ range: ≈136 years)
  (not affected by CPU, resets, and operates in all low-power modes)
Language requirements for interfaces

☞ Specify the device interface (protocol and formats) in **all detail**
   (candidates: Modula-2, Ada95, CHILL, ERLANG, …
   or Macro-Assemblers and C/C++ (if platform-independence is not required))

☞ Handling asynchronous messages from the device
   many different methods to implement a context-switch
   (candidates: all languages with some process-orientation:
   PEARL, CHILL, ERLANG, Ada95, RT-Java, POSIX, …)

Language requirements for interfaces

☞ Specify the device interface (protocol and formats) in **all detail**
   (candidates: Modula-2, Ada95, CHILL, ERLANG, …
   or Macro-Assemblers and C/C++ (if platform-independence is not required))

☞ Handling asynchronous messages from the device
   many different methods to implement a context-switch
   (candidates: all languages with some process-orientation:
   PEARL, CHILL, ERLANG, Ada95, RT-Java, POSIX, …)

☞ The term “**high-level languages**” in the real-time interface context:
   Allow for **abstractions while being specific**
   down to the actual level of interface realities

Basic sampling control mechanisms

• **Status driven**: the computer polls for information
  (used in dedicated µcontrollers and pre-scheduled hard real-time environments)

• **Interrupt driven**: The data generating device may issue an interrupt
  when new data had been detected / converted or when internal buffers are full
  • **Program controlled**: The interrupts are handled by the CPU directly
    (by changing tasks, calling a procedure, raising an exception,
    free tasks on a semaphore, sending a message to a task, …)
  • **Program initiated**: The interrupts are handled by a DMA-controller.
    No processing is performed. Depending on the DMA setup,
    cycle stealing can occur and needs to be considered for the worst case computing times.

• **Channel program controlled**: The interrupts are handled by a dedicated channel
  device. The data is transferred and processed. Optional memory-based communication
  with the CPU. The channel controller is usually itself a dedicated µengine / µcontroller.

Handling device responses

After setting up an interface device …

☞ … responses from devices can be:
  • immediate
  • with a constant delay
  • or within a defined time-frame
  • or unpredictable

☞ The device handler may thus:
  • perform a ‘busy-wait’ for the response
  • reschedule the device-process by a constant delay
  • schedule the device-process periodically and employ different time-slots for sending
    control / data and receiving status / data
  • activate / trigger / call / signal the device-process by interrupts
    from the specific device
Interfaces & Controllers

Handling device responses

☞ The central issue for hard real-time environments:

*How to embed the unpredictable in predictable systems?*

Obviously: a infinite flood of interrupts cannot be accepted, so:

- either the asynchronous events need to be synchronized with the remaining system-tasks or
- there are exclusive processing resources (e.g. a dedicated µcontroller) for a specific device.

Converters & Interfaces

- Analogue signal chain in a digital system
  - Sampling data ≠ aliasing ≠ Nyquist's criterion ≠ oversampling
  - Quantization (LSB, rms noise voltage, SNR, ENOB)
  - Missing codes, DNL, INL.

- A/D converters: flash, pipelined-flash, SAR, Σ-Δ, n-th order Σ-Δ

- Examples:
  - Fast and simple A/D converter example: National Semiconductor ADC08200
  - Multi-channel A/D data logging interface example: National Semiconductor LM12L458
  - Simple 8-bit µcontroller example: Motorola MC68HC05
  - Complex 32-bit µcontroller example: Motorola MPC565
    - TPU: µprogramming, atomic states, µengine scheduling, max. latency analysis
    - NEXUS debugging port

- General device handling / sampling control / language requirements
Notions of time and space

What is time? / What is embodiment?

Interfacing with time

Specifying timing requirements

Satisfying timing requirements

What is time?

Do we exist in time, or is time part of our existence?

☞ Is time an intrinsic property of nature? ⇒ Platonism

Time is an external phenomenon. Thus simultaneous events happen at the same exact absolute time. There is the underlying assumption that time is progressing, even if no changes can be observed.

☞ Is time a construct which is based on observable events? ⇒ Reductionism

Time is the observation of distinguishable events. If the observed events are ‘regular’, a useful time-reference can be constructed. If all possible observers detect one event before another one, they are said to be in sequence. If this cannot be assumed for all possible observers, they are said to be simultaneous. Therefore the notion of time is reduced to a notion of causality.

☞ Is time ‘linear’ between observable events?

A mathematical notion of time

- Transitivity: \( x < y \land y < z \Rightarrow x < z \)
- Linearity: \( x < y \lor x > y \Rightarrow x \neq y \)
- Irreflexivity: \( \neg(x < x) \)
- Density: \( x < y \Rightarrow \exists z \mid x < z \land z < y \)

☞ Real clocks have limited resolutions and are running asynchronously!

What is time? – A physical notion of time

- 1676: Rømer proofs the existence of the speed of light (and measures it).
- 1687: Newton’s “Principia Mathematica” assumes an absolute time, independent of space itself and independent of events.
- 1905: The concept of absolute time is destroyed first by Einstein and (a few weeks later) by Poincaré.
- 1915: Einstein’s general theory of relativity eliminates the independence of time (space) and events in time (space).

☞ One principal consequence for measurements of time:

Clocks under higher gravity or in faster observation frames are slower.

Practical consequences: clocks in satellites need to be adjusted accordingly.
Notions of time and space

What time is it?

| UT0 | 1884: Mean solar time at Greenwich meridian |
| UT1 | Corrected UT0, \(\approx\) polar motion |
| UT2 | Corrected UT1, \(\approx\) variations in the speed of rotation of the earth |
| IAT (International Atomic Time) | a caesium 133 atomic clock (current accuracies: one miss in \(10^{13}\) ticks, e.g. approximately once every 300000 years) |
| UTC (Universal Time Coordinated) | a IAT clock, which is synchronized to UT2 (by introducing occasional leap ticks) – difference between UTC and IAT is < 0.5 s |

\[ \text{'1 sec.'} \]

1/86400 of a mean solar day … or … 1/31566925.9747 of the tropical year for 1900 (Ephemeris Time defined 1955) … or … 9192631770 periods of the radiation corresponding to the transition between two hyperfine levels of the ground state of the caesium 133 atom

Generating a time frame?

• by a timer generating a regular interrupt
• by employing a RTC-module

Using an existing time-frame?

• by employing time-stamps or sequence numbers of received sensor-readings
• by a radio receiver for UTC or IAT (available in some countries)

Common programming languages guarantee a 'resolution' and 'accuracy' of time (in reference to 'a second'), not its origin or meaning.

How real is real-time?

• 'Real'-time \(\approx\) usually: time as given by external sources.

Engineering: it is of no importance how time is defined and understood, as long as an 'external reference' is given and used as 'real'-time.
Real-Time & Embedded Systems

Notions of time and space

What is embodiment?

Working hypothesis:

☞ Embodied phenomena are those that by their very nature occur in real time and real space.

Phenomenology:

☞ The phenomena of experience as the central aspects and building blocks of understanding.

(applied to and trying to combine aspects of)

☞ Ontology (about the nature of being and categories of existence)

☞ Epistemology (the study of knowledge)

Notions of time and space – Phenomenology

What is embodiment?

Edmund Husserl (1859-1938, Vienna, Halle, Göttingen, Freiburg):

☞ Founder of the phenomenological tradition

☞ as a trial to establish modern science which is firmly grounded on the phenomena of experience (instead of being an abstract mathematical construct).

☞ Phenomenology originally as a method to examine the nature of intentionality

☞ Coined the terms

☞ Noema: the objects of consciousness

☞ Noesis: the mental experiences of those objects

☞ Lebenswelt (life-world): the inter-subjective world of everyday-experience

☞ Husserl rejected pure abstract and formalized reasoning

Martin Heidegger (1889-1976, Freiburg):

☞ Moved phenomenology from a discussion about mental phenomena separated from the physical world (Cartesian dualism)

☞ to a discussion about connected physical and mental phenomena.

☞ Moved the central questions from epistemology to ontology (‘Being and Time’, 1927)

The meaning is not ‘in the head’ but in the world!

☞ Coined terms:

☞ Dasein (being-in-the-world): ‘being as inseparable from the world in which it occurs’

☞ being as always purposeful and active

☞ the world as an unconscious but accessible background

☞ Zuhanden (ready-to-hand): ‘equipment as a part of actual interaction with the world’

☞ Vorhanden (present-at-hand): ‘equipment as a conscious model’
What is embodiment?

Maurice Merleau-Ponty (1908-1961, Paris (Sorbonne)):

- ‘The Phenomenology of Perception’ (1945)
- Embodiment has three implications:
  - a body as a physical entity
  - a body as a set of physical skills and situated responses gained from the physical world
  - a body as a set of ‘cultural skills’ gained from the cultural world in which it is embedded
- Embodied perception as a bi-directional sensation and a basis for empathy
  (Perception in itself does not exist).

☞ see also: Phenomenology of Jean-Paul Sartre

Recent works in robotics (and insights about biological sensors) blur the line between action and perception even further.

Implications:

☞ There is no such thing as

  ‘intelligence’, ‘autonomy’ or any other cognitive process,
  which is independent of a physical environment.

☞ There is no such thing as a

  universal system or body (mechanical design, robot, device, …)
  which is operational in all physical environments.

Meaningfully embedded systems are part of an ‘ecological niche’ (Rolf Pfeifer):

- The operational environment is supportive and employed by the system
- The embedded system is constructed as a part of the operational environment
  and according to the task
- The task is meaningful considering the morphology and cognitive ability of the system
  as well as the response from the environment

☞ i.e. being situated, embodied, and self-sufficient
Real-Time & Embedded Systems

Notions of time and space

What is embodiment?

- Embodied skills depend on a tight coupling between perception and action… up to the level where the distinction between both can become difficult.
- Tight coupling between perception and action means to operate under real-time constraints… and to construct meaningful morphologies.

Real-Time & Embedded Systems

Notions of time and space

The big topics:

What is time? / What is embodiment?

Interfacing with time

Specifying timing requirements

Satisfying timing requirements

Real-Time & Embedded Systems

Interfacing with time

What time is it in …

<table>
<thead>
<tr>
<th>Resolution (syntactical)</th>
<th>Range (all time variables)</th>
<th>Requested resolution (clock ticks)</th>
<th>Actual resolution (detectable?)</th>
</tr>
</thead>
<tbody>
<tr>
<td>JAVA</td>
<td>ms</td>
<td>undefined</td>
<td>undefined</td>
</tr>
<tr>
<td>RT JAVA</td>
<td>ms, ns</td>
<td>undefined</td>
<td>undefined</td>
</tr>
<tr>
<td>Ada95</td>
<td>ms, µs, ns</td>
<td>&gt;50 years</td>
<td>&lt;1ms</td>
</tr>
<tr>
<td>POSIX threads</td>
<td>ms, ns</td>
<td>undefined</td>
<td>&lt;20ms</td>
</tr>
<tr>
<td>C</td>
<td>integer (as seconds)</td>
<td>undefined</td>
<td>undefined</td>
</tr>
</tbody>
</table>

Real-Time & Embedded Systems

Notions of time in Ada

Ada.Real-Time package

```
package Ada.Real_Time is
  type Time is private;
  Time_First : constant Time;
  Time_Last  : constant Time;
  Time_Unit  : constant := 10**1.0#E-9; -- ns
  type Time_Span is private;
  Time_Span_First : constant Time_Span;
  Time_Span_Last  : constant Time_Span;
  Time_Span_Zero  : constant Time_Span;
  Time_Span_Unit  : constant Time_Span;
  Tick : constant Time_Span; -- actual clock resolution < 1 ms
  function Clock return Time;
  -- operations on and conversions with time and time_span ...
end Ada.Real_time;
```
Notions of time in RT-Java

RT-Java time classes

Time root class:

```java
public abstract class HighResolutionTime implements java.lang.Comparable
```

direct known subclasses: `AbsoluteTime, RelativeTime, RationalTime`

- similar to Ada.Real-Time, but no requested accuracy
- adds the concept of frequency (‘rational time’),
  but does not guarantee for equidistant instantiations.

Clock Class:

```java
public abstract class Clock {
    public static Clock getRealtimeClock ();
    public abstract RelativeTime getResolution ();
    public abstract void setResolution (RelativeTime resolution);
}
```

Notions of time in POSIX

Real-time clock interface in POSIX

```c
#define CLOCK_REALTIME ...
struct timespec {
    time_t tv_sec; /* number of seconds */
    long tv_nsec; /* number of nanoseconds */
};
typedef ...
clockid_t;
int clock_gettime (clockid_t clock_id, struct timespec *tp);
int clock_settime (clockid_t clock_id, const struct timespec *tp);
int clock_getres (clockid_t clock_id, struct timespec *res);
int clock_gettime (clockid_t clock_id, struct timespec *tp);
int clock_settime (clockid_t clock_id, const struct timespec *tp);
int clock_getres (clockid_t clock_id, struct timespec *res);
int clock_gettime (clockid_t clock_id, struct timespec *tp);
int clock_settime (clockid_t clock_id, const struct timespec *tp);
int clock_getres (clockid_t clock_id, struct timespec *res);
int clock_gettime (clockid_t clock_id, struct timespec *tp);
int clock_settime (clockid_t clock_id, const struct timespec *tp);
int clock_getres (clockid_t clock_id, struct timespec *res);
int clock_gettime (clockid_t clock_id, struct timespec *tp);
```

Interfacing with time

Programming primitive ‘Delay’

- Alternative for ‘busy-wait’ and interrupts: Suspend a task for a fixed time
- but all these process or thread delays are precise only in their lower bound!
Interfacing with time

Programming primitive ‘Delay’

- Alternative for ‘busy-wait’ and interrupts: Suspend a task for a fixed time
  - **local drift:** sum of all additional actual delays

```
intended delay
```

```
actual normal delay
```

```
local drift
```

### Task 

```
task T;
task body T is
begin
  loop
    Action;
    delay 5.0; -- sec.
  end loop;
end T;
```

- This loop will delay at least for 5 seconds:
  - local and cumulative drift effects here!

Relative delay

(Ada95)

### Task 

```
task body T is
begin
  Interval: constant Duration:= 5.0; -- sec.
  Start_Time: Time;
  loop
    Start_Time:= Clock;
    Action;
    delay Interval - (Clock - Start_Time);
  end loop;
end T;
```

- Delay time calculation is not atomic!
Relative as absolute delay?

(task body T)

Interval: constant Duration := 5.0; -- sec.
Start_Time: Time;
begin
  loop
    Start_Time := Clock;
    Action;
    delay Interval - (Clock - Start_Time);
  end loop;
end T;

Absolute delay

(task body T)

Interval: constant Duration := 5.0; -- sec.
Next_Time: Time;
begin
  Next_Time := Clock + Interval;
  loop
    Action;
    delay until Next_Time;
    Next_Time := Next_Time + Interval;
  end loop;
end T;

Zero delay?

(task body T)

begin
  loop
    Action;
    delay 0.0;
  end loop;
end T;

Delays

Absolute & relative delays are available in:
- Real-Time Java, Pearl, Ada95, ... and many other process-oriented languages

Only relative delays are available in some ‘low-level’ environments:
- POSIX: `nanosleep` (absolute delays need to be constructed employing timers and signals)

Only absolute delays are available in some ‘harder’ real-time environments:
- Occam2, Ada95 (Ravenscar profile), ...
As a third alternative to busy-waiting and infinite blocking, timeouts are implemented in:

- **Shared variable** communications
  - Semaphore
  - Conditional critical regions
  - Monitors
  - Protected objects
- **Message passing** between processes
  - Asynchronous and synchronous message transfers
  - Remote procedure calls
  - Remote objects
- **Actions**

POSIX:

```c
if (sem_timedwait (&call, &timeout) < 0) {
    if (errno == ETIMEDOUT) {      /* timeout occurred, try something else */   }   else {       /* some other error occurred, do something about it */   } }else {   /* semaphore is locked successfully, go ahead */
```

Suspend current process until the semaphore `call` is open or the time-span `timeout` has passed.

Timeouts on entry calls

```c
(task body Sensor is
  T : Temperature;
  begin
    loop
      select
        Controller.Call(T);
      or
        delay 0.5; -- sec.
        -- action if temperature could not be delivered in time
      end select;
      end loop;
    end Sensor;
```

Try calling for 500 ms

Timeouts on incoming calls

```c
(task body Controller is
  Current_Temp : Temperature;
  begin
    loop
      select
        accept Call (T: Temperature) do
          Current_Temp:= T;
        end Call;
      or
        delay 1.0; -- sec.
        -- action if the temperature was not available in time
      end select;
      end loop;
    end Controller;
```

Accept calls for a limited time-span of 1s
### Interfacing with time

#### Timeouts on incoming calls

```vhdl
task body Controller is
    -- declarations
    begin
        loop
            select
                accept Call (T: Temperature) do
                    Current_Temp:= T;
                    end Call;
            or
                delay until Deadline;
                -- no further calls before the deadline
            end select;
            -- normal processing
        end loop;
    end Controller;
```

- **Timeouts on incoming calls**
  - Accept any number of calls until a closing time 'Deadline' for this entry

#### No-wait on incoming calls

```vhdl
task body Controller is
    -- declarations
    begin
        loop
            select
                accept Shutdown do
                    -- termination actions
                    exit;
                end Shutdown;
            else
                -- normal operation
            end select;
            -- synchronize
        end loop;
    end Controller;
```

- **No-wait on incoming calls**
  - Synchronous alternative for an interrupt acceptance

#### No-wait on entry calls

```vhdl
(task same for task-entry calls (message passing) and protected object calls (monitors))
```

```vhdl
task body Sensor is
    T : Temperature;
    begin
        loop
            select
                Controller.Call(T);
            else
                -- action if temperature can not yet be delivered,
                -- e.g. further refine the measurement
            end select;
        end loop;
    end Sensor;
```

- **No-wait on entry calls**
  - Try delivering, else refine the result further

#### Timeout on actions

- All above timeouts suspend / activate a process / task / thread at a synchronization point.
- Up to now, there wasn't any abortion of on-going actions, due to a timeout

To achieve this there need to be an

- **asynchronous change of control flows**
  - on the level of code-blocks: Timeout on actions
  - on the level of processes / tasks / threads: (investigated later)
Monitor a code block:

```ada
select
    delay until Deadline;
    -- computations did not finish in time: take measures
then abort
    -- hard to predict sequence of computations
end select;
```

**Common RT-systems concept:**

1. **Get a first approximation** in fixed amount of time and well before the deadline
2. Inspect the deadline and if there is enough spare time:
3. **Improve the result** and keep a record of improvements (while keeping an eye on the deadline)

   3-a If the most precise result is delivered before the deadline occurs: ✔
   3-b else use the closest approximation so far: ✔

**The deadline is fulfilled and there is a result in any case**

Get a first approximation and employ spare time for refinements:

```ada
Deadline := _ -- set an absolute deadline for the computations
-- compulsory computations (save first result)
select
    delay until Deadline;
    Precise_Result := False;
then abort
    -- optimising computations (save results after each iteration)
while Result_Can_Be_Improved loop
    improve, improve, improve, ...
end loop;
Precise_Result := True;
end select;
-- use result
```

Time-base can also be given externally, e.g., via a protected call:

```ada
loop
    select
        Get_New_Data (Current_Sensor_Data);
        -- employ precise results based on previous data
        -- compulsory computations (save first result)
        -- employ first result on current data
    then abort
        -- optimising computations (save results after each iteration)
        improve, – improve, – improve, ...
    while Result_Can_Be_Improved loop
        improve, – improve, – improve, ...
end loop;
end select;
```

**Timeliness is often more important than Precision**

1. Get a first approximation in fixed amount of time and well before the deadline
2. Inspect the deadline and if there is enough spare time:
3. Improve the result and keep a record of improvements (while keeping an eye on the deadline)

   3-a If the most precise result is delivered before the deadline occurs: ✔
   3-b else use the closest approximation so far: ✔

**The deadline is fulfilled and there is a result in any case**

Take a first guess

Continue in time and with the best result possible
Interfacing with time

Timeout on actions

(RT-Java)

Similar methods in RT-Java:

```java
public class Timed extends AsynchronouslyInterruptedException
    implements java.io.Serializable
{
    public Timed (HighResolutionTime time) throws IllegalArgumentException;
    public boolean doInterruptible (Interruptible logic);
    public void resetTime (HighResolutionTime time);
}
```

☞ see full example of imprecise computations in RT-Java in the course textbook.

(timeouts on actions in POSIX need to be emulated employing timers and signals)

Notions of time and space

What is time? / What is embodiment?

Interfacing with time

Specifying timing requirements

Satisfying timing requirements

Temporal scopes

Common attributes:

- Minimal & maximal delay after creation
- Maximal elapsed time
- Maximal execution time
- Absolute deadline

Common attributes:

- Minimal & maximal delay after creation
- Maximal elapsed time
- Maximal execution time
- Absolute deadline
Real-Time & Embedded Systems

Specifying timing requirements

Temporal scopes

Common attributes:
- Minimal & maximal delay after creation
- Maximal elapsed time
- Maximal execution time
- Absolute deadline

Some common scope attributes

Temporal Scopes can be:
- Periodic – e.g. controllers, samplers, monitors
- Aperiodic – e.g. ‘periodic on average’ tasks, burst requests
- Sporadic / Transient – e.g. mode changes, occasional services

Deadlines (absolute, elapse, or execution time) can be:
- Hard – single failure leads to severe malfunction
- Firm – results are meaningless after the deadline
- Soft – only multiple or permanent failures threaten the whole system

Distinction is not so obvious in practical systems

Hard
- single failure leads to severe malfunction
- deadline

Firm
- only multiple or permanent failures threaten the whole system

Soft
- results may still be useful after the deadline

Real-Time & Embedded Systems

Specifying timing requirements

Temporal scopes

Common attributes:
- Minimal & maximal delay after creation
- Maximal elapsed time
- Maximal execution time
- Absolute deadline

Some common scope attributes

Temporal Scopes can be:
- Periodic – e.g. controllers, samplers, monitors
- Aperiodic – e.g. ‘periodic on average’ tasks, burst requests
- Sporadic / Transient – e.g. mode changes, occasional services

Deadlines (absolute, elapse, or execution time) can be:
- Hard – single failure leads to severe malfunction
- Firm – results are meaningless after the deadline
- Soft – only multiple or permanent failures threaten the whole system

Distinction is not so obvious in practical systems
Specifying timing requirements

Language support for specifying temporal scopes

Real-time Euclid

Real-time Euclid

Language features:

- Recursions and goto-statements are prohibited.
- Loops are restricted to time bounded loops.
- Processes are static and non-nested.

Time scopes:

- periodic <Frameinfo> first activation <TimeOrEvent>
- atEvent <ConditionalId> <Frameinfo>

Periodic process example:

```
realTimeUnit := 1.0 % = 1 seconds

var Reactor : module
var startMonitoring : activation condition atLocation 16#A10D

process TempController : periodic
frame 60
first activation atTime 600 or atEvent startMonitoring

% import list
% execution part
end TempController

end Reactor
```

Ada-emulation of the above RT-Euclid program

```
task body Temp_Controller is

Next_Release : Duration;

begin
select accept Start_Monitoring;
or delay 600.0; -- sec.
end select;

Next_Release := Clock + 60.0; -- sec.
loop
-- execution part
  delay until Next_Release;
  Next_Release := Next_Release + 60.0; -- sec.
end loop;

end Temp_Controller;
```
Real-time Euclid

task body Temp_Controller is
    Next_Release : Duration;
begin
    select accept Start_Monitoring;
    or delay 600.0; -- sec.
    end select;
    Next_Release := Clock + 60.0; -- sec.
    loop
        -- execution part
        delay until Next_Release;
        Next_Release := Next_Release + 60.0; -- sec.
    end loop;
end Temp_Controller;

Ada-emulation of the above RT-Euclid program

State:

Real-Time Euclid was suggested by Kligerman and Stoyenko in 1986
☞ Additional schedulability analysis modules became available
☞ Stayed in an academic context, but influenced many more recent RT-systems.

CRL

(a language for complex real-time systems)

Ways to handle ‘dangerous constructs’ (loops, recursions, synchronisations, …):

- Exclude them
  ☞ Real-time Euclid
- Expand them to become individually safe
  (e.g. by adding mandatory timeout)
- Attribute the code with additional constraints,
  enabling a full pre-runtime analysis ☞ CRL

Constraints in CRL on:

- **Time:**
  - `timeconstraint` (also relative constraints)
  - `nosoonerthan` | `nolaterthan <abs_time>`
  - `endtimeconstraint`

- **Iterations:** assert lower and upper limits for the number of iterations per block

- **Activations:**
  - `activationdeactivationconstraint`
  - `periodic <Frameinfo> firstactive <TimeOrEvent> | atEvent <ConditionalId> <Frameinfo>`
  - `endactivationdeactivationconstraint`

- **Direct recursions:** assert lower and upper recursion limits (general recursions are not covered).
Specifying timing requirements

CRL
(a language for complex real-time systems)

Evaluation of the constraints and assertions:

- **Timing:**
  - verified at compile-time
- **Activation / Deactivation:**
  - checked for schedulability at compile-time and enforced by the scheduler at run-time.
- **Iterations and recursion:**
  - either verified at compile-time or checked at run-time.

State:

CRL was suggested by Stoyenko, Marlowe and Younis in 1995

- Full featured language
- Compiled to attributed C++

☞ Stayed also in an academic context.

Pearl

Explicit time-scope expressions:

- **TaskStart** ::= [StartCondition] ACTIVATE <task>
- **StartCondition** ::= AT <time> [Frequency] | AFTER <duration> [Frequency] | WHEN <interrupt> [AFTER <duration> [Frequency]] | Frequency

- **Frequency** ::= ALL <duration> | {[UNTIL <time> ] | [DURING <duration>]}]

☞ Schedulability analysis (at compile-time or run-time) possible
(although not defined by the language)
☞ Pearl refinement: a combination of Pearl and Real-time Euclid = High-Integrity Pearl

DSP: Distributed Programming System

Explicit time-scope expressions at the ‘statement level’:

- e.g. time-scope for a software-engineer:
  - from 11:00 to 19:30 every 45
  - start elapse 10 do
    - setup_coffee_machine
    - power_coffee_machine up
    - find_favrite_cup
    - put_coffee_in_favrite_cup
    - clean_coffee_machine
  - end
  - start after 3 elapse 25 by 20:00 do
    - drink_coffee
  - end

☞ DSP-compiler: breaks the sequences down in processes and schedules
Real-time Java comes with:

- multiple sets of predefined time-scope parameters
- a scheduler class (with a predefined priority scheduler)
- Schedulability (feasibility) analysis possible.

```
public abstract class Scheduler {
    protected Scheduler ();
    protected abstract boolean addToFeasibility (Schedulable s);
    public abstract void fireSchedulable (Schedulable s);
    public abstract boolean isFeasible ();
    protected abstract boolean removeFromFeasibility (Schedulable s);
    public boolean setIfFeasible (Schedulable s, ReleaseParameters r, MemoryParameters m);
}
```

Formulates an on-line schedulability analysis!

```
public class PriorityScheduler extends Scheduler {
    public static final int MAX_PRIORITY;
    public static final int MIN_PRIORITY;
    protected PriorityScheduler ();
    protected boolean addToFeasibility (Schedulable s);
    public void fireSchedulable (Schedulable s);
    public boolean isFeasible ();
    protected boolean removeFromFeasibility (Schedulable s);
    public boolean setIfFeasible (Schedulable s, ReleaseParameters r, MemoryParameters m);
}
```

This PriorityScheduler is the only requested instantiation.

```
public class RealtimeThread extends java.lang.Thread implements Schedulable {
    public RealtimeThread (SchedulingParameters s, ReleaseParameters r, MemoryParameters m, MemoryArea a);
    public synchronized void addToFeasibility ();
    public synchronized void addIfFeasible ();
    public static RealtimeThread currentRealtimeThread () throws ...;
    public synchronized void deschedulePeriodic ();
    public synchronized void interrupt ();
    public synchronized void sleep (...) throws ...;
}
```

Priority, Periodic, aperiodic, or sporadic parameters
Real-time Java

```java
public class NoHeapRealtimeThread extends RealtimeThread {
    public RealtimeThread (SchedulingParameters s,
                           ReleaseParameters r,
                           MemoryArea a) throws ...;
}
```

This thread is allowed to interrupt any garbage collector at any time!
(since it doesn't depend on it itself)

Real-time Java

```java
public abstract class SchedulingParameters {
    public SchedulingParameters ();
}
```

```java
public class PriorityParameters extends SchedulingParameters {
    public PriorityParameters (int priority);
    public int getPriority ();
    public void setPriority (int priority) throws ...;
}
```

'Priority' is the only default scheduling parameter

Real-time Java

```java
public abstract class ReleaseParameters {
    protected ReleaseParameters (RelativeTime cost,
                                  RelativeTime deadline,
                                  AsyncEventHandler overrunHandler,
                                  AsyncEventHandler missHandler);
    public RelativeTime getCost();
    public AsyncEventHandler getCostOverrunHandler();
    public RelativeTime getDeadline();
    public AsyncEventHandler getDeadlineMissHandler();
}
```

Cost is an estimate of the max. execution time

Measuring execution time is not requested, i.e. the overrunHandler might never be activated!

Real-time Java

```java
public class PeriodicParameters extends ReleaseParameters {
    public PeriodicParameters (HighResolutionTime start,
                                RelativeTime period,
                                RelativeTime cost,
                                RelativeTime deadline,
                                AsyncEventHandler overrunHandler,
                                AsyncEventHandler missHandler);
    public RelativeTime getPeriod ();
    public HighResolutionTime getStart ();
    public void setPeriod (RelativeTime period);
    public void setStart (HighResolutionTime start);
}
```

most frequently used release parameters
Specifying timing requirements

Real-time Java

```java
public class AperiodicParameters extends ReleaseParameters {
    public AperiodicParameters (RelativeTime cost, RelativeTime deadline,
        AsyncEventHandler overrunHandler, AsyncEventHandler missHandler);
}
```

these are the minimum release parameters (while cost might be used for feasibility analysis only)

```java
public class SporadicParameters extends AperiodicParameters {
    public SporadicParameters (RelativeTime minInterarrival, RelativeTime cost,
        RelativeTime deadline, AsyncEventHandler overrunHandler,
        AsyncEventHandler missHandler);
    public RelativeTime getMinimumInterarrival ();
    public void setMinimumInterarrival (RelativeTime minimum);
}
```

Sporadic events are not allowed to come in bursts!

The minimal required implementation supplies:

☞ Priority scheduling
☞ On-line schedulability analysis
☞ Deadline violation handlers
☞ (max. execution time deadline checks are suggested but not required)
☞ a sporadic scheduler is not required (although the sporadic release parameter set is).

Hard real-time environments require the exclusive usage of ‘No heap real-time threads’ (synchronization with an 'object-oriented' thread invalidates the feasibility assurances)

Real-Time & Embedded Systems

Specifying timing requirements

Real-time Java

The minimal required implementation supplies:

☞ Priority scheduling
☞ On-line schedulability analysis
☞ Deadline violation handlers
☞ (max. execution time deadline checks are suggested but not required)
☞ a sporadic scheduler is not required (although the sporadic release parameter set is).

Hard real-time environments require the exclusive usage of ‘No heap real-time threads’ (synchronization with an 'object-oriented' thread invalidates the feasibility assurances)

Ada95

Ada95 has no explicit time-scope expressions at task-level.

Ada95 offers …

• tasking
• a priority scheduler (the only required scheduler)
• synchronisation and communication primitives
• asynchronous transfer of control, timed calls, timeout on actions … … … etc. pp.
• … but no hardware timers!

…but no automatic schedulability analysis!

… to create the basis for most kinds of hard real-time scopes manually.
Since Esterel is a synchronous language, …
… all actions and communications take zero time by definition.
☞ There is no expression for continuous, non-zero time-scopes.
☞ Time is interpreted in the reductionistical way as a sequence of events
☞ Time-scopes translate to signal-relations and signal-counters
Continuous time scopes need to be taken into account while
1. analysing and reducing the problem to a zero-time atomic system
2. implementing the synchronous system on an actual system.
☞ Continuous time-scopes for the …
... validation of the zero-time assumption!

POSIX
the usual: use timers!
common combination:
☞ usage of Ada95 together with POSIX timers as a basis for hard-real-time-scopes and schedulers

The big topics:

What is time? / What is embodiment?
Interfacing with time
Specifying timing requirements
Satisfying timing requirements

Two paths towards fulfilling rt-requirements:
☞ Real-time logic approach
  formal, correct in its specifications, & offers calculus for asynchronous, real-time systems
  needs to ignore most real world effects, like jitters, drifts, failures, interferences, etc. pp.
  gives a correct solution according to the specification
☞ Complex systems oriented approach
  deals with existing computer systems, sensors, & offers a set of approximating methods
  not complete or correct in any formal sense
  deals with real-world systems, gives ‘robust’ systems, passes rigorous experiments
Satisfying timing requirements

Fulfilling rt-requirements:

Complex systems oriented approach:
- System identification and compile-time analysis:
  - Calculate or limit statement durations
  - Calculate or limit iterations and recursions
  - Analyse potential dead- or life-locks ⇔ chapter 8
  - Calculate schedulability ⇔ chapter 7
- Run-time analysis and checks:
  - Dynamic scheduling schemes: Re-validate schedulability ⇔ chapter 7
  - Check for all constraints and assertions at run-time ⇔ chapter 9
- Supply fault-tolerant behaviours:
  - Error recoveries, mode changes, … ⇔ chapter 9

Real-time logic approach:
- Reduce the problem:
  - Reduce any asynchronous, analogue, dynamical, fractal, jitter-, drift, or failure-affected parts of the system to a fully synchronous and discrete system ⇔ chapters 5 and 6
  - Formulate the specification on the basis of the reduced synchronous system.
- Verify the reduced system:
  - Verify the reduced synchronous against the specifications ⇔ not covered in this course
- Compile the reduced system to an actual system:
  - The resulting actual system will be executable on a real machine and employ real devices
  ⇔ Re-check the actual system (e.g. by means of a complex systems-approach) ⇔ ...

Summary

Time & Space

- What is time? / What is embodiment?
  - Approaches by different faculties to understand the basis for this course
- Interfacing with time
  - Formulating local time-dependent constraints
  - Access time, delay processes, detect timeouts (in different languages)
- Specifying timing requirements
  - Formulating global timing-constraints
  - Understanding time-scope parameters (and expressing them in different languages)
- Satisfying timing requirements
  - Real-time logic and complex systems approach
Asynchronism

Interrupts

Required mechanisms for interrupt driven programming:

- **Interrupt control**: grouping, encoding, prioritising, and en-/disabling interrupt sources
- **Context switching**: mechanisms for cpu-state saving and restoring + task-switching
- **Interrupt identification**: interrupt vectors, interrupt states

```
hardware-supported
```

## Asynchronism

### Interrupts

**Interrupt control**:

- ... at the individual device level
- ... at the system interrupt controller level
- ... at the operating system level

---

### References for this chapter

[Bollella01] Greg Bollella, Ben Brosbol, Steve Furr, David Hardin, Peter Dibble, James Gosling, Mark Turnbull & Rudy Belliardi

Real-Time Systems and Programming Languages

Addison Wesley, third edition, 2001

[Ada95RM] (link to on-line version)

Ada Working Group

ISO/IEC JTC1/SC22/WG 9

Ada 95 Reference Manual

– Language and Standard Libraries


[Burns01] Alan Burns and Andy Wellings

Real-Time Systems and Programming Languages

Addison Wesley, third edition, 2001

all references and links are available on the course page

---

### LM12L458 – accessible registers

<table>
<thead>
<tr>
<th>Address</th>
<th>Purpose</th>
<th>Type</th>
<th>D15</th>
<th>D14</th>
<th>D13</th>
<th>D12</th>
<th>D11</th>
<th>D10</th>
<th>D9</th>
<th>D8</th>
<th>D7</th>
<th>D6</th>
<th>D5</th>
<th>D4</th>
<th>D3</th>
<th>D2</th>
<th>D1</th>
<th>D0</th>
</tr>
</thead>
<tbody>
<tr>
<td>0 0 0 0</td>
<td>Instruction RAM (RAM Pointer = 00)</td>
<td>R/W</td>
<td>Acquisition Time</td>
<td>Write-</td>
<td>0/T</td>
<td>Timer Sync</td>
<td>VDD</td>
<td>VSS</td>
<td>Pause</td>
<td>Loop</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>0 0 0 0</td>
<td>Instruction RAM (RAM Pointer = 01)</td>
<td>R/W</td>
<td>Don’t Care</td>
<td>&gt;</td>
<td>Sign</td>
<td>Limit #1</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>0 0 0 0</td>
<td>Instruction RAM (RAM Pointer = 10)</td>
<td>R/W</td>
<td>Don’t Care</td>
<td>&gt;</td>
<td>Sign</td>
<td>Limit #2</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>1 0 0 0</td>
<td>Configuration Register</td>
<td>R/W</td>
<td>Don’t Care</td>
<td>D/A</td>
<td>Test</td>
<td>0</td>
<td>ROM Pointer</td>
<td>IO</td>
<td>Auto</td>
<td>Zero</td>
<td>Chan</td>
<td>Mask</td>
<td>Stand</td>
<td>Cal</td>
<td>Run</td>
<td>Top</td>
<td>Reset</td>
<td>Start</td>
</tr>
<tr>
<td>1 0 0 1</td>
<td>Interrupt Enable Register</td>
<td>R/W</td>
<td>Number of Conversions in Conversion FIFO to Generate INT2</td>
<td>Sequence Address to Generate INT1</td>
<td>INT7</td>
<td>Don’t Care</td>
<td>INT5</td>
<td>INT4</td>
<td>INT3</td>
<td>INT2</td>
<td>INT1</td>
<td>INT0</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>1 0 1 0</td>
<td>Interrupt Status Register</td>
<td>R</td>
<td>Actual Number of Conversion Results</td>
<td>Conversion FIFO</td>
<td>Address of Sequencer</td>
<td>Instruction being Executed</td>
<td>INST7</td>
<td>INST5</td>
<td>INST4</td>
<td>INST3</td>
<td>INST2</td>
<td>INST1</td>
<td>INST0</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>1 0 1 1</td>
<td>Timer Register</td>
<td>R/W</td>
<td>Timer Preset High Byte</td>
<td>Timer Preset Low Byte</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>1 1 0 0</td>
<td>Conversion FIFO</td>
<td>R</td>
<td>Address or Sign</td>
<td>Conversion Data: MSBs</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>1 1 0 1</td>
<td>Limit Status Register</td>
<td>R</td>
<td>Link #2: Status</td>
<td>Limit #1: Status</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>
**Interrupts**

**LM12L458**  
(National Semiconductor)

- **Select interrupt sources** (interrupt enable register, 15 bits):
  - **Watchdog**: limit conditions are fulfilled
  - **Instruction**: the instruction pointer equals a pre-programmed value (bits 8-10)
  - **Conversions**: a specified number of conversions (bits 11-15) have been performed
  - **Auto-Zero**: short calibration has been performed
  - **Full-Calibration**: long calibration has been performed
  - **Pause**: Sequencer arrived at a pause state
  - **Active**: Controller returned from power-down to active mode

- **Read interrupt status** (interrupt status register, 15 bits):
  - indicates the current interrupt conditions, incl. the actual number of conversions and the currently processed instruction

☞ all of the status bits (0-5, 7) are reset with any read access to this register!

- **Interrupt control**:
  - ... at the individual device level
  - ... at the system interrupt controller level
  - ... at the operating system level

☞ only one interrupt signal line available!

☞ in order to identify the interrupt reason, an additional read cycle is required!
System interrupt controller:

- collects the interrupt signal lines from all managed devices
- identifies the device status, if not given already as an interrupt vector
- encodes all interrupt signals into a common interrupt vector or status scheme
- orders and masks all interrupts according to priority levels
- triggers the actual CPU or controller interrupt

☞ three common forms:
  - the interrupt controller as an intrinsic part of a complex µcontroller
  - a dedicated interrupt controller for a set of similar or identical devices
    (e.g. a hard disk array)
  - a universal interrupt controller (usually unable to fetch interrupt status information)

Interrupt control:

... at the individual device level

... at the system interrupt controller level

... at the operating system level

- beyond task-level
- communicating interrupts to task
- transforming interrupts to signals

Interrupt service routines (available only in some OSs, e.g. VxWorks)

Purpose:

- Allow full access to the interrupt controller (interrupt vectors, priorities).
- Change to an interrupt service routine in a predictable amount of time.

☞ Cannot operate on the level of threads or tasks!

☞ Limitations regarding the accessibility of some OS-facilities (task level system calls).

- Real-time-operating systems and real-time-languages provide this access.

Some VxWorks OS entries:

<table>
<thead>
<tr>
<th>Call</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>intConnect</td>
<td>Connect a routine to an interrupt vector</td>
</tr>
<tr>
<td>intLevelSet</td>
<td>Set the interrupt mask level</td>
</tr>
<tr>
<td>intLock</td>
<td>Disable interrupts (besides NMI)</td>
</tr>
<tr>
<td>intUnlock</td>
<td>Enable interrupts</td>
</tr>
<tr>
<td>intVecBaseSet</td>
<td>Set the interrupt vector base address</td>
</tr>
<tr>
<td>intVecBaseGet</td>
<td>Get the interrupt vector base address</td>
</tr>
<tr>
<td>intVecGet</td>
<td>Set an interrupt vector</td>
</tr>
<tr>
<td>intVecGet</td>
<td>Get an interrupt vector</td>
</tr>
</tbody>
</table>

these calls are employed by the language run-time environment or used directly from ‘C’-code
Asynchronism

Interrupt service routines
(available only in some OSs, e.g. VxWorks)

Minimal hardware support (supplied by the cpu):
- save essential CPU registers (IP, condition flags)
- jump to the vectorized interrupt service routine

Minimal wrapper (supplied by the real-time-os):
- save remaining CPU registers
- save stack-frame
- execute user level interrupts service code
- restore stack-frame
- restore CPU registers
- restore IP

Interrupt service routine to task communication methods:
- Shard memory and ring buffers:
  most low level communication scheme (should be avoided)
- Semaphore:
  trigger a semaphore, where a task has been blocked before.
- Monitors:
  free a task, which is blocked at a monitor entry (standard Ada-method: protected object).
- Message queues:
  Send messages to a task (if queue is not full).
- Pipes:
  Write to a pipe (if pipe is not full).
- Signals:
  indicate an asynchronous task switch to the scheduler
☞ in all of the above: the interrupt service routines cannot block!

Interrupts ☞ ‘Signals’

Interrupt control:

... at the individual device level
... at the system interrupt controller level
... at the operating system level
- beyond task-level
- communicating interrupts to task
- transforming interrupts to signals

Some characteristics of signals:
- Involve a full task-switch operation
☞ Hard to predict timing behaviour
- Limited information about the interrupt-source
- Traditionally used to ‘kill’ processes
- Concept stems from a time before thread models,
  therefore the signal-to-thread propagation is implementation dependent and sometimes tricky.
Some common UNIX OS entries:

<table>
<thead>
<tr>
<th>POSIX 1003.1b</th>
<th>BSD-UNIX</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>signal (...)</td>
<td>signal (...)</td>
<td>Specify the handler associated with a signal</td>
</tr>
<tr>
<td>sigaction (...)</td>
<td>sigvec (...)</td>
<td>Examine or set the signal handler for a signal</td>
</tr>
<tr>
<td>kill (...)</td>
<td>kill (...)</td>
<td>Send a signal (overwrite all other pending signals)</td>
</tr>
<tr>
<td>sigqueue (...)</td>
<td>N/A</td>
<td>Send a queued signal</td>
</tr>
<tr>
<td>sigsuspend (...)</td>
<td>pause (...)</td>
<td>Wait for a signal</td>
</tr>
<tr>
<td>sigwaitinfo (...)</td>
<td>sigtimedwait (...)</td>
<td>Wait for a signal, but do not involve the handler</td>
</tr>
<tr>
<td>sigemptyset (...)</td>
<td>sigsetmask (...)</td>
<td>Manipulate and set the mask of blocked signals</td>
</tr>
<tr>
<td>sigprocmask (...)</td>
<td>sigblock (...)</td>
<td>Add to a set of blocked signals</td>
</tr>
</tbody>
</table>

• Signals are originally process-level synchronization methods (‘kill’) and have been expanded to be used for everything from hardware-interrupts and timers to asynchronous task messaging.

☞ Signals are passed through a global task-scheduler.

☞ In many OSs: unpredictable ‘work-arounds’ for missing direct hardware interrupt propagation.

☞ Make sure that you understand the attached strings in your OS, before employing any signals.

☞ RT-environments always impose restrictions for the interrupt handler.

The handler then either...

• ... deals with the situation itself (employing its limited capabilities)

• ... initiates a general change in the control flow (involving other parts of the system).

☞ Formulate restrictions with respect to the interruptible code!

e.g. Ada: ‘asynchronous transfer of control’, Real-Time Java: ‘Interrupted exceptions’

(exception handling concepts and atomic actions will be introduced first, then the discussion about asynchronous transfer of control methods is continued)

Wish list for exception handling in real-time programming languages:

1. Exception facilities should not obscure the understanding of the normal, exception-free control flow.

2. Exception facilities should produce no or minimal run-time overhead, until an exception actually occurs.

3. Exceptions should produce predictable run-time overhead when an exception occurs.

4. Exceptions indicated by the run-time environment and by the program itself should be treated uniformly.

5. The exception mechanism should be applicable to asynchronous and synchronous exceptions (might be hard to achieve with respect to wish 3).

6. The exception mechanism should allow for appropriate recoveries (supply sufficient information and appropriate re-entry possibilities).
Asynchronism

Exceptions

Historic exception handling methods:

- Use an ‘unusual return value’ and a global variable convention: ‘C’:

```c
if (function_call (parameters) < NORMAL_RETURN_VALUE) {
    if (errno == SOMEKNOWN_ERROR) {
        /* react to the known error condition */
    } else {
        /* try to improvise something */
    }
} else {
    -- normal control flow
}
```

☞ inflexible (exceptions from the environment cannot be detected)
☞ error-prone (lots of chances to forget checking or to use the wrong constants)
☞ obstructive (all fragments are commingled)

Assembler level exception handling methods:

- Provide a jump table and manipulate the return address on the stack:

Caller:
```
jsr pc, PRINT_CHAR
jmp IO_ERROR
jmp DEVICE_NOT_ENABLED
```

Subroutine:
```
% indicate an exception:
% increment the return address on the stack
% by the exception number
% to employ the caller-provided exception handling
% indicate normal operation:
% increment the return address on the stack
% by the max. exception number + 1
```

Emulating exception handling methods (in older languages):

- Unrecoverable exceptions: ⇣ provide a jump label.
- Recoverable exception: ⇣ provide a procedure variable.

☞ Historic, since all current real-time suitable languages provide some means of dedicated exception handling.
### Exception indication

Four cases of modern exception indication:

<table>
<thead>
<tr>
<th>Raised:</th>
<th>From:</th>
<th>Application</th>
</tr>
</thead>
<tbody>
<tr>
<td>Synchronously</td>
<td>Run-time environment</td>
<td>Application</td>
</tr>
<tr>
<td>Asynchronously</td>
<td>*</td>
<td>*</td>
</tr>
</tbody>
</table>

Ada95:

<table>
<thead>
<tr>
<th>Raised:</th>
<th>From:</th>
<th>Application</th>
</tr>
</thead>
<tbody>
<tr>
<td>Synchronously</td>
<td>Run-time environment</td>
<td>Exceptions</td>
</tr>
<tr>
<td>Asynchronously</td>
<td>Interrupt/signal handler</td>
<td>Asynchronous transfer of control</td>
</tr>
</tbody>
</table>

Real-time Java:

<table>
<thead>
<tr>
<th>Raised:</th>
<th>From:</th>
<th>Application</th>
</tr>
</thead>
<tbody>
<tr>
<td>Synchronously</td>
<td>Exceptions</td>
<td></td>
</tr>
<tr>
<td>Asynchronously</td>
<td>Asynchronous exceptions</td>
<td></td>
</tr>
</tbody>
</table>

POSIX:

<table>
<thead>
<tr>
<th>Raised:</th>
<th>From:</th>
<th>Application</th>
</tr>
</thead>
<tbody>
<tr>
<td>Synchronously</td>
<td>Run-time environment</td>
<td>N/A</td>
</tr>
<tr>
<td>Asynchronously</td>
<td>Signals</td>
<td>(Signals)</td>
</tr>
</tbody>
</table>
Asynchronism

Exception granularity

at block level:

Ada95:

```ada
begin
   do something dangerous
exception
   when E: Constraint_Error => Deal with it;
end;
```

or in RT-Java

```java
try {
   // do something dangerous
}
catch (ExceptionType e) {
   // handle exception e
}
```

all exceptions need to be (or are) declared!

Ada95:

```ada
begin
   do something dangerous
exception
   when E: others => Deal with it;
end;
```

or in RT-Java

```java
try {
   // do something dangerous
}
catch (Exception e) {
   // handle exception e
}
```

handlers can catch all! but don't need to catch any

Asynchronism

Exception granularity

large blocks:

```ada
declare
subtype Temperature is Integer range 0 .. 100;
subtype Pressure is Integer range 0 .. 50;
subtype Flow is Integer range 0 .. 200;
begin
   -- read temperature sensor and calculate its value
   -- read pressure sensor and calculate its value
   -- read flow sensor and calculate its value
   -- adjust temperature, pressure and flow
      -- according to requirements
exception
   -- handler for Constraint_Error
end;
```

might be too unspecific

Asynchronism

Exception granularity

small blocks:

```ada
declare
begin
   -- read temperature sensor and calculate its value
exception -- handler for Constraint_Error for temperature
end;
begin
   -- read pressure sensor and calculate its value
exception -- handler for Constraint_Error for pressure
end;
begin
   -- read temperature sensor and calculate its value
exception -- handler for other possible exceptions
end;
```
Asynchronism

Exception granularity

small blocks:

```
declare
begin
    begin
        -- read temperature sensor and calculate its value
        exception -- handler for Constraint_Error for temperature
    end;
    begin
        -- read pressure sensor
        exception -- handler for Constraint_Error for pressure
    end;
    exception -- handler for other possible exceptions
end;
```

becomes quickly unreadable

statement level:

```
e.g. supported in CHILL

A : temperature;
B, C : integer;
begin
    A := B + C on
        (overflow) : …;
        (rangefail) : …;
        else …;
    end;
```

Exceptions can also be handled at block, procedure, or process level

```
All exceptions in CHILL need to be handled
```

```
Exceptions are not propagated
```

```
Handlers are determined at compile-time
```

Instead of catching exceptions after each statement:

- Exceptions can have parameters with information about its source:
  - Real-time Java: Exception can carry any number of user-defined parameters
  - Ada95: The environment automatically attaches additional information to exceptions, which are indicating the position of the exception-occurrence and other observed conditions (implementation dependent, inflexible, useful for debugging).

1. All procedures and functions declare every potentially raised exceptions (requested in CHILL, requested in Real-time Java for user-defined exceptions):
   - If an appropriate exception handler can not be determined at compile-time:
     Either ...
     - … treat it as a programmer error and stop compilation (CHILL)
     or ...
     - … propagate the exception at run-time outside its static scope (Real-time Java)
   - Exceptions are declared for whole modules — not specifically for methods (Ada95):
     - The handler is determined at run-time in any case (either by propagation or in the static scope).
Asynchronism

Exception propagation
(tasks)

If the handler can neither be found by propagation or in the static context of a task:
☞ Stop the specific task — exceptions are not propagated to the parent-process (Ada95).
or
☞ Propagate the exception to the parent-task (and potentially stall the whole system).

Common ‘safety-line’ for each process:
☞ Use a ‘catch all’-handler (possible in Ada95 and Real-time Java) at the highest level.
   (often last emergency level for this process ⇔ common reaction: PANIC!)

Asynchronism

Exception handling: resumption or termination?

Resumption-model (offered in: Pearl, Mesa):
1. Find an exception handler.
2. Execute the exception handler (and potentially raise another exception ⇔ recursion).
3. After completion of the exception handler (and possibly other handlers):
   return to the invoker and try resume processing as if nothing happened.

Feature:
• In case of an asynchronous exception, there is little impact on the current control-flow.

Problems:
• some errors cannot be ‘repaired’ (especially all timing related errors).
• exceptions can be raised in the middle of evaluations, which will be hard to restore.

‘Block resumption’-model (offered in: Eiffel):
☞ Re-execute the complete code-block after exception handling

Feature:
• Intended to keep the formulated contract for this method.

Problems:
• Local variables must not be re-initialized (otherwise the exception will probably occur again)
• The code needs to be aware of all possible combinations of half-evaluated processing states.
• Trying the same method again (and again) is usually not the suitable way for real-time systems.

Termination-model (only model in Ada95, Real-time Java; offered in: Pearl, Mesa):
• The control is not returned to the point of invocation.
• Instead the block / function / procedure is assumed to terminated in an exceptional state, and the control is returned to the calling or enclosing scope of the activated exception handler.
• If the calling block wish to re-try the same operation, it need to start over at the visible entry-points and with re-initialized local variables.

Feature:
• The method of choice, if exceptions imply that the operation (statement, block, process) was not successful and something else need to be done now ⇔ real-time systems.

Problem:
• There is no way to continue, in case that the exception could be identified as of minor impact.
Exception handling: resumption or termination?

Hybrid-model (offered in: Mesa):

☞ The exception handler can decide at run-time whether to terminate or to resume.

Cleaning up before exception-handling:

Assuming a block is holding a number of resources, and occurring exceptions need to be handled at the caller level:

```plaintext
procedure Allocate (Number : Devices) is
begin
  -- request each device be allocated in turn
  -- noting which requests are granted
  exception
    when others =>
      -- e.g. deallocate those devices allocated
      raise; -- re-raise the exception
  end Allocate;
```

☞ helpful to keep a consistent system-state and to avoid dead-locks (all-or-nothing allocation).

… in Real-time Java: the ‘finally’ clause takes care of block consistent finalization.

Issues when handling exceptions in Ada95:

☞ Exceptions are declared at package level, i.e. it is unclear, which functions may raise which exceptions!
☞ Exceptions may be propagated outside the scope of their declaration, i.e. only ‘when others’ can handle them (might also be further propagated back in scope again).
☞ Parameter passing limited to one string.
☞ Exception in task bodies are never propagated to the parent task, i.e. if there could not be any handler identified in the task, the task will ‘die silently’.
☞ Exception in task declarations are always propagated to the parent task.
☞ Exceptions in task rendezvous, which are not handled in the accept statement, are propagated to both involved tasks.

☞ Traps, which need to be taken care of!
most expensive not caught exception up to now: half a billion dollars (maiden crash of Ariane 5, ‘96)

Exception handling in Real-time Java

```
Errors
  Throwable
    Exceptions
      Linkage
      VM
      RunTime
    User def.
 .checked => need to be declared per method (runtime exceptions can occur undeclared)
  unchecked (errors are unrecoverable)
```
Exception handling in Real-time Java

Exceptions are objects in Real-time Java:
☞ Exceptions have hierarchical relations
☞ Exceptions handlers can catch
  • one individual exception
  • all exceptions out of a finite list of exceptions
  • all exceptions of a certain class
  • all exceptions
☞ The kinds of exceptions which are handled at a certain point can be described precisely, completely, and safely.
(exceptions are not part of the class-hierarchy in Ada95 or Eiffel)

Exception handling in ‘C’ / POSIX

☞ there is no exception handling in ‘C’ / POSIX

possible work arounds by using POSIX long jumps or signals.
☞ see also macro-assembler or ‘old language’ exception handling methods.

Exception handling: compare sheet

<table>
<thead>
<tr>
<th>Language</th>
<th>Terminating?</th>
<th>Handler</th>
<th>Decl. per:</th>
<th>Decl. as:</th>
<th>Scope</th>
</tr>
</thead>
<tbody>
<tr>
<td>Real-time Java</td>
<td>termination</td>
<td>dynamical / propagating</td>
<td>method</td>
<td>classes</td>
<td>‘try-block’</td>
</tr>
<tr>
<td>Ada95</td>
<td>termination</td>
<td>dynamical / propagating</td>
<td>package</td>
<td>static names</td>
<td>any block</td>
</tr>
<tr>
<td>CHILL</td>
<td>termination</td>
<td>static</td>
<td>via handler</td>
<td>static handler</td>
<td>statement / block / task</td>
</tr>
<tr>
<td>Mesa</td>
<td>resumption or termination</td>
<td>dynamical / propagating</td>
<td>procedure</td>
<td>static procedures</td>
<td>block</td>
</tr>
<tr>
<td>Pearl</td>
<td>resumption or termination</td>
<td>static</td>
<td>process</td>
<td>static names</td>
<td>task</td>
</tr>
<tr>
<td>‘C’ / POSIX</td>
<td>class-retry</td>
<td>dynamical / propagating</td>
<td>(contract violation)</td>
<td>-</td>
<td>‘do-rescue-block’</td>
</tr>
</tbody>
</table>

Atomic actions: definitions:

An action is atomic if the processes performing it …
• ... are not aware of the existence of any other active process, and no other active process is aware of the activity of the processes during the time the processes are performing the action.
• ... do not communicate with other processes while the action is being performed.
• ... cannot detect any outside state change and do not reveal their own state changes until the action is complete.
☞ … can be considered to be indivisible and instantaneous.
Atomic actions: implications:

An atomic action …
• … is either performed fully, or not at all.
• … is declared as failed, if any part of the action fails.

Thus all parts of an atomic action need to be prepared:
☞ to be interrupted (due to the failure of one of them)
☞ and to reset to their initial state at any time
   (‘no effect’ is visible to the outside)

Nested atomic actions:

Action actions can be nested, iff …
• … all processes involved in the nested atomic actions
  are a true subset of the processes involved in the enclosing atomic action.
Asynchronism

Atomic actions – Requirements for real-time environments:

- Well-defined boundaries
  - A start, end and a side boundary:
    - Define clear entry and exit points for all processes involved in the atomic action.
    - Processes can enter at different time, but need are released from the action at once.
  - Separate the involved processes from the rest of the system ('side boundary').

- Indivisibility (Isolation)
  - Prohibit or restrict communications to outside processes and resources.
  - Employing results from one atomic action in another one implies strict serialisation.

- Nesting
  - Atomic actions may be nested, if they form a true enclosing relation.

- Concurrency
  - Independent atomic actions may be executed in any order and concurrently.

Failure in one action part:
☞ 'clean up' (restore the initial states)

Failure in one part of a nested atomic action:
Asynchronism

Atomic actions

Failure in one part of a nested atomic action (iterative failure propagation):

Failure in one part of a nested atomic action (immediate failure propagation):

Nested atomic actions

Failure in one part of a nested atomic action:
Real-Time & Embedded Systems

Asynchronism

Atomic actions

Failure in an enclosing atomic action:

Failure in an enclosing atomic action (no communication with inner action):

Real-Time & Embedded Systems

Asynchronism

Atomic actions

Failure in an enclosing atomic action (no communication with inner action):
Asynchronism

Atomic actions

Failure in an enclosing atomic action:

1. Atomic actions

Failure in an enclosing atomic action (revoking activation condition for inner action):

2. Atomic actions

No mainstream language is supplying such a construct.

Implementing atomic actions by creating dedicated tasks:

3. Atomic actions
Asynchronism

Atomic actions in Ada95:

```
with Atomic_Action_Types; use Atomic_Action_Types;
generic
   Actions : in Action_Parts;
package Generic_Atomic_Action is
   procedure Perform;
   Failure_State, Time_Out_State, Late_Activation_State : exception;
end Generic_Atomic_Action;
☞ Scope mechanism is employed to limit communication possibilities of atomic action parts.
☞ The perform-call is atomic for the caller and invisible for others.
☞ Failures in one part are automatically propagated in the whole atomic action.
```

```
with:
   type Action_Part_Time_Scope is record
      Start_Delay_Min : Time_Span := Time_Span_Zero;
      Start_Delay_Max : Time_Span := Time_Span_Last;
      Max_Elapse : Time_Span := Time_Span_Last;
      Deadline : Time := Time_Last;
      end record;
   type Action_Part_Proc is access procedure;
   type Action_Part_Procs is record
      Action, Cleanup : Action_Part_Proc;
      Scope : Action_Part_Time_Scope;
      end record;
   type Action_Parts is array (Positive range <>) of Action_Part_Procs;
☞ The Cleanup procedure is meant to restore the initial state! (undoing all effects of Action)
```

```
package Actions is new Generic_Atomic_Action (Actions);
procedure Perform is
   begin
      Atomic_Action.Perform;
   exception
      when _
      end Perform;
```

```
task body Action_Task is
   then abort
      begin
         Monitor.Check_In (Task_Id);
         select
            Monitor.Failed (Task_Id);
            Actions (Task_Id).Cleanup.all;
   begin
         then abort
            begin
               Monitor.Check_Out (Task_Id);
               delay To_Duration (Actions (Task_Id).Scope.Start_Delay_Max);
               Deadline := Time_Last;
```

```
begin
   Monitor.Check_In (Task_Id);
   select
      Monitor.Failed (Task_Id);
      Actions (Task_Id).Cleanup.all;
      Atomic_Action.Monitor.Check_Out (Task_Id);
   then abort
      begin
         delay To_Duration (Actions (Task_Id).Scope.Start_Delay_Max);
         Deadline := Time_Last;
```

```
begin
   Monitor.Check_In (Task_Id);
   select
      Monitor.Failed (Task_Id);
      Actions (Task_Id).Cleanup.all;
      Atomic_Action.Monitor.Check_Out (Task_Id);
   then abort
      begin
         delay To_Duration (Actions (Task_Id).Scope.Start_Delay_Max);
         Deadline := Time_Last;
```
Asynchronism

Atomic actions in Ada95:

```ada
-- select
  delay until Real_Deadline; -- based on Max_Elapse & Deadline
raise Time_Out_State;
then abort
  Actions (Task_Id).Action.all;
end select;
Atomic_Action.Monitor.Check_Out (Task_Id);
exception
  when Time_Out_State        => Monitor.Fail (Time_Out);
  when Late_Activation_State => Monitor.Fail (Late_Activation);
  when others                => Monitor.Fail (Other_Exception);
end;
end Action_Task;
```

Atomic actions in Ada95:

```ada
protected Monitor is
  entry Check_In (Task_Id : in Task_Ids);
  entry Fail (Condition : in Atomic.Condition);
  entry Failed (Task_Id : in Task_Ids);
  -- blocking until Fail is called
  entry Check_Out (Task_Id : in Task_Ids);
  -- blocking until all parts are completed or all are cleaned up
  entry Action_Result (Condition : out Atomic.Condition);
private
  Check_List      : Task_List        := Check_List_Out;
  State           : Atomic_State     := Checking_In;
  Final_Condition : Atomic.Condition := Succeeded;
end Monitor;
```

Mechanism can be extended to allow any task to dedicate itself to one part of an atomic action as well as to allow for nested atomic actions: ⇢ laboratories!

Backward error recovery in real-time environments

Since once some parts of the action might have failed:
- some action-parts might re-execute the same method again,
- or execute alternative methods, even if the original method was locally valid.

In a real-time environment, failed atomic actions are often identical with some kind of a disaster:
- Tracking back and re-trying the same atomic action with modified parameters or methods is rarely useful, considering timing constraints.
- A ‘mode change’ and a complete different set of (atomic) actions (and goals!) might be more useful in many cases.
Asynchronism

Atomic actions:
Forward error recovery in real-time environments

- Backtracking is often hardly possible in real-time systems!
- ... and even if it is, it would be rarely useful in a real-time context (see above).

- Forward error recovery is more common in real-time systems.
  (more on forward error recovery in chapter 10)

Asynchronous Transfer of Control vs. Interrupts

From interrupts (sub-process level) employed in ...
- ... communication with slow / asynchronous / sporadic devices
- ... sampling / control loops
- ... closely coupled reflective systems

... to asynchronous transfer of control (process level):
- Error recovery — supporting atomic actions and forward recovery
- Mode changes — sudden changes from 'normal' operations to emergency measures
- Partial / imprecise computations — whenever timeliness is more important than precision
- Operator intervention — User triggered mode changes

Real-time Java

supplies three basic features:

1. Binding of handlers to internal and external asynchronous events.
   1-a Specifying asynchronous events
   1-b Specifying adequate handlers

2. Asynchronous transfer of control.

3. Asynchronous termination of threads.

Real-time Java: Asynchronous events

Whenever an instance of AsyncEvent occurs:
- all .run() methods of all instances of AsyncEventHandler, which are bound to this AsyncEvent are scheduled for execution.
- multiple AsyncEvents may have different impacts on the scheduling.
- an event counter (FireCount) is supplied.
- AsyncEvents and AsyncEventHandler may be created and used by any program logic.
- More than one handler can be attached to one event.
- More than one event can be attached to one handler.

☞ Flexible, but handled as a schedulable object
Real-Time & Embedded Systems

Asynchronism

Real-time Java: Asynchronous events

```java
public class AsyncEvent {
    public AsyncEvent ();
    public synchronized void addHandler (AsyncEventHandler handler);
    public synchronized void removeHandler (AsyncEventHandler handler);
    public void setHandler (AsyncEventHandler handler);
    public void bindTo (java.lang.String happening);
    public ReleaseParameters createReleaseParameters ();
    public void fire ();
}
```

```java
public abstract class AsyncEventHandler implements Schedulable {
    public AsyncEventHandler (SchedulingParameters scheduling, ReleaseParameters release, MemoryParameters memory, MemoryArea area, ProcessingGroupParameters group);
    public void addToFeasibility ();
    public void removeFromFeasibility ();
    protected int getAndClearPendingFireCount ();
    public abstract void handleAsyncEvent (); -- called by run while FireCount ≠ 0
    public final void run ();
}
```

```java
public class BoundAsyncEventHandler extends AsyncEventHandler {
    public BoundAsyncEventHandler (); // inherit modes from the current thread
    public BoundAsyncEventHandler (SchedulingParameters scheduling, ReleaseParameters release, MemoryParameters memory, MemoryArea area, ProcessingGroupParameters group, boolean nonheap, java.lang.Runnable logic);
}
```

☞ bind the current thread to an AsyncEventHandler (otherwise the AsyncEventHandler will execute in a thread of its own).

Real-Time & Embedded Systems

Asynchronism

Real-time Java: Asynchronous events

While an AsyncEvent just schedules a handler for execution, other event classes are needed to alter the control flow more directly:

Standard Java:
The `InterruptedException` indicates the wish to interrupt a thread, which itself needs to call the `isInterrupted` method to find out, whether it is supposed to be interrupted. If the thread is currently executing and ignoring the flag:
☞ there is no effect on the actual control flow!
If the thread, which is to be interrupted is currently blocking:
☞ it is activated and receives an `InterruptedException`.
☞ too weak to be employed in an asynchronous transfer of control!
Real-time Java: Asynchronously interrupting exception

While an AsyncEvent just schedules a handler for execution, other event classes are needed to alter the control flow more directly:

Real-time Java:
the AsynchronouslyInterruptedException might intercept a control flow directly, while:

- The code regions, which are interruptible need to be indicated.
- Synchronized blocks, task creations and finalizations are not interruptible
- The response time of a thread must be within 'a bounded number of bytecodes' (which itself need to be documented).
- If the interruptible thread is currently blocked (in a java.io.* operation) then the thread is either released or kept in the block (definition of a reasonable blocking time or situation is implementation dependent).

AsynchronouslyInterruptedException

public class AsynchronouslyInterruptedException extends java.lang.InterruptedException {
    ...
    public synchronized void disable();
    public synchronized boolean enable();
    public synchronized boolean fire();
    public boolean doInterruptible (Interruptible logic);
    public happened (boolean propagate);
    public static AsynchronouslyInterruptedException getGeneric();
    // returns the AsynchronouslyInterruptedException which is generated when RealtimeThread.interrupt() is invoked
    public void propagate();
}

Real-time Java: Asynchronously interrupting exception

Cases of operations if a AsynchronouslyInterruptedException (AIE) occurs:

1. If control is in a synchronized section (within an interruptible section):
   * the AIE is put into a pending state

2. If control is in an interruptible section:
   * the control is transferred to the nearest dynamically enclosing catch clause, which is in a non-interruptible section.

3. If control is in wait, sleep, or join:
   * the thread is activated and 1. or 2. is followed.

4. If control is in a non-interruptible section:
   * nothing happens until an interruptible section is reached

5. If an interruptible section is reached while propagating an exception:
   * the original exception is discarded (!) and replaced by the AIE.

import NonInterruptibleServices.*;
public class InterruptibleService {
    public AIE stopNow = AIE.getGeneric();
    public boolean Service() throws AIE {
        try {
            // code interdispersed with calls to NonInterruptibleServices
        } catch AIE AI {
            if (stopNow.happened (true)) {
                // handle the ATC
            }
        }
    }
}
Real-time Java: Asynchronously interrupting exception

```
import NonInterruptibleServices.*;
public class InterruptibleService {
   public AIE stopNow = AIE.getGeneric();
   public boolean Service() throws AIE {
      try {
         // code interdispersed with calls to NonInterruptibleServices
      }
      catch AIE AI {
         if (stopNow.happened (false))               { // handle the ATC
            // handle the stopNow
         } else { // cleanup
            AI.propagate
         }
      }
   }
```

interruptible code section

handle the stopNow
otherwise: clean up and propagate

Real-time Java: Asynchronous exception propagation

While asynchronously interrupting event handling is part of the standard exception handling mechanism, there are nevertheless differences in exception propagation:

- A standard exception is not propagated by default when caught by any 'catch' statement (an explicit re-raising of the exception is necessary to pass it on).
- A `AsynchronouslyInterruptedException` is propagated further, even if caught (an explicit propagation stop is necessary to avoid its further propagation).

☞ realized in the `AIE.happened` method

Rationale:
- local catch-all clauses might not be aware of a potential asynchronous interrupting exception.
- practical compromise, but destroying some integrity of the Java-exception concept.

Real-time Java: Interruptible interface

```
public interface Interruptible {
   public void interruptAction (AsynchronouslyInterruptedException exception);
   public void run (AsynchronouslyInterruptedException exception) throws AsynchronouslyInterruptedException;
}
```

• An object may declare an interruptible method explicitly by implementing the above interface.

☞ only purpose: pass this implementation to a `doInterruptible` method of a AIE class.

☞ The `AIE.doInterruptible` can only be called by one thread at a time!

Real-time Java: Timeout on actions

```
public class Timed extends AsynchronouslyInterruptedException implements java.io.Serializable {
   public Timed (HighResolutionTime time) throws IllegalArgumentException;
   public boolean doInterruptible (Interruptible logic);
   public void resetTime (HighResolutionTime time);
}
```

• The timer is started sometime between the invocation of the `doInterruptible` itself and the `run` method of the interruptible interface.
• A generic `interrupt()` is thrown at the expiration of the timer.
• `time` can be absolute or relative.
Asynchronism

Ada95

provides:

• Exception handling (synchronous only).
• Asynchronous transfer of control
• Task aborts
• Interrupt handling (close to the hardware).

a set of different methods to handle different kinds of events!

package Ada.Interrupts is
    type Interrupt_ID is implementation-defined;
    type Parameterless_Handler is access protected procedure;
    function Is_Reserved (Interrupt : Interrupt_ID) return Boolean;
    function Is_Attached (Interrupt : Interrupt_ID) return Boolean;
    function Current_Handler (Interrupt : Interrupt_ID) return Parameterless_Handler;
    procedure Attach_Handler (New_Handler : in Parameterless_Handler; ...
        Interrupt   : in Interrupt_ID);
    procedure Detach_Handler (Interrupt   : in Interrupt_ID);
    function Reference (Interrupt : Interrupt_ID) return System.Address;
end Ada.Interrupts;

Protected procedures need to qualify as an interrupt handler:
1. use pragma Interrupt_Handler
2. let the compiler evaluate the suitability of the routine as an interrupt handler.

Protected procedures can also be attached statically to an interrupt:
package Ada.Interrupts is
  type Interrupt_ID is implementation-defined;
  type Parameterless_Handler is access protected procedure;
  function Is_Reserved (Interrupt : Interrupt_ID) return Boolean;
  function Is_Attached (Interrupt : Interrupt_ID) return Boolean;
  function Current_Handler (Interrupt : Interrupt_ID) return Parameterless_Handler;
  procedure Attach_Handler (New_Handler : in Parameterless_Handler; Interrupt : in Interrupt_ID);
  procedure Detach_Handler (Interrupt : in Interrupt_ID);
  function Reference (Interrupt : Interrupt_ID) return System.Address;
end Ada.Interrupts;

The mechanism to invoke an interrupt handler may be different from calling a protected procedure from a task.

Implementation advice: Whenever possible, the implementation should allow interrupt handlers to be called directly by the hardware.

Direct access to the invocation address:
May be used to connect task-entries to interrupts
\textit{== risky!} — use with special care.

Metrics: The implementation shall document to worst case overhead for an interrupt handler invocation (in clock cycles).

asynchronous_select ::= select
  triggering_alternative                        then abort                           abortable_part                        end select;

triggering_alternative ::= triggering_statement \[sequence_of_statement\]trigging_statement   ::= entry_call_statement | delay_statement abortable_part         ::= sequence_of_statements

== cannot contain an accept statement.
Asynchronism

Ada95: Asynchronous Transfer of Control

Execute the trigger (entry.call or delay), then:

1. If the trigger is going through and can be completed:
   the optional statements following the trigger are executed and the select statement is completed
   (the abortable part is never started).

2. If the trigger is blocked or requeued to a blocked entry:
   the statements in the abortable part are executed:

   2-a If the abortable part completes before the trigger
       is completed, an attempt is made to revoke the triggering statement.
       The select statement is completed after the cancelled or completed triggering statement.

   2-b If the trigger is completed before the abortable part is completed,
       the abortable part is stopped, the optional statements following the trigger are executed
       and the select statement is completed.

Exception handling:

Both parts of a select-then-abort statement
can raise exceptions, but ...

☞ ... in case of an interruption of the abortable part,
the exceptions from the abortable part are lost!

select <entry-call | delay>  
[ ... statements ... ]
then abort  
... statements ...  
end select;

task body A is
T : Time;
D : Duration;
begin
...  
select  
delay until T;  
then abort  
delay D;  
end select;  
end A;

☞ are these equivalent?

task body B is
T : Time;
begin
...  
select  
delay D;  
then abort  
delay until T;  
end select;  
end A;

... are these equivalent?.
Asynchronism

Ada95: Asynchronous Transfer of Control

Task body A is

\[
\begin{align*}
T : \text{Time};& \\
\text{begin};& \\
\text{select};& \\
\text{delay until } T; & \\
\text{then abort};& \\
\text{Server.Entry1;} & \\
\text{end select;} & \\
\text{end A;}& \\
\end{align*}
\]

... if rendezvous starts before but finishes after timeout.

Task body B is

\[
\begin{align*}
T : \text{Time};& \\
\text{begin};& \\
\text{select};& \\
\text{Server.Entry1;} & \\
\text{then abort};& \\
\text{delay until } T; & \\
\text{end select;} & \\
\text{end B;}& \\
\end{align*}
\]

... if rendezvous starts and completes before timeout.

Task body C is

\[
\begin{align*}
T : \text{Time};& \\
\text{begin};& \\
\text{select};& \\
\text{Server.Entry1;} & \\
\text{or};& \\
\text{delay until } T; & \\
\text{end select;} & \\
\text{end C;}& \\
\end{align*}
\]

... if rendezvous starts and completes before timeout.
Real-Time & Embedded Systems

Asynchronism

Ada95: Asynchronous Transfer of Control

task body A is
T : Time;
begnin
begin
select
delay until T;
then abort
Server.Entry1;
end select;
end A;

end body B is
T : Time;
begnin
begin
select
Server.Entry1;
then abort
server.Entry1;
end select;
end B;

end body C is
T : Time;
begnin
begin
select
Server.Entry1;
then abort
server.Entry1;
end select;
end C;

... timeout occurs before the rendezvous starts.

Asynchronism in Ada95 and Real-time Java:
(Common features)

- ATC-enabled regions must be declared.
- Some regions are always deferred from asynchronous transfer of control (task/thread communication / finalization).
- Exceptions from the run-time environment as well as user-defined exceptions are supported.
- Asynchronous events may be triggered by the environment as well as from a task.
Asynchronism

Asynchronism in Ada95 and Real-time Java:

(Differences)

- Mechanisms:
  - In **Real-time Java** asynchronism is embedded into the synchronous exception scheme
  - In **Ada95** interrupts are interrupts and ATC is embedded in the 'select' scheme.

- Asynchronous transfer of control regions:
  - **Real-time Java** declares ATC-enabled regions per method and any asynchronous event is deferred until the next ATC-enabled method is executing.
  - **Ada95** assumes that all code which is called from within an ATC-enabled region is ATC-enabled.

- Handler identification:
  - **Real-time Java** delivers asynchronous events to all enrolled handlers and propagates an asynchronous interrupting event through the closest handlers.
  - **Ada95** delivers an interrupt to one global handler and each ATC-enabled region has exactly one exit point.

Summary

Asynchronism

- **Interrupts / Signals**
  - Device / system / language / operating-system level interrupt control
  - Characteristics of interrupts and signals

- **Exceptions**
  - Exception classes / granularity / parametrisation / propagation
  - Resumption and termination, specific language issues

- **Atomic Actions**
  - Definition / requirements / failure cases / implementation / error recovery

- **Asynchronous transfer of control / Interrupts in context**
  - Interrupts and ATC in real-time Java and Ada95

References for this chapter

[Ada95RM] (link to on-line version)
Ada Working Group
ISO/IEC JTC1/SC 22/WG 9
Ada 95 Reference Manual
– Language and Standard Libraries

[Bollella01]
Greg Bollella, Ben Brosgol, Steve Furr, David Hardin, Peter Dibble, James Gosling, Mark Turnbull & Rudy Belliardi
The Real-Time Specification for Java
http://www.rtj.org

[Burns01]
Alan Burns and Andy Wellings
Real-Time Systems and Programming Languages
Addison Wesley, third edition, 2001

[Cohen96]
Norman H. Cohen
Ada as a second language
McGraw-Hill series in computer science, 2nd edition

all references and links are available on the course page
Synchronization methods

• Shared memory based synchronization
  - Semaphores
  - Conditional critical regions
  - Monitors
  - Mutexes & conditional variables
  - Synchronized methods
  - Protected objects
  - 'C', POSIX — Dijkstra
  - Edison (experimental)
  - Modula-1, Mesa — Dijkstra, Hoare, …
  - POSIX
  - Real-time Java
  - Ada95

• Message based synchronization
  - Asynchronous messages
  - Synchronous messages
  - Remote invocation, remote procedure call
  - Synchronization in distributed systems
  - e.g. POSIX, …
  - e.g. Ada95, CHILL, Occam2
  - e.g. Ada95, …
  - e.g. CORBA, …

Synchronization in real-time systems

There are many concurrent entities in a real-time systems:

• Interrupt handlers
• Tasks
• Dispatchers
• Timers
• …

... and ... real-time systems are often complex and is possibly expanded at a later stage ...

Thus all data is declared ...

... either local (and protected by language-, os-, or hardware-mechanisms)
... or it is 'out in the open' and all access need to be synchronized!

Synchronization: the run-time overhead?

Is the potential overhead justified for simple data-structures:

```plaintext
int i;
...
i++;  {in one thread}  |  i=0;  {in another thread}
```

• Are those operations atomic?
• Do we really need to introduce full featured synchronization methods here?

Depending on the hardware and the compiler, it might be atomic, it might be not:

• Handling a 64-bit integer on a 8- or 16-bit controller will not be atomic
  ... but perhaps it is an 8-bit integer.

• Any manipulations on the main memory will not be atomic
  ... but perhaps it is a register.

• Broken down to a load-operate-store cycle, the operations will not be atomic
  ... but perhaps the processor supplies atomic operations for the actual case.

• Assuming that all 'perhapses' are applying: how to expand this code?
Synchronization

Synchronization in real-time systems

```c
int i;
    ....
i++;  {in one thread}  |  i=0;  {in another thread}
```

Unfortunately: the chances that such programming errors turn out are usually small and some implicit by chance synchronization in the rest of the system might prevent them at all.

- Many effects stemming from asynchronous memory accesses are interpreted as (hardware) ‘glitches’, since they are rare and effect usually only some parts of the data.
- On assembler level: synchronization by employing knowledge about the atomicity of CPU-operations and interrupt structures is nevertheless possible and done frequently.

In anything higher than assembler level on small, predictable µcontrollers:

☞ Measures for synchronization are required!

Some synchronization terms:

- **Condition synchronization:** synchronize a task with an event given by another task.

- **Critical sections:**
  - code fragments which contain access to shared resources and need to be executed without interference with other critical sections, sharing the same resources.

- **Mutual exclusion:**
  - protection against asynchronous access to critical sections.

- **Atomic operations:**
  - the set of operations, which atomicity is guaranteed by the underlying system (e.g. hardware).

☞ there must be a set of atomic operations to start with!

Synchronization by flags

Word-access atomicity:

Assuming that any access to a word in the system is an atomic operation:

- e.g. assigning two values (not wider than the size of word) to a memory cell simultaneously:

  ```c
  Task 1:  x := 0;
  Task 2:  x := 5;
  ```

will result in either \( x = 0 \) xor \( x = 5 \) — and no other value is ever observable.
**Condition synchronization by flags**

```
var Flag : boolean := false;

process P1;
    statement X;
    repeat until Flag;
    statement Y;
end P1;

process P2;
    statement A;
    Flag := true;
    statement B;
end P2;
```

Sequence of operations: \([A \mid X] \rightarrow [B \mid Y]\)

---

**Synchronization by semaphores**

(∗Dijkstra 1968∗)

Assuming further that there is a shared memory between two processes:

- A set of processes agree on a (word-size) atomic variable operating as a flag to indicate synchronization conditions:
  
  Memory flag method is ok for simple condition synchronization, but ...

  ⇒ ... is not sufficient for general mutual exclusion in critical sections!

  ⇒ ... busy-waiting is required to poll the synchronization condition!

  ⇒ More powerful synchronization operations are required for critical sections

```
var sync : semaphore := 0;

process P1;
    statement X;
    wait (sync);
    statement Y;
end P1;

process P2;
    statement A;
    signal (sync);
    statement B;
end P2;
```

Sequence of operations: \([A \mid X] \rightarrow [B \mid Y]\)
Synchronization

Mutual exclusion by semaphores

```ada
var mutex : semaphore := 1;

process P1;
statement X;
wait (mutex);  
statement Y;
signal (mutex); 
statement Z;
end P1;

process P2;
statement A;
wait (mutex); 
statement B;
signal (mutex); 
statement C;
end P2;
```

Sequence of operations: [A | X] → [B → Y xor Y → B] → [C | Z]

Types of semaphores:

- **General semaphores** (counting semaphores): non-negative number; (range limited by the system)

  - `P` and `V` increment and decrement the semaphore by one.

- **Binary semaphores**: restricted to [0, 1]; Multiple `V` (Signal) calls have the same effect than 1 call.
  - Binary semaphores are sufficient to create all other semaphore forms.
  - Atomic ‘test-and-set’ operations at hardware level are usually binary semaphores.

- **Quantity semaphores**: The increment (and decrement) value for the semaphore is specified as a parameter with `P` and `V`.

Semaphores in Ada95

```ada
package Ada.Synchronous_Task_Control is
  type Suspension_Object is limited private;
  procedure Set_True  (S : in out Suspension_Object);
  procedure Set_False (S : in out Suspension_Object);
  function Current_State (S : Suspension_Object) return Boolean;
  procedure Suspend_Until_True (S : in out Suspension_Object);
private
  _ -- not specified by the language
  end Ada.Synchronous_Task_Control;
```

- only one task can be blocked at `Suspend_Until_True!` (strict version of a binary semaphore) 
  (Program_Error will be raised with the second task trying to suspend itself)
- no queues! => minimal run-time overhead

```ada
package Ada.Synchronous_Task_Control is
  type Suspension_Object is limited private;
  procedure Set_True  (S : in out Suspension_Object);
  procedure Set_False (S : in out Suspension_Object);
  function Current_State (S : Suspension_Object) return Boolean;
  procedure Suspend_Until_True (S : in out Suspension_Object);
private
  _ -- not specified by the language
  end Ada.Synchronous_Task_Control;
```

- only one task can be blocked at `Suspend_Until_True!` (strict version of a binary semaphore) 
  (Program_Error will be raised with the second task trying to suspend itself)
- no queues! => minimal run-time overhead
### Semaphores in POSIX

```c
int sem_init (sem_t *sem_location, int pshared, unsigned int value);
int sem_destroy (sem_t *sem_location);
int sem_wait      (sem_t *sem_location);
int sem_trywait   (sem_t *sem_location);
int sem_timedwait (sem_t *sem_location, const struct timespec *abstime);
int sem_post      (sem_t *sem_location);
int sem_getvalue  (sem_t *sem_location, int *value);
```

- **generate semaphore for usage between processes**
  (otherwise for threads of the same process only)

### Semaphores in POSIX

```c
int sem_init (sem_t *sem_location, int pshared, unsigned int value);
int sem_destroy (sem_t *sem_location);
int sem_wait      (sem_t *sem_location);
int sem_trywait   (sem_t *sem_location);
int sem_timedwait (sem_t *sem_location, const struct timespec *abstime);
int sem_post      (sem_t *sem_location);
int sem_getvalue  (sem_t *sem_location, int *value);
```

- delivers the number of waiting processes as a negative integer,
  if there are processes waiting on this semaphore

### Synchronization

#### Semaphores in POSIX

```c
void allocate (priority_t P) {
    sem_wait (&mutex);
    if (busy) {
        sem_post (&mutex);
        sem_wait (&cond[P]);
    }    
    busy = 1;
    sem_post (&mutex);
    ...
}
```

```c
void deallocate (priority_t P) {
    sem_wait (&mutex);
    busy = 0;
    sem_getvalue (&cond[high], &waiting);
    if (waiting > 0) {
        sem_post (&cond[high]);
    } else {
        sem_getvalue (&cond[low], &waiting);
        if (waiting < 0) {
            sem_post (&cond[low]);
        } else {
            sem_post (&mutex);
        }
    }
}
```

- could raise a `Program_Error` in Ada95.
- produces a potential deadlock when implemented with general semaphores.
- Deadlocks can be generated by all kinds of synchronization methods.
Criticism of semaphores

- Semaphores are not bound to any resource or method or region
  - Adding or deleting a single semaphore operation some place might stall the whole system

- Semaphores are scattered all over the code
  - Hard to read, error-prone

- Semaphores are considered not adequate for the real-time domain.
  (all concurrent and real-time languages offer more abstract and safer synchronization methods).

Conditional critical regions

Basic idea:

- Critical regions are a set of code sections in different processes, which are guaranteed to be executed in mutual exclusion:
  - Shared data structures are grouped in named regions and are tagged as being private resources.
  - Processes are prohibited from entering a critical region, when another process is active in any associated critical region.

- Condition synchronisation is provided by guards:
  - When a process wishes to enter a critical region it evaluates the guard (under mutual exclusion). If the guard evaluates false, the process is suspended / delayed.
  - As with semaphores, no access order can be assumed.

```plaintext
buffer : buffer_t;
resource critial_buffer_region : buffer;

process producer;
  loop
    region critial_buffer_region
      when buffer.size < N do
        -- place in buffer etc.
      end region
    end loop;
  end producer

process consumer;
  loop
    region critial_buffer_region
      when buffer.size > 0 do
        -- take from buffer etc.
      end region
    end loop;
  end consumer
```

Criticism of conditional critical regions

- All guards need to be re-evaluated, when any conditional critical region is left:
  - All involved processes are activated to test their guards
  - There is no order in the re-evaluation phase "potential livelocks"

- As with semaphores the conditional critical regions are scattered all over the code.
  - On a larger scale: same problems as with semaphores.

The language Edison uses conditional critical regions for synchronization in a multiprocessor environment (each process is associated with exactly one processor).
Synchronization

Monitors

(Modula-1, Mesa — Dijkstra, Hoare)

Basic idea:

- Collect all operations and data-structures shared in critical regions in one place, the monitor.
- Formulate all operations as procedures or functions.
- Prohibit access to data-structures, other than by the monitor-procedures.
- Assure mutual exclusion of the monitor-procedures.

Monitor

monitor buffer;
export append, take;
var (* declare protected vars *)
procedure append (I : integer);
    ...
procedure take (var I : integer);
    ...
begin
    (* initialisation *)
end;

How to realize conditional synchronization?

Monitors with condition synchronization

(Hoare)

Hoare-monitors:

- Condition variables are implemented by semaphores (Wait and Signal).
- Queues for tasks suspended on condition variables are realized.
- A suspended task releases its lock on the monitor, enabling another task to enter.

☞ More efficient evaluation of the guards:
    the task leaving the monitor can evaluate all guards and the right tasks can be activated.

☞ Blocked tasks may be ordered and livelocks prevented.

Monitor with condition synchronization

monitor buffer;
    export append, take;
    var Buf [ ] of integer;
    top, base : 0..size-1;
    NumberInBuffer : integer;
    spaceavailable, itemavailable : condition;
    procedure append (I : integer);
        begin
            if NumberInBuffer = size then
                wait (spaceavailable);
            begin
                Buf[top] := I; NumberInBuffer := NumberInBuffer+1;
                top := (top+1) mod size;
                signal (itemavailable)
            end;
        end;
    ...
    ...

Monitors with condition synchronization

procedure take (var I : integer);
begin
  if NumberInBuffer = 0 then
    wait (itemavailable);
  end if;
  I := BUF[base];
  base := (base+1) mod size;
  NumberInBuffer := NumberInBuffer-1;
  signal (spaceavailable);
end take;

begin (* initialisation *)
  NumberInBuffer := 0;
  top := 0;
  base := 0
end;

The signalling and the waiting process are both active in the monitor!

Suggestions to overcome the multiple-tasks-in-monitor-problem:

- A signal is allowed only as the last action of a process before it leaves the monitor.
- A signal operation has the side-effect of executing a return statement.
- Hoare, Modula-1, POSIX: a signal operation which unblocks another process has the side-effect of blocking the current process; this process will only execute again once the monitor is unlocked again.
- A signal operation which unblocks a process does not block the caller, but the unblocked process must gain access to the monitor again.

Monitors in Modula-1

• wait (s, r):
  delays the caller until condition variable s is true (r is the rank (or ‘priority’) of the caller).
• send (s):
  If a process is waiting for the condition variable s, then the process at the top of the queue of the highest filled rank is activated (and the caller suspended).
• awaited (s):
  check for waiting processes on s.

INTERFACE MODULE resource_control;
DEFINER allocate, deallocate;
VAR busy : BOOLEAN; free : SIGNAL;
PROCEDURE allocate;
BEGIN
  IF busy THEN WAIT (free) END;
  busy := TRUE;
END;
PROCEDURE deallocate;
BEGIN
  busy := FALSE;
  SEND (free); -- or: IF AWAITED (free) THEN SEND (free);
END;
BEGIN
  busy := false;
END.
Monitors in ‘C’ / POSIX
(types and creation)

Synchronization between POSIX-threads:

typedef ... pthread_mutex_t;
typedef ... pthread_mutexattr_t;
typedef ... pthread_cond_t;
typedef ... pthread_condattr_t;

int pthread_mutex_init (      pthread_mutex_t *mutex,
                             const pthread_mutexattr_t *attr);

int pthread_mutex_destroy (      pthread_mutex_t *mutex);

int pthread_cond_init (      pthread_cond_t      *cond,
                               const pthread_condattr_t  *attr);

int pthread_cond_destroy (      pthread_cond_t      *cond);

…

Attributes include:
• semantics for trying to lock a mutex which
  is locked already by the same thread
• sharing of mutexes and
  condition variables between processes
• priority ceiling
• clock used for timeouts
• … … …

Monitors in ‘C’ / POSIX
(operators)

Synchronization between POSIX-threads:

int pthread_mutex_lock (      pthread_mutex_t *mutex);

int pthread_mutex_trylock   (      pthread_mutex_t *mutex);

int pthread_mutex_timedlock (      pthread_mutex_t *mutex,
                                     const struct timespec     *abstime);

int pthread_mutex_unlock (      pthread_mutex_t *mutex);

int pthread_cond_wait (      pthread_cond_t      *cond,
                               pthread_mutex_t     *mutex);

int pthread_cond_timedwait  (      pthread_cond_t      *cond,
                                     pthread_mutex_t     *mutex,
                                     const struct timespec     *abstime);

int pthread_cond_signal (      pthread_cond_t      *cond);

int pthread_cond_broadcast (      pthread_cond_t      *cond);

...
Monitors in ‘C’ / POSIX
(operators)

```c
int pthread_mutex_lock (      pthread_mutex_t     *mutex);
int pthread_mutex_trylock   (      pthread_mutex_t     *mutex);
int pthread_mutex_timedlock (      pthread_mutex_t     *mutex,                             const struct timespec     *abstime);
int pthread_mutex_unlock (      pthread_mutex_t     *mutex);
int pthread_cond_wait (      pthread_cond_t      *cond,                             pthread_mutex_t     *mutex);int pthread_cond_timedwait  (      pthread_cond_t      *cond,                                   pthread_mutex_t     *mutex,                              const struct timespec     *abstime);
int pthread_cond_signal (      pthread_cond_t      *cond);
int pthread_cond_broadcast (      pthread_cond_t      *cond);
```

Synchronization

Monitors in ‘C’ / POSIX
(operators)

```c
int pthread_mutex_lock (      pthread_mutex_t     *mutex);
int pthread_mutex_trylock   (      pthread_mutex_t     *mutex);
int pthread_mutex_timedlock (      pthread_mutex_t     *mutex,                             const struct timespec     *abstime);
int pthread_mutex_unlock (      pthread_mutex_t     *mutex);
int pthread_cond_wait (      pthread_cond_t      *cond,                             pthread_mutex_t     *mutex);int pthread_cond_timedwait  (      pthread_cond_t      *cond,                                   pthread_mutex_t     *mutex,                              const struct timespec     *abstime);
int pthread_cond_signal (      pthread_cond_t      *cond);
int pthread_cond_broadcast (      pthread_cond_t      *cond);
```

Synchronization

Monitors in ‘C’ / POSIX
(operators)

```c
#define BUFF_SIZE 10
typedef struct {
    pthread_mutex_t mutex;
    pthread_cond_t buffer_not_full;   pthread_cond_t buffer_not_empty;
    int             count, first, last;   int             buf[BUFF_SIZE];} buffer;
```
Monitors in ‘C’ / POSIX

(int item, buffer *B) {

PTHREAD_MUTEX_LOCK (&B->mutex);
while (B->count == BUFF_SIZE) {
    PTHREAD_COND_WAIT (
        &B->buffer_not_full,
        &B->mutex);
}

PTHREAD_MUTEX_UNLOCK (&B->mutex);
return 0;

} 

Monitors in Real-time Java

Java provides two mechanisms to construct monitors:

• Synchronized methods and code blocks
  all methods and code blocks which are using the synchronized tag
  are mutually exclusive with respect to the addressed class.

• Notification methods: wait, notify, and notifyAll
  can be used only in synchronized regions and are waking any or all threads,
  which are waiting in the same synchronized object.

Considerations:

1. Synchronized methods and code blocks:
   • In order to implement a monitor all methods in an object need to be synchronized.
     ✦ any other standard method can break the monitor and enter at any time.
   • Methods outside the monitor-object can synchronize at this object.
     ✦ it is impossible to analyse a monitor locally, since lock accesses can exist all over the system.
   • Static data is shared between all objects of a class.
     ✦ access to static data need to be synchronized over the whole class.

2. Notification methods: wait, notify, and notifyAll
   • wait suspends the thread and releases the local lock only
     ✦ nested wait-calls will keep all enclosing locks.
   • notify and notifyAll does not release the lock.
     ✦ methods, which are activated via notification need to wait for lock-access.
   • wait-suspended threads are hold in a queue, thus notifyAll is waking the threads in order
     ✦ livelocks are prevented at this level (in opposition to Java).
   • There are no explicit conditional variables.
     ✦ every notified thread needs
to wait for the lock to be released and to re-evaluate its entry condition
Synchronization

Monitors in Real-time Java
(multiple-readers-one-writer-example)

each of the readers uses these monitor.calls:  

\[
\text{startRead ();} \\
\text{// read the shared data only} \\
\text{stopRead ();}
\]

each of the writers uses these monitor.calls:

\[
\text{startWrite ();} \\
\text{// manipulate the shared data} \\
\text{stopWrite ();}
\]

☞ construct a monitor, which allows multiple readers  
or one writer  
at a time inside the critical regions

Synchronization

Monitors in Real-time Java
(multiple-readers-one-writer-example: wait-notifyAll method)

```java
public class ReadersWriters {
    private int readers        = 0;
    private int waitingWriters = 0;
    private boolean writing        = false;
    ...  
    public synchronized void StartWrite () throws InterruptedException   {
        while (readers > 0 || writing)       {
            waitingWriters++;       
            wait();       
            waitingWriters--;  
        }      
        writing = true;   }
    public synchronized void StopWrite()   {
        writing = false;
        if (readers == 0) notifyAll();
    }
    ...  
}
```

whenever a synchronized region is left:  
• all thread are notified  
• all threads are re-evaluating their guards
Synchronization

Monitors in Real-time Java

Standard monitor solution:

- declare the monitored data-structures private to the monitor object (non-static).
- introduce a class `ConditionVariable`:
  ```java
  public class ConditionVariable {
      public boolean wantToSleep = false;
  }
  ```
- introduce synchronization-scopes in monitor-methods:
  ✓ synchronize on the adequate conditional variables first and
  ✓ synchronize on the monitor-object second.
- make sure that all methods in the monitor are implementing the correct synchronizations.
- make sure that no other method in the whole system is synchronizing on this monitor-object.

```
public class ReadersWriters {
    private int readers = 0;
    private int waitingReaders = 0;
    private int waitingWriters = 0;
    private boolean writing = false;
    ConditionVariable OkToRead = new ConditionVariable();
    ConditionVariable OkToWrite = new ConditionVariable();
    ...
}
```

```
  public void StartWrite () throws InterruptedException {
    synchronized (OkToWrite) {
      synchronized (this) {
        if (writing | readers > 0) {
          waitingWriters++;
          OkToWrite.wantToSleep = true;
        } else {
          writing = true;
          OkToWrite.wantToSleep = false;
        }
      }
      if (OkToWrite.wantToSleep) OkToWrite.wait();
    }
  }
```

```
  public void StopWrite () {
    synchronized (OkToRead) {
      synchronized (OkToWrite) {
        synchronized (this) {
          if (waitingWriters > 0) {
            waitingWriters--;
            OkToWrite.notify(); // wakeup one writer
          } else {
            writing = false;
            OkToRead.notifyAll (); // wakeup all readers
            readers = waitingReaders;
            waitingReaders = 0;
          }
        }
      }
    }
  }
```
Monitors in Real-time Java

```java
public void StartRead () throws InterruptedException {
    synchronized (OkToRead) {
        synchronized (this) {
            if (writing | waitingWriters > 0) {
                waitingReaders++;
                OkToRead.wantToSleep = true;
            } else {
                readers++;
                OkToRead.wantToSleep = false;
            }
        }
        if (OkToRead.wantToSleep) OkToRead.wait();
    }
}
```}

Monitors in Real-time Java

```java
public void StopRead () {
    synchronized (OkToWrite) {
        synchronized (this) {
            readers--;
            if (readers == 0 & waitingWriters > 0) {
                waitingWriters--;
                OkToWrite.notify();
            }
        }
    }
```}

Object-orientation and synchronization

Since mutual exclusion, notification, and condition synchronization schemes need to be designed and analysed considering the implementation of all involved methods and guards:

☞ new methods cannot be added without re-evaluating the whole class!

In opposition to the general re-usage idea of object-oriented programming, the re-usage of synchronized classes (e.g. monitors) need to be considered carefully.

☞ The parent class might need to be adapted in order to suit the global synchronization scheme.

☞ Inheritance anomaly (Matsuoka & Yonezawa '93)

Methods to design and analyse expandible synchronized systems exist, but are fairly complex and are not provided in any current object-oriented language.
Synchronization

Nested monitor calls

Assuming a thread in a monitor is calling an operation in another monitor and is suspended at a conditional variable there:

☞ the called monitor is aware of the suspension and allows other threads to enter.
☞ the calling monitor is possibly not aware of the suspension and keeps its lock!
☞ the unjustified locked calling monitor reduces the system performance and leads to potential deadlocks.

Suggestions to solve this situation:
- Maintain the lock anyway: e.g. POSIX, Real-time Java
- Prohibit nested procedure calls: e.g. Modula-1
- Provide constructs which specify the release of a monitor lock for remote calls, e.g. Ada95

Criticism of monitors

- Mutual exclusion is solved elegantly and safely.
- Conditional synchronization is on the level of semaphores still
  ☞ all criticism on semaphores apply
  ☞ mixture of low-level and high-level synchronization constructs.

Synchronization by protected objects

Combine
- the encapsulation feature of monitors
- the coordinated entries of conditional critical regions

☞ Protected objects
- all controlled data and operations are encapsulated
- all operations are mutual exclusive
- entry guards are attached to operations
- the protected interface allows for operations on data
- no protected data is accessible (other than by defined operations)
- tasks are queued (according to their priorities)

Protected objects in Ada95

Some read-only operations do not need to be mutual exclusive:

protected type Shared_Data (Initial : Data_Item) is
  function  Read return Data_Item;
  procedure Write (New_Value : in Data_Item);
private
  The_Data : Data_Item := Initial;
end Shared_Data_Item;

☞ protected functions can have ‘in’ parameters only and are not allowed to alter the private data (enforced by the compiler).
☞ protected functions allow simultaneous access (but mutual exclusive with other operations).
- there is no defined priority between functions and other protected operations in Ada95.
Synchronization by protected objects in Ada95

Condition synchronization is realized in the form of protected procedures combined with boolean conditional variables (barriers): entries in Ada95:

```
Buffer_Size : constant Integer := 10;
type Index is mod Buffer_Size;
subtype Count is Natural range 0 .. Buffer_Size;
type Buffer_T is array (Index) of Data_Item;

protected type Bounded_Buffer is
  entry Get (Item : out Data_Item);
  entry Put (Item : in Data_Item);
private   First  : Index := Index'First;   Last   : Index := Index'Last;   Num    : Count := 0;   Buffer : Buffer_T;
end Bounded_Buffer;
```

```
protected body Bounded_Buffer is
  entry Get (Item : out Data_Item)
  when Num > 0 is
    begin
      Item  := Buffer (First);         First := First + 1;         Num   := Num - 1;
    end Get;

  entry Put (Item : in Data_Item)
  when Num < Buffer_Size is
    begin
      Last          := Last + 1;         Buffer (Last) := Item;         Num           := Num + 1;
    end Put;
end Bounded_Buffer;
```

Protected entries are used like task entries:

```
select
  Buffer : Bounded_Buffer;
  Buffer.Put (Some_Data);
or  delay 10.0;
  -- do something after 10 s.
end select;
select
  delay 10.0;
  then abort
    Buffer.Put (Some_Data);
  -- try to enter for 10 s.
end select;
select
  Buffer.Get (Some_Data);
else  -- do something else
end select;
```

Barrier evaluations and task activations:

- on calling a protected entry, the associated barrier is evaluated (only those parts of the barrier which might have changed since the last evaluation).
- on leaving a protected procedure or entry, related barriers with tasks queued are evaluated (only those parts of the barriers which might have been altered by this procedure / entry or which might have changed since the last evaluation).

Barriers are not evaluated while inside a protected object or on leaving a protected function.
Synchronization by protected objects in Ada95

(operations on entry queues)

The count attribute indicates the number of tasks waiting at a specific queue:

protected Blocker is
        entry Proceed;
        when Proceed'count = 5
        or Release is
        begin
                Release := Proceed'count > 0;
                end Proceed;
        end Blocker;

protected body Blocker is
        entry Proceed
        when Proceed'count = 5
        or Release is
        begin
                Release := Proceed'count > 0;
                end Proceed;
        end Blocker;

protected type Broadcast is
        entry Receive (M: out Message);
        procedure Send (M: in  Message);
        private
                New_Message : Message;
                Arrived     : Boolean := False;
        end Blocker;

protected body Broadcast is
        entry Receive (M: out Message)
        when Arrived
        is
        begin
                M := New_Message;
                Arrived := Receive'count > 0;
                end Proceed;
                procedure Send (M: in  Message) is
                begin
                        New_Message := M;
                        Arrived := Receive'count > 0;
                        end Send;
        end Blocker;

Further refinements on task control by:

- **Entry families**: a protected entry declaration can contain a discrete subtype selector, which can be evaluated by the barrier (other parameters cannot be evaluated by barriers) and implements an array of protected entries.

- **Requeue facility**: protected operations can use 'requeue' to redirect tasks to other internal, external, or private entries. The current protected operation is finished and the lock on the object is released.

- **Internal progress first'-rule': internally requeued tasks are placed at the head of the waiting queue!

- **Private entries**: protected entries which are not accessible from outside the protected object, but can be employed as destinations for requeue operations.
Synchronization by protected objects in Ada95
(requeue & private entries)

How to implement a queue, at which every task can be released only once per triggering event?
☞ e.g. by employing two entries:

```
package Single_Release is
  entry Wait;
  procedure Trigger;
private
  Front_Door, Main_Door : Boolean := False;
  entry Queue;
end Single_Release;
```

```
package body Single_Release is
  entry Wait
  when Front_Door is
  begin
    if Wait'Count = 0 then
      Front_Door := False;
      Main_Door  := True;
    end if;
    requeue Queue;
  end Wait;

  entry Queue
  when Main_Door is
  begin
    if Queue'count = 0 then
      Main_Door := False;
    end if;
    end Queue;
    procedure Trigger is
    begin
      Front_Door := True;
    end Trigger;
  end Queue;
end Single_Release;
```

opening the main door before requeuing?

Synchronization by protected objects in Ada95
(restrictions applying to protected operations)

Code inside a protected procedure, function or entry is bound to non-blocking operations (which would keep the whole protected object locked).

Thus the following operations are prohibited:
- entry call statements
- delay statements
- task creations or activations
- calls to sub-programs which contains a potentially blocking operation
- select statements
- accept statements

☞ The requeue facility allows for a potentially blocking operation, but releases the current lock!

Summary

Shared memory based synchronization

General

Criteria:
- level of abstraction
- centralized vs. distributed concepts
- support for consistency and correctness validations
- error sensitivity
- predictability
- efficiency
Real-Time & Embedded Systems

Summary

Shared memory based synchronization

POSIX
- all low level constructs available.
- no connection with the actual data-structures.
- error-prone.
- non-determinism introduced by ‘release some’ semantics of conditional variables (cond_signal).

Semaphores
(atomic P, V ops)
Flags
(atomic word access)
Conditional variables
Synchronized methods
(mutual exclusion)
Monitors
Data structure encapsulation
Conditional critical regions
Guards (barrier)
Protected objects

Real-time Java
- mutual exclusion (synchronized methods) as the only support.
- general notification feature (no conditional variables)
- non-restricted object oriented extension introduces hard to predict timing behaviours.

Semaphores
(atomic P, V ops)
Flags
(atomic word access)
Conditional variables
Synchronized methods
(mutual exclusion)
Monitors
Data structure encapsulation
Conditional critical regions
Guards (barrier)
Protected objects

Modula-1, CHILL
- full monitor implementation (Dijkstra-Hoare monitor concept).
… no more, no less, …
☞ all features of and criticism about monitors apply.

Ada95
- complete synchronization support
- low-level semaphores for very special cases.
- predictable timing (☞ scheduler).
☞ most memory oriented synchronization conditions are realized by the compiler or the run-time environment directly rather then the programmer.
(Ada95 is currently without any mainstream competitor in this field)

Semaphores
(atomic P, V ops)
Flags
(atomic word access)
Conditional variables
Synchronized methods
(mutual exclusion)
Monitors
Data structure encapsulation
Conditional critical regions
Guards (barrier)
Protected objects
Message-based synchronization

- Synchronization model
  - Asynchronous
  - Synchronous
  - Remote invocation

- Addressing (name space)
  - direct communication
  - mail-box communication

- Message structure
  - arbitrary
  - restricted to ‘basic’ types
  - restricted to un-typed communications

Asynchronous messages

If there is a listener:
☞ send the message directly

If the receiver becomes available at a later stage:
☞ the message need to be buffered

Synchronous messages

Delay the sender:
• until the receiver got the message
**Synchronization**

**Message-based synchronization**

**Synchronous messages**

Delay the sender:
- until the receiver got the message
- two asynchronous messages required

- async. send
- async. receive
- async. send
- async. receive

If the receiver becomes available at a later stage:
- messages need to be buffered

- async. send
- async. receive
- async. send
- async. receive

Delay the sender, until:
- a receiver got the message
- a receiver executed an addressed routine

- rem. invoc.
- invocation
**Message-based synchronization**

Remote invocation

Delay the sender, until:
- a receiver got the message
- a receiver executed an addressed routine

Asynchronous remote invocation

Delay the sender, until:
- a receiver becomes available
- a receiver got the message
- a receiver executed an addressed routine
Synchronization

Message-based synchronization

Asynchronous remote invocation

Delay the sender, until:
- a receiver becomes available
- a receiver got the message

Purpose ‘synchronization’: synchronous messages / remote invocations
Purpose ‘in-time delivery’: asynchronous messages / asynchronous remote invocations

☞ ‘Real’ synchronous message passing in distributed systems requires hardware support.
☞ Asynchronous message passing requires the usage of (infinite?) buffers.

• Synchronous communications are emulated by a combination of asynchronous messages in some systems.
• Asynchronous communications can be emulated in synchronized message passing systems by introducing ‘buffer-tasks’ (de-coupling sender and receiver as well as allowing for broadcasts).

Synchronous vs. asynchronous communications

Addressing (name space)

Direct vs. indirect:

send <message> to <process-name>
wait for <message> from <process-name>
send <message> to <mailbox>
wait for <message> from <mailbox>

Asymmetrical addressing:

send <message> to ...
wait for <message>
☞ Client-server paradigm

Communication medium:

<table>
<thead>
<tr>
<th>Connections</th>
<th>Functionality</th>
</tr>
</thead>
<tbody>
<tr>
<td>one-to-one</td>
<td>buffer, queue, synchronization</td>
</tr>
<tr>
<td>one-to-many</td>
<td>multicast</td>
</tr>
<tr>
<td>one-to-all</td>
<td>broadcast</td>
</tr>
<tr>
<td>many-to-one</td>
<td>local server, synchronization</td>
</tr>
<tr>
<td>all-to-one</td>
<td>general server, synchronization</td>
</tr>
<tr>
<td>many-to-many</td>
<td>general network- or bus-system</td>
</tr>
</tbody>
</table>
Synchronization

Message structure

• Machine dependent representations need to be taken care of in a distributed environment.
• Communication system is often outside the typed language environment.

Most communication systems are handling streams (packets) of a basic element type only.

Conversion routines for data-structures other then the basic element type are supplied ...
...
... manually (POSIX)
... semi-automatic (Real-time CORBA)
... automatic and are typed-persistent (Ada95)

Reading and writing values of any type to a stream:

procedure S'Write(
    Stream : access Ada.Streams.Root_Stream_Type'Class; Item : in  T);
procedure S'Class'Write(
    Stream : access Ada.Streams.Root_Stream_Type'Class; Item : in  T'Class);
procedure S'Read(
    Stream : access Ada.Streams.Root_Stream_Type'Class; Item : out T);
procedure S'Class'Read(
    Stream : access Ada.Streams.Root_Stream_Type'Class; Item : out T'Class)

Reading and writing values, bounds and discriminants of any type to a stream:

procedure S'Output(
    Stream : access Ada.Streams.Root_Stream_Type'Class; Item : in T);
function S'Input(
    Stream : access Ada.Streams.Root_Stream_Type'Class) return T;

Message-based synchronization

Practical message-passing systems:

POSIX:

- ordered indirect [asymmetrical | symmetrical] asynchronous byte-level many-to-many message passing

CHILL:

- ordered indirect [asymmetrical | symmetrical | synchronous | asynchronous] typed [many-to-many | many-to-one] message passing

Occam2:

- indirect symmetrical synchronous fully-typed one-to-one message passing

Ada95:

- (extended) rendezvous:
  - ordered direct asynchronous [synchronous | asynchronous] fully-typed many-to-one remote invocation

Real-time Java:

- no communication via messages available
Real-Time & Embedded Systems

Message-based synchronization

Practical message-passing systems:

<table>
<thead>
<tr>
<th>Method</th>
<th>Ordered</th>
<th>Symmetrical</th>
<th>Asynchronous</th>
<th>Direct</th>
<th>Indirect</th>
<th>Contents</th>
<th>One-to-one</th>
<th>Many-to-many</th>
</tr>
</thead>
<tbody>
<tr>
<td>POSIX</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td>bytes</td>
<td></td>
<td></td>
</tr>
<tr>
<td>CHILL</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td>typed</td>
<td></td>
<td></td>
</tr>
<tr>
<td>Occam2</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td>fully typed</td>
<td></td>
<td></td>
</tr>
<tr>
<td>Ada95</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td>fully typed</td>
<td></td>
<td></td>
</tr>
<tr>
<td>Real-time Java</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

Real-time Java: no communication via messages available

Message-based synchronization in Occam2

Communication is ensured by means of a ‘channel’, which:

- can be used by one writer and one reader process only
- and is synchronous:

CHAN OF INT SensorChannel:

PAR
  INT reading:
  SEQ i = 0 FOR 1000
  SEQ -- generate reading
  SensorChannel! reading
  INT data:
  SEQ i = 0 FOR 1000
  SEQ SensorChannel? data -- employ data

Real-time Java: no communication via messages available

Message-based synchronization in CHILL

CHILL is the ‘CCITT High Level Language’, where CCITT is the Comité Consultatif International Télégraphique et Téléphonique. The CHILL language development was started in 1973 and standardized in 1979.

- strong support for concurrency, synchronization, and communication (monitors, buffered message passing, synchronous channels)

dcl SensorBuffer buffer (32) int;
... send SensorBuffer (reading);
... receive case (SensorBuffer in data) : ...
esac;

signal SensorChannel = (int) to consumertype;
... send SensorChannel (reading) to consumer
... receive case (SensorChannel in data) : ...
esac;
Synchronization

Message-based synchronization in CHILL

CHILL is the 'CCITT High Level Language', where CCITT is the Comité Consultatif International Télégraphique et Téléphonique. The CHILL language development was started in 1973 and standardized in 1979.

- strong support for concurrency, synchronization, and communication (monitors, buffered message passing, synchronous channels)

```chill
dcl SensorBuffer buffer (32) int;
  ...
  send SensorBuffer (reading);
  ...
  receive case ...
    (SensorBuffer in data) : ...
  esac;

signal SensorChannel = (int) to consumer type;
  ...
  send SensorChannel (reading)
  ...
  receive case ...
    (SensorChannel in data): ...
  esac;
```

Message-based synchronization in Ada95

Ada95 supports remote invocations ((extended) rendezvous) in form of:

- entry points in tasks
- full set of parameter profiles supported

If the local and the remote task are on different architectures, or if an intermediate communication system is employed:

- parameters incl. bounds and discriminants are 'tunneled' through byte-stream-formats.

Synchronization:

- both tasks are synchronized at the beginning of the remote invocation ("rendezvous")
- the calling task if blocked until the remote routine is completed ("extended rendezvous")
Message-based synchronization in Ada95

(Rendezvous)

\[
\text{accept } \langle \text{entry\_name} \rangle \{\langle \text{index} \rangle \} \langle \text{parameters} \rangle \text{;}
\]

\[
\text{--- waiting for synchronization}
\]

\[
\begin{align*}
\text{synchronized} & \quad \text{accept } \langle \text{entry\_name} \rangle \{\langle \text{index} \rangle \} \langle \text{parameter\_profile} \rangle; \\
\text{--- waiting for synchronization}
\end{align*}
\]

Message-based synchronization in Ada95

(Extended rendezvous)

\[
\text{accept } \langle \text{entry\_name} \rangle \{\langle \text{index} \rangle \} \langle \text{parameters} \rangle \text{ do}
\]

\[
\begin{align*}
\text{--- blocked} & \quad \text{--- remote invocation} \\
\text{return results} & \quad \text{end } \langle \text{entry\_name} \rangle;
\end{align*}
\]

\[
\text{synchronized}
\]

\[
\text{synchronized}
\]

\[
\text{synchronized}
\]

\[
\text{return results}
\]

\[
\text{end } \langle \text{entry\_name} \rangle;
\]
Message-based synchronization in Ada95

Some things to consider for task-entries:
- In contrast to protected-object-entries, task-entries can call other blocking operations.
- Accept statements can be nested (but need to be different).
  ☞ helpful e.g. to synchronize more than two tasks.
- Accept statements can have a dedicated exception handler (like any other code-block).
  Exceptions, which are not handled during the rendezvous phase are propagated to all involved tasks.
- Parameters cannot be direct ‘access’ parameters, but can be access-types.
- ‘count on task-entries is defined, but is only accessible from inside the tasks owning the entry.
- Entry families (arrays of entries) are supported.
- Private entries (accessible for internal tasks) are supported.

Selective waiting in Occam2

ALT
  Guard1
    Process1
  Guard2
    Process2

- Guards are referring to boolean expressions and/or channel input operations.
- The boolean expressions are local expressions, i.e. if none of them evaluates to true at the time of the evaluation of the ALT-statement, then the process is stopped.
- If all triggered channel input operations evaluate to false, the process is suspended until further activity on one of the named channels.
- Any Occam2 process can be employed in the ALT-statement.
- The ALT-statement is non-deterministic (there is also a deterministic version: PRI ALT).

Selective waiting

Dijkstra’s guarded commands:

\[
\begin{aligned}
& \text{if } x \leq y \rightarrow m := x \\
& \text{fi}
\end{aligned}
\]

\[
\begin{aligned}
& \text{if } x \geq y \rightarrow m := y \\
& \text{fi}
\end{aligned}
\]

☞ the programmer needs to design the alternatives as ‘parallel’ options: all cases need to be covered and overlapping conditions need to lead to the same result

Extremely different philosophy: ‘C’-switch:

\[
\begin{aligned}
\text{switch } (x) \{ \\
& \text{case 1: } r := 3; \\
& \text{case 2: } r := 2; \text{ break; } \\
& \text{case 3: } r := 1;
\}
\end{aligned}
\]

☞ the sequence of alternatives has a crucial role.

CSP (Hoare) also supports non-deterministic selective waiting

\[
\begin{aligned}
& \text{ALT: NumberInBuffer } < \text{ Size } \& \text{ Append } ? \text{ Buffer [Top] } \\
& \text{SEQ: NumberInBuffer } := \text{ NumberInBuffer } + 1 \\
& \text{Top } := (\text{Top } + 1) \text{ REM Size } \\
& \text{NumberInBuffer } > 0 \& \text{ Request } ? \text{ ANY } \\
& \text{SEQ: Take } ! \text{ Buffer [Base] } \\
& \text{NumberInBuffer } := \text{ NumberInBuffer } - 1 \\
& \text{Base } := (\text{Base } + 1) \text{ REM Size }
\end{aligned}
\]

- synchronization on input-channels only:
  ☞ to initiate the sending of data (Take ! Buffer [Base]), a request need to be made first (Request ? ANY)
Message-based selective synchronization in Ada95

Forms of selective waiting:

- `select_statement ::= selective_accept | conditional_entry_call | timed_entry_call | asynchronous_select`

... underlying concept: Dijkstra's guarded commands

`selective_accept` implements...

- ... wait for more than a single rendezvous at any one time
- ... time-out if no rendezvous is forthcoming within a specified time
- ... withdraw its offer to communicate if no rendezvous is available immediately
- ... terminate if no clients can possibly call its entries

Basic forms of selective synchronization

(select-or)

```
select
  accept ... do ...
end ...
or
  accept ... do ...
end ...
or
  accept ... do ...
end ...
or
  accept ... do ...
end ...
end select;
```

- If none of the named entries have been called, the task is suspended until one of the entries is addressed by another task.
- The selection of an accept is non-deterministic, in case that multiple entries are called.
- The selection can be controlled by means of the real-time systems annex.
- The select statement is completed, when at least one of the entries has been called and those accept-block has been executed.

(select-guarded)

```
select
  when <condition> =>
    accept ... do ...
end ...
or
  when <condition> =>
    accept ... do ...
end ...
or
  when <condition> =>
    accept ... do ...
end ...
or
  when <condition> =>
    accept ... do ...
end ...
end select;
```

• Analogue to Dijkstra's guarded commands
• all accepts closed will raise a Program_Error

☞ set of conditions need to be complete
**Synchronization**

**Basic forms of selective synchronization**
*(guarded select-or-else)*

```
select
  [ when <condition> => ]
    accept ... do ...
  end ...
or
  [ when <condition> => ]
    accept ... do ...
  end ...
or
  [ when <condition> => ]
    accept ... do ...
  end ...
else
  <statements>
end select;
```

- If none of the open entries can be accepted immediately, the else alternative is selected.
- There can be only one else alternative and it cannot be guarded.

**Basic forms of selective synchronization**
*(guarded select-or-delay)*

```
select
  [ when <condition> => ]
    accept ... do ...
  end ...
or
  [ when <condition> => ]
    delay ...
  <statements>
  or
  [ when <condition> => ]
    delay ...
  <statements>
  else
  <statements>
end select;
```

- If none of the open entries has been called before the amount of time specified in the earliest open delay alternative, this delay alternative is selected.
- There can be multiple delay alternatives if more than one delay alternative expires simultaneously, either one may be chosen.
- `delay` and `delay until` can be employed.

**Basic forms of selective synchronization**
*(guarded select-or-terminate)*

```
select
  [ when <condition> => ]
    accept ... do ...
  end ...
end select;
```

- The terminate alternative is chosen if none of the entries can ever be called again, i.e.:
  - all tasks which can possibly call any of the named entries are terminated.
  - all remaining active tasks which can possibly call any of the named entries are waiting on selective terminate statements and none of their open entries can be called any longer.
  - This task and all its dependent waiting-for-termination tasks are terminated together.

```
select
  [ when <condition> => ]
    delay-terminate
    alternatives
cannot be mixed!
  end ...
```

- `else-delay-terminate` alternatives cannot be mixed!
Non-determinism in selective synchronizations

- If equal alternatives are given, then the program correctness (incl. the timing specifications) must not be affected by the actual selection.
- If alternatives have different priorities, this can be expressed e.g. by means of the Ada real-time annex.
- Non-determinism in concurrent systems is or can be also introduced by:
  - non-ordered monitor or other queues
  - buffering / routing message passing systems
  - non-deterministic schedulers
  - timer quantization
  - ... any form of asynchronism

Conditional & timed entry-calls

conditional_entry_call ::= select
  entry_call_statement
  [sequence_of_statements]
  else
  sequence_of_statements
  end select;

select
  Light_Monitor.Wait_for_Light;
  Lux := True;
else
  Lux := False;
end;

timed_entry_call ::= select
  entry_call_statement
  [sequence_of_statements]
  or
delay_alternative
  end select;

select
  Controller.Request (Medium)
  (Some_Item);
  -- process data
  or
delay 45.0;
  -- try something else
  end select;

There is only one entry call and either one 'else' or one 'or delay'
...
Summary

Synchronization

• Shared memory based synchronization
  • Flags, condition variables, semaphores, ...
  • Guard evaluation times, nested monitor calls, deadlocks, ...
  • ... simultaneous reading, queue management.
  • Synchronization and object orientation, blocking operations and re-queuing.

• Message based synchronization
  • Synchronization models, addressing modes, message structures
  • Selective accepts, selective calls
  • Indeterminism in message based synchronization

References for this chapter

[Burns01] Alan Burns and Andy Wellings
  Real-Time Systems and Programming Languages
  Addison Wesley, third edition, 2001

[Murthy01] C. Siva Ram Murthy, G. Manimaran
  Resource Management in Real-time Systems and Networks

all references and links are available on the course page

Scheduling in Real-Time Systems

☞ Concurrency may lead to non-determinism
☞ Non-determinism may make it harder to predict the timing behaviour

☞ RT-Scheduling schemes reduce non-determinism
Scheduling

A scheduling scheme provides two features:

- Ordering the use of resources (e.g. CPUs, networks)
- Predicting the worst-case behaviour of the system when the scheduling algorithm is applied

The prediction can then be used
- at compile-run: to confirm the overall temporal requirements of the application
- or
- at run-time: to permit acceptance of additional usage/reservation requests.

Scheduling schemes

- **Static**
  - all predictions and schedules are done off-line
  - often better predictability = most hard real-time systems

- **Dynamic**
  - run-time situation is taken into account
  - more flexible, more efficient = most soft real-time systems

Real-time scheduling as task queuing

pre-emption or cycle done

create → batch → ready → CPU → term.

ready, suspended
blocked, suspended
blocked

block or synchronize

admitting new tasks if set is still schedulable

dispatching according to deadlines, priorities, or values
A simple process model

- The number of processes in the system is fixed.
- All processes are periodic and all periods are known.
- All processes are independent.
- The task-switching overhead is negligible.
- All deadlines are identical with the process cycle times (periods).
- The worst case execution time is known for all processes.
- All processes are released at once.

This model can only be applied to a specific group of hard real-time systems. (extensions to this model will be discussed later in this chapter).

Example: Requested times

Example: Deadlines

Earliest deadline first (EDF)

1. Determine (one of) the process(es) with the closest deadline.
2. Execute this process
   2-a until it finishes
   2-b or until another process' deadline is found closer than the current one.

= Pre-emptive scheme
= Dynamic scheme, since the dispatched process is selected at run-time, due to the current deadlines.
1. Schedule the earliest deadline first

2. Avoid task switches (in case of equal deadlines)

worst case response times $R_i$ (maximal time in which the request from task $T_i$ is served).

$R_i$ can be close or identical to deadlines.

small or none spare capacity, if any task misses its expected computation time.
Fixed Priority Scheduling (FPS), rate monotonic

1. Each process is assigned a fixed priority according to its cycle time $T_i$:
   \[ T_i < T_j \Rightarrow P_i > P_j \]

2. At any point in time: dispatch the process with the highest priority

- Pre-emptive scheme
- Static scheme, since the order dispatch order of processes is fixed and calculated at off-line.

Rate monotonic priorities:

- Rate monotonic ordering is optimal (in the framework of fixed priority schedulers), i.e. if a process set is schedulable under a FPS-scheme, it is also schedulable by applying rate monotonic priorities.

\begin{align*}
(16,8) & \quad (12,3) & \quad (4,1) \\
(T_i, C_i) & \quad (T_i, C_i) & \quad (T_i, C_i)
\end{align*}

\begin{align*}
1 & \quad 2 & \quad 3 \\
5 & \quad 10 & \quad 15 \\
15 & \quad 20 & \quad 25 \\
25 & \quad 30 & \quad 35 \\
35 & \quad 40 & \quad 45 \\
45 & \quad 50 & \quad 1
\end{align*}
Static scheduling: Fixed Priority Scheduling (FPS), rate monotonic

Rate monotonic priorities

Rate monotonic priorities (reduced requests)

max. utilization test: \[ \sum_{i=1}^{n} \frac{C_i}{T_i} \leq N \left( \frac{1}{2^N} - 1 \right) \]  
⇒ sufficient, but not necessary test!

utilization test: \[ \sum_{i=1}^{n} \frac{C_i}{T_i} = 1 > 0.779 \approx \left( \frac{1}{N} \right) \left( \frac{1}{2^N} - 1 \right) \]  
⇒ not guaranteed!

max. utilization test: \[ \sum_{i=1}^{n} \frac{C_i}{T_i} \leq N \left( \frac{1}{2^N} - 1 \right) \]  
⇒ utilization: \[ \frac{6}{16} + \frac{3}{12} + \frac{1}{4} = 0.875 > 0.779 \approx \left( \frac{1}{N} \right) \left( \frac{1}{2^N} - 1 \right) \]  
⇒ not guaranteed!
Real-Time & Embedded Systems

Static scheduling: Fixed Priority Scheduling (FPS), rate monotonic

Rate monotonic priorities (reduced requests)

\[
\begin{align*}
(16,6), (12,3), (4,1) \\
(T_i, C_i)
\end{align*}
\]

\[
\text{utilization: } \frac{6}{16} + \frac{5}{12} + \frac{1}{4} = 0.875 > 0.779 = 3 \left( \frac{1}{2} - 1 \right) \quad \text{\(\Rightarrow\) not guaranteed!}
\]

Rate monotonic priorities (further reduced requests)

\[
\begin{align*}
(16,4), (12,3), (4,1) \\
(T_i, C_i)
\end{align*}
\]

\[
\text{utilization: } \frac{4}{16} + \frac{3}{12} + \frac{1}{4} = 0.75 \leq 0.779 = 3 \left( \frac{1}{2} - 1 \right) \quad \text{\(\Rightarrow\) guaranteed!}
\]
Static scheduling: Fixed Priority Scheduling (FPS), rate monotonic

**Utilisation based Analysis for FPS rate monotonic**

\[ U ≔ \sum_{i=1}^{n} \frac{C_i}{T_i} \leq N \left( 2^{-N} - 1 \right) \equiv U_{\text{max}} \]

- with \( C_i \) the computation time and \( T_i \) the length of the period for task \( i \) out of \( N \) tasks and assuming that the deadline \( D_j = T_j \)

⇒ sufficient, but not necessary
⇒ \( O(n) \) complexity

**Response time analysis (further reduced requests)**

- calculate the worst case response times for each task individually.

⇒ for the highest priority task: \( R_3 = C_3 \)

⇒ for other tasks: \( R_i = C_i + I_i = \text{computation} C_i + \text{interference} I_i \)
Response time analysis

\[ R^k_{i+1} = C_i + \sum_{j=i}^{k} \left( \frac{R_j}{T_j} \right) C_j \]

⇔ fixed-point equation!

\[ R^0_i = C_i \]

⇔ starting with \( R^0_i = C_i \)

⇔ iterate (1) until \( R^{k+1}_i = R^k_i \) or \( R^{k+1}_i > T_i \)

Example (further reduced requests):

• set of tasks: \( \{(T_p, C_j)\} = \{(16, 4); (12, 3); (4, 1)\} \) at priorities \( \{1; 2; 3\} \); \( R^0_3 = 1 \)

\[ R^0_3 = 1(\checkmark) \]

• set of tasks: \( \{(T_p, C_j)\} = \{(16, 4); (12, 3); (4, 1)\} \) at priorities \( \{1; 2; 3\} \); \( R^0_2 = 3 \)

\[ R^1_2 = 3 + 4 = 7 \]

\[ R^2_2 = 3 + 4 = 7(\checkmark) \]
Real-Time & Embedded Systems

Static scheduling: Fixed Priority Scheduling (FPS), rate monotonic

Response time analysis

\[ R_j^{k+1} = C_j + \sum_{i \geq j} \left[ \frac{R_i^k}{T_j} \right] C_j \] (1); \( R_j^0 = C_j - \text{iterate (1)} \) until \( R_j^{k+1} = R_j^k \) or \( R_j^{k+1} > T_j \)

Example (full requests):
- set of tasks: \( \{(T_p, C_p)\} = \{(16, 8); (12, 3); (4, 1)\} \) at priorities \{1; 2; 3\}; \( R_1^0 = 8 \)
  - \( R_1^1 = 8 + \left[ \frac{8}{12} \right] 3 + \left[ \frac{8}{4} \right] 1 = 13 \)
  - \( R_1^2 = 8 + \left[ \frac{13}{12} \right] 3 + \left[ \frac{13}{4} \right] 1 = 18(\times) \)
  - \( R_1^3 = 8 + \left[ \frac{18}{12} \right] 3 + \left[ \frac{18}{4} \right] 1 = 19(\times) \)

Response time analysis

\[ R_j^{k+1} = C_j + \sum_{i \geq j} \left[ \frac{R_i^k}{T_j} \right] C_j \] (1); \( R_j^0 = C_j - \text{iterate (1)} \) until \( R_j^{k+1} = R_j^k \) or \( R_j^{k+1} > T_j \)

Example (reduced requests):
- set of tasks: \( \{(T_p, C_p)\} = \{(16, 6); (12, 3); (4, 1)\} \) at priorities \{1; 2; 3\}; \( R_1^0 = 6 \)
  - \( R_1^1 = 6 + \left[ \frac{6}{12} \right] 3 + \left[ \frac{6}{4} \right] 1 = 11 \)
  - \( R_1^2 = 6 + \left[ \frac{11}{12} \right] 3 + \left[ \frac{11}{4} \right] 1 = 12 \)
  - \( R_1^3 = 6 + \left[ \frac{12}{12} \right] 3 + \left[ \frac{12}{4} \right] 1 = 12(\checkmark) \)

Dynamic scheduling: Earliest Deadline First (EDF)

Response time analysis

The worst case for EDF is not necessarily when all tasks are released at once!
- all possible combinations in a full hyper-cycle need to be considered!
  - The response times are bounded by the cycle times as long as the maximal utilization is \( \leq 1 \).
  - Other tasks need to be considered only, if their deadline is closer or equal to the current task.
Dynamic scheduling: Earliest Deadline First (EDF)

Response time analysis

\[ R_j(a) = \left[ \frac{a}{T_i} + 1 \right] C_i + \sum_{j \neq i \min} \left\{ \left[ \frac{R_j(a)}{T_j} \right] \max \left\{ 0, \frac{a + T_i - T_j}{T_j} + 1 \right\} \right\} C_j \]

\[ R_j^{k+1}(a) = \left[ \frac{a}{T_i} + 1 \right] C_i + \sum_{j \neq i \min} \left\{ \left[ R_j^k(a) \right] \max \left\{ 0, \frac{a + T_i - T_j}{T_j} + 1 \right\} \right\} C_j \]

⇒ starting with \( R_j^0(a) = a + C_i \)

⇒ Iterate (2) until \( R_j^{k+1}(a) = R_j^k(a) \) (or \( R_j^{k+1}(a) > a + T_i \Rightarrow utilization\ beyond\ 100\%!) \n
⇒ \( R_j = \max \{ R_j(a) - a \}_{a \in A} \), where \( A = \text{scm} \{ T_i \} \)

Static scheduling: Fixed Priority Scheduling (FPS), rate monotonic

Response time analysis (further reduced requests)

\[ R_j = C_i + \sum_{j > i} \left[ \frac{R_j(a)}{T_j} \right] C_j; R_3 = 1 \nu; R_2 = 4 \nu; R_1 = 10 \nu \text{ and } \sum_{i=1}^{N} \frac{C_i}{T_i} \leq N \left( \frac{1}{2^n - 1} \right) \nu \]
Real-Time & Embedded Systems

Static scheduling: Fixed Priority Scheduling (FPS), rate monotonic

Response time analysis (reduced requests)

\[ R_i = C_i + \sum_{j > i} \left \lceil \frac{R_j}{T_i} \right \rceil C_j; R_3 = 1 \nu; R_2 = 4 \nu; R_1 = 12 \nu \text{ but } \sum_{i=1}^{n} \frac{C_i}{T_i} > N\left(\frac{1}{2^N - 1}\right) \times \]

Dynamic scheduling: Earliest Deadline First (EDF)

Response time analysis (full requests)

\[ R_i = C_i + \sum_{j > i} \left \lceil \frac{R_j}{T_i} \right \rceil C_j; R_3 = 1 \nu; R_2 = 4 \nu; R_1 = 19 \nu \text{ and } \sum_{i=1}^{n} \frac{C_i}{T_i} > N\left(\frac{1}{2^N - 1}\right) \times \]

© 2009 Uwe R. Zimmer, The Australian National University

Page 593 of 769 (chapter 7: to 667)
Dynamic scheduling: Earliest Deadline First (EDF)

Response time analysis (reduced requests)

- Relaxed task-set changes:
  \[ R_-: 16 \rightarrow 12 \leq 16 \nu = T_-; \quad R_-: 12 \rightarrow 8 \leq 12 \nu = T_-; \quad R_-: 4 \rightarrow 1 \leq 4 \nu = T_- \]

Further relaxed task-set changes:

- Further relaxed task-set changes:
  \[ R_-: 12 \rightarrow 10 \leq 16 \nu = T_-; \quad R_-: 8 \rightarrow 6 \leq 12 \nu = T_-; \quad R_-: 1 \rightarrow 1 \leq 4 \nu = T_- \]

Real-time scheduling

Response time analysis (comparison)

<table>
<thead>
<tr>
<th>Task Set</th>
<th>Fixed Priority Scheduling</th>
<th>Earliest Deadline First</th>
</tr>
</thead>
<tbody>
<tr>
<td>((T_P, C_i)) = {(16, 8); (12, 3); (4, 1)}</td>
<td>✗ (1.000)</td>
<td>✓ (1.000)</td>
</tr>
<tr>
<td>(= {16, 12, 4})</td>
<td>✓ (1.000)</td>
<td>(= {16, 12, 4})</td>
</tr>
<tr>
<td>((T_P, C_i)) = {(16, 6); (12, 3); (4, 1)}</td>
<td>✗ (0.875)</td>
<td>✓ (0.875)</td>
</tr>
<tr>
<td>(= {12, 8, 1})</td>
<td>✓ (0.875)</td>
<td>(= {12, 8, 1})</td>
</tr>
<tr>
<td>((T_P, C_i)) = {(16, 4); (12, 3); (4, 1)}</td>
<td>✓ (0.750)</td>
<td>✓ (0.750)</td>
</tr>
<tr>
<td>(= {10, 6, 1})</td>
<td>✓ (0.750)</td>
<td>(= {10, 6, 1})</td>
</tr>
</tbody>
</table>

- EDF can handle higher (full) utilization than FPS.
- FPS is easier to implement and implies less run-time overhead.
- Graceful degradation features (resource is over-booked):
  - FPS: processes with lower priorities will always miss their deadlines first.
  - EDF: any process can miss its deadline \(\nu\) and can trigger a cascade of failed deadlines.
- Response time analysis and utilization tests:
  - FPS: \(O(n)\) utilization test — response time analysis: fixed point equation
  - EDF: \(O(n)\) utilization test — response time analysis: fixed point equation in hyper-cycle

Fixed Priority Scheduling \(\leftrightarrow\) Earliest Deadline First

- EDF can handle higher (full) utilization than FPS.
- FPS is easier to implement and implies less run-time overhead.
- Graceful degradation features (resource is over-booked):
  - FPS: processes with lower priorities will always miss their deadlines first.
  - EDF: any process can miss its deadline \(\nu\) and can trigger a cascade of failed deadlines.
- Response time analysis and utilization tests:
  - FPS: \(O(n)\) utilization test — response time analysis: fixed point equation
  - EDF: \(O(n)\) utilization test — response time analysis: fixed point equation in hyper-cycle
Constraints which we used up to here:

- tasks are periodic
- deadlines are identical with task’s period time \((D = T)\)
- pre-emptive scheduling
- worst case execution times are known
- tasks are independent

Extensions which we will introduce:

- tasks are periodic
  - we will introduce sporadic and aperiodic processes
- deadlines are identical with task’s period time \((D = T)\)
  - we will introduce arbitrary deadlines
- pre-emptive scheduling
  - we will introduce (briefly) cooperative scheduling
- worst case execution times are known
  - we will introduce fault tolerant scheduling
- tasks are independent
  - we will introduce schedules for interacting tasks

Including
aperiodic, sporadic & ‘soft’ real-time tasks
Introducing soft real-time tasks

Static scheduling: Fixed Priority Scheduling (FPS), rate monotonic

- set can be scheduled using average computation and period times
- hard real-time tasks can be scheduled under worst case conditions (including worst case behaviours of soft real-time tasks)

Introducing a server task

Server is established at a high priority

Deferrable Server (DS): Capacity replenished every $T_s$ (here: 8)
Real-Time & Embedded Systems

Static scheduling: FPS, rate monotonic + server

Introducing a server task: Sporadic Server

Sporadic Server (SS): Capacity replenished $T_s$ units after $t_s \approx$ POSIX

Real-Time & Embedded Systems

Static scheduling: Fixed Priority Scheduling (FPS), dual-priorities

Introducing dual priorities

start hard rt-tasks in low priorities; promote them in time to higher ones

Real-Time & Embedded Systems

Dynamic scheduling: Earliest Deadline First + aperiodic server

Introducing a server task to EDF

Real-Time & Embedded Systems

Dynamic scheduling: Earliest Deadline First + aperiodic server

Introducing a server task to EDF
Dynamic scheduling: Earliest Deadline First + aperiodic tasks

Switching between EDF & Earliest Deadline Last (EDL)

Scheduling — real-world considerations

Including

tasks with deadlines shorter than their cycle time

Tasks with D < T
(Deadline earlier than inter-arrival period)

In fixed priority scheduling (FPS): change from:
Rate Monotonic Priority Ordering (RMPO)
to:
Deadline Monotonic Priority Ordering (DMPO)

Lemma

Any task set Q which is schedulable by a FPS scheme W
is also schedulable by DMPO!

Proof

1. i, j are two tasks in Q, with \( (P_i = P_j + 1) \wedge (D_i > D_j) \) in W \( \Rightarrow \) DMPO
2. Generate W' by swapping \( P_i \) and \( P'_j \), i.e. \( (P'_j < P_j) \wedge (D_j > D_j) \) \( \Rightarrow \) DMPO
3. W' is scheduling Q because:
   3-a all \( t_k \in Q \) with \( P_k > P_i \) or \( P_k < P_j \) are unaffected
   3-b \( t_j \) is schedulable in W' because \( P'_j > P_j \Rightarrow R'_j \leq R_j \leq D_j 
   3-c \( t_i \) is schedulable in W' because …
**Real-Time & Embedded Systems**

**Proof**

1. \( t_i, t_j \) are two tasks in \( Q \), with \( P_i > P_j \) and \( D_i > D_j \) in \( W \Rightarrow \neg \text{DMPO} \)

2. Generate \( W' \) by swapping \( P_i \) and \( P_j \), i.e. \((P_i' < P_j') \land (D_i > D_j) \Rightarrow \text{DMPO} \)

3. \( W' \) is scheduling \( Q \) because:
   - 3-a all \( t_k \in Q \) with \( P_k > P_i \) or \( P_k < P_j \) are unaffected
   - 3-b \( t_j \) is schedulable in \( W' \) because \( P_j' > P_j \Rightarrow R_j' \leq R_j \leq D_j \)
   - 3-c \( t_i \) is schedulable in \( W' \) because
     - in \( W' \): \( R_j \leq D_j < D_i \leq T_i \Rightarrow R_j < T_i \) i.e. \( t_i \) interfered only once with \( t_j \)
     - in \( W' \): \( t_j \) released once in \( R_j \) and \( R_i < R_j \)
     - in \( W' \): \( t_j \) interferes only once with \( t_i \Rightarrow R_i' = R_j \leq D_j < D_i \Rightarrow R_i' < D_i \)

**Static scheduling: Fixed Priority Scheduling (FPS), DMPO**

Any task set \( Q \) which is schedulable by a FPS scheme \( W \), is also schedulable by DMPO.

- Swap all \( t_i, t_j \) in \( Q \), with \( P_i > P_j \) and \( D_i > D_j \) in \( W \) resulting in all \( t_i, t_j \) in \( Q \), with \( P_i > P_j \) have \( D_i < D_j \)
  \( \Rightarrow \) **Deadline Monotonic Priority Ordering (DMPO)**

- Since the each swapping operation keeps schedulability, the final priority scheme (DMPO) is also schedulable.

\( \Rightarrow \) If FPS-DMPO is not schedulable, there is no schedulable FPS-scheme.
Scheduling — real-world considerations

Including

task interdependencies

Schedule for independent tasks

Synchronized via lock

(independent task set)

(independent task set)

(interdependent task set ⇒ lock [shared between ] and )
**Scheduling: Interdependencies**

**Synchronized via lock**

Priority inversion

(Interdependent task set \( \text{lock} \) shared between \( \text{and} \) )

**Priority inheritance**

Task \( t_i \) inherits the priority of \( t_j \) if:

1. \( P_i < P_j \)
2. Task \( t_i \) has locked a resource \( Q \)
3. Task \( t_j \) is blocked waiting for resource \( Q \) to be released

Maximal blocking time for task \( t_i \): \( B_i = \sum_{r=1}^{R} usage(r, i)C(r) \)

- with \( R \) the number of critical sections
- \( usage(r, i) \) a boolean \( (0/1) \) function indicating that \( r \) is used by at least one \( t_j \) with \( P_j < P_i \) and at least one \( t_k \) with \( P_k \geq P_i \)
- \( C(r) \) is the worst case computation time in critical section \( r \)

A task can only be blocked once for each employed resource!
Without priority inheritance

Priority inversion
(interdependent task set lock shared between  and )

A more complex example

(independent task set)

A more complex example

(independent task set)
Scheduling: Interdependencies

**Priority inheritance**

( and inherit priority of , when in lock and is dispatched)

**Priority inversion**

( and inherit priority of , when in lock and is dispatched)

One additional lock request

( and inherit priority of , when in lock and is dispatched)
Scheduling: Interdependencies

One additional lock request

⇒ Deadlock

Scheduling: Interdependencies: Priority ceiling protocols

Immediate ceiling priority protocol (POSIX, Ada, RT-Java)

- Each task $t_i$ has static default priority $P_i$.
- Each resource (lock, monitor) $R_k$ has a static ceiling priority $C_k$, which is the maximum of priorities of the tasks $t_i$ which employ this resource.

$$C_k = \max \{ employ(i, k) \cdot P_i \}$$

- Each task $t_i$ has a dynamic priority $P_i^D$, which is the maximum of its own static priority and the ceiling priorities of any resource it has locked.

$$P_i^D = \max \{ P_i, \max_k \{ locked(i, k) \cdot C_k \} \}$$

(, , and inherit the ceiling priority of or when entering the lock)
Real-Time & Embedded Systems

Scheduling: Interdependencies: Priority ceiling protocols

Immediate ceiling priority protocol (POSIX, Ada, RT-Java)

- Tasks are dispatched only if all employed resources are available.
- Deadlocks are prevented
- Number of context switches is reduced

Maximal blocking time:

- with $R$ the number of critical sections
- $\text{usage}(r, i)$ a boolean (0/1) function indicating that $r$ is used by at least one $t_j$ with $P_j < P_i$ and at least one $t_k$ with $P_k \geq P_i$
- $C(r)$ is the worst case computation time in critical section $r$

A task can only be blocked once by any lower priority task!

Scheduling — real-world considerations

Considering non-pre-emptive scheduling

Cooperative Scheduling

- All schemes up to here used pre-emptive dispatching.
- In interdependent task sets maximal blocking times $B_i$ can be determined for each task $t_i$ when employing a priority ceiling protocol.
- If the overall maximal blocking time $B_{\text{max}}$ can be accepted by all tasks, the number of pre-emption can significantly reduced by:

**Deferred pre-emption – Cooperative Scheduling**

- Each task $t_i$ is divided in $k$ non-pre-emptive blocks of $C_{ik} \leq B_{\text{max}}$
- All critical sections are completely enclosed in one code-block
- Each task calls a ‘de-scheduling’ kernel routine at the end of each code-block, i.e., ‘offering’ a task-switch.
Cooperative Scheduling

Deferred pre-emption – Cooperative Scheduling

Number of task switches is reduced
Caches, pre-fetching, and pipelines are more efficient
Execution times are (a bit) easier to predict
Schedules are simpler
Interdependent task sets are schedulable deadlock free

Deferred pre-emption – Cooperative Scheduling

Response times:

\[ R_i = R_i^n + F_i, \text{ with } R_i^{k+1} = B_{\text{max}} + C_j - F_i + \sum_{j > i} \left[ \frac{R_i^k}{T_j} \right] C_j \]

and \( F_i \) the execution time of the final code-block

... in the simplified case \( C = C_i = C_j = F_i = B_{\text{max}} \):

\[ R_i = R_i^n, \text{ with } R_i^{k+1} = C + \sum_{j > i} \left[ \frac{R_i^k}{T_j} \right] C \]

... and with even \( T = T_i \forall i \):

\[ R_i = C + \sum_{j > i} C \]

Cooperative Scheduling

Deferred pre-emption – Cooperative Scheduling

What’s the cost?

- Code block division need to be done very thoroughly.
- Additional protection against badly behaving (non-cooperative) tasks:
  - Scheduler pre-empts tasks, which fail to offer a ‘de-schedule’ themselves.
- Due to a central \( B_{\text{max}} \) value, additional tasks need to be engineered to participate in a specific cooperative schedule.
- Requires that a \( B_{\text{max}} \) value can be accepted by all tasks
  - very short and reactive tasks are excluded or will be treated specially.

Scheduling — real-world considerations

Considering

deadlines beyond the release period
Tasks with $D > T$
(Deadline later than inter-arrival period)
(a cross-over of a hard, periodic and a soft real-time task)

Assuming that a task $t_j$ is released only after a former release of $t_j$ is completed.

☞ In case that $R > T$ for a specific scheduling situation, the following release of task $t_j$ is delayed until completion of the former release.

☞ Mind that $R > T$ cannot hold for all release situations, otherwise the task is not schedulable.

☞ The worst case response time $R_i$ might thus be longer than $T_i$ but must still be shorter than $D_i$.

Since the response time $R$ can now be potentially greater than the cycle time $T$:

more than one release $q$ of the task $t_i$ needs to be considered:

$$R_i(q) = B_i + q C_i + \sum_{j > i}^{T_j} \left[ \frac{R_j(q)}{T_j} \right] C_j$$

where $\forall q | (R_i(q) - (q - 1) T_i \leq D_i)!$

with $q$: number of releases

$$R_i = \max \{ R_i(q) - (q - 1) T_i | q \in \{1, \ldots q_{\max}\} \}$$

and $q_{\max} = \min \left\{ q \left| \frac{R_i(q)}{q} \leq T_i \right. \right\}$

Scheduling — real-world considerations

Considering
‘fault-tolerance’
(additional CPU-time for exception handling and recovery)

Task $t_i$ needs extra CPU-time $C_i^f$ for error recovery or exception handling (done at $P_i$) and the minimum inter-arrival time between faults is $T_f$:

$$R_i = B_i + C_i + \sum_{j > i}^{T_j} \left[ \frac{R_j}{T_j} \right] C_j + \max_{k \geq i} \left\{ \left[ \frac{R_k}{T_f} \right] C_k^f \right\}$$

if error recovery or exception handling is performed at the highest priority:

$$R_i = B_i + C_i + \sum_{j > i}^{T_j} \left[ \frac{R_j}{T_f} \right] C_j + \max_{k \geq i} \left\{ \left[ \frac{R_k}{T_f} \right] C_k^f \right\}$$

Fault Tolerance

Exceptions and Recoveries

Task $t_i$ needs extra CPU-time $C_i^f$ for error recovery or exception handling (done at $P_i$) and the minimum inter-arrival time between faults is $T_f$:

$$R_i = B_i + C_i + \sum_{j > i}^{T_j} \left[ \frac{R_j}{T_j} \right] C_j + \max_{k \geq i} \left\{ \left[ \frac{R_k}{T_f} \right] C_k^f \right\}$$

if error recovery or exception handling is performed at the highest priority:

$$R_i = B_i + C_i + \sum_{j > i}^{T_j} \left[ \frac{R_j}{T_f} \right] C_j + \max_{k \geq i} \left\{ \left[ \frac{R_k}{T_f} \right] C_k^f \right\}$$
Scheduling — real-world considerations

Considering
task sets with release offsets

Some task sets can be scheduled by introducing offsets to the release times, but ...

☞ without any further restrictions this problem is NP-hard!

…

by introducing further assumptions about the granularity of the period times and deadlines:

☞ the schedulability analysis’ complexity can be reduced to be realistic.

Scheduling support

in different real-time languages / environments

Language Support for Scheduling : Ada95

Ada95 provides:

• Task and interrupt priorities (static, dynamic, active)
• Task attributes
• Prioritized entry queues
• Priority ceiling locking (ICPP)
• Schedulers (at least FIFO within priorities (pre-emptive) is requested)

Ada95 does not provide:

~ Earliest Deadline First (EDF)
~ Sporadic servers (a Ada95-implementation of a sporadic server is on the course page)
~ Direct task execution time measurements ~ e.g. POSIX or VxWorks timers
package System is
    subtype Any_Priority is Integer
        range implementation-defined;
    subtype Priority is Any_Priority
        range Any_Priority'First .. Any_Priority'Last;
    subtype Interrupt_Priority is Any_Priority
        range Priority'Last+1 .. Any_Priority'Last;
    Default_Priority : constant Priority :=
        (Priority'First + Priority'Last)/2;
end System;

package Ada.Dynamic_Priorities is
    procedure Set_Priority
        (Priority : in System.Any_Priority;
    function Get_Priority
        (T        : Ada.Task_Identification.Task_ID
        return System.Any_Priority;
end Ada.Dynamic_Priorities;

POSIX provides:
- Task and interrupt priorities (static, dynamic, active)
- Prioritized message queues
- Priority ceiling locking (ICPP)
- Schedulers, priority based with at least:
  - FIFO, Round-Robin, Sporadic Server, possibly others
- Threads can be
  - 'system contented' or
  - 'process contented' (priority scheduling unclear in this case)
- Timers

Real-Time Java provides:
- Task priorities (static, dynamic, active)
- Prioritized message queues
- Priority ceiling locking (ICPP)
- Schedulable objects (associated with threads) with
  - memory, release, and scheduling parameters
- Pre-emptive priority-oriented dispatching, possibly with a feasibility analysis
- An extendible scheduler class = dynamic scheduling

Real-Time Java does not (necessarily) provide:
- Earliest Deadline First (EDF)
- Sporadic servers
- Direct task execution time measurements (might be provided)
public class PriorityScheduler extends Scheduler
{
    public static final int MAX_PRIORITY;
    public static final int MIN_PRIORITY;
    protected PriorityScheduler();
    protected boolean addToFeasibility(Schedulable s);
    public void fireSchedulable(Schedulable s);
    public boolean isFeasible();
    protected boolean removeFromFeasibility(Schedulable s);
    public boolean setIfFeasible(Schedulable s, ReleaseParameters r, MemoryParameters m);
}

This PriorityScheduler is the only requested instantiation.

public abstract class SchedulingParameters
{
    public SchedulingParameters();
}

public class PriorityParameters extends SchedulingParameters
{
    public PriorityParameters(int priority);
    public int getPriority();
    public void setPriority(int priority) throws ...;
}

'Priority' is the only default scheduling parameter.

public abstract class ReleaseParameters
{
    protected ReleaseParameters(RelativeTime cost,
                                 RelativeTime deadline,
                                 AsyncEventHandler overrunHandler,
                                 AsyncEventHandler missHandler);
    public RelativeTime getCost();
    public AsyncEventHandler getCostOverrunHandler();
    public RelativeTime getDeadline();
    public AsyncEventHandler getDeadlineMissHandler();
}

Cost is an estimate of the max. execution time.

Measuring execution time is not requested, i.e. the overrunHandler might never be activated.

public class PeriodicParameters extends ReleaseParameters
{
    public PeriodicParameters(HighResolutionTime start,
                               RelativeTime period,
                               RelativeTime cost,
                               RelativeTime deadline,
                               AsyncEventHandler overrunHandler,
                               AsyncEventHandler missHandler);
    public RelativeTime getPeriod();
    public HighResolutionTime getStart();
    public void setPeriod(RelativeTime period);
    public void setStart(HighResolutionTime start);
}

most frequently used release parameters.
public class AperiodicParameters extends ReleaseParameters
{
    public AperiodicParameters (RelativeTime cost, RelativeTime deadline, AsyncEventHandler overrunHandler, AsyncEventHandler missHandler);
}

these are the minimum release parameters (while cost might be used for feasibility analysis only)

the deadline-missHandler need to be supplied in any implementation

public class SporadicParameters extends AperiodicParameters
{
    public SporadicParameters (RelativeTime minInterarrival, RelativeTime cost, RelativeTime deadline, AsyncEventHandler overrunHandler, AsyncEventHandler missHandler);
    public RelativeTime getMinimumInterarrival ();
    public void         setMinimumInterarrival (RelativeTime minimum);
}

Sporadic events are not allowed to come in bursts!

Summary

Scheduling

• Basic real-time scheduling
  • Fixed Priority Scheduling (FPS) with Rate Monotonic (RMPO) Deadline Monotonic Priority Ordering (DMPO)
  • Earliest Deadline First (EDF)

• Real-world extensions
  • Aperiodic, sporadic, soft real-time tasks
  • Deadlines shorter than period
  • Cooperative and deferred pre-emption scheduling
  • Fault tolerance in terms of exception handling considerations
  • Synchronized talks (priority inheritance, priority ceiling protocols)

• Language support
  • Ada95, POSIX ↔ static, off-line analysis mostly — RT-Java ↔ on-line, dynamic scheduling

Resource control

Uwe R. Zimmer – The Australian National University
References for this chapter

[Ada95RM] (link to on-line version)
Ada Working Group
ISO/IEC JTC1/SC 22/WG 9
Ada 95 Reference Manual
– Language and Standard Libraries
ISO/IEC 8652:1995(E) with COR.1:2000,
June 2001

[Ada95RM] (link to on-line version)
Ada Working Group
ISO/IEC JTC1/SC 22/WG 9
Ada 95 Reference Manual
– Language and Standard Libraries
ISO/IEC 8652:1995(E) with COR.1:2000,
June 2001

[Bloom79] Toby Bloom
Evaluating synchronization mechanisms
Proceedings of the seventh ACM Symposium
on Operating systems principles, 1979

[Bloom79] Toby Bloom
Evaluating synchronization mechanisms
Proceedings of the seventh ACM Symposium
on Operating systems principles, 1979

[Bloom79] Toby Bloom
Evaluating synchronization mechanisms
Proceedings of the seventh ACM Symposium
on Operating systems principles, 1979

[Bloom79] Toby Bloom
Evaluating synchronization mechanisms
Proceedings of the seventh ACM Symposium
on Operating systems principles, 1979

[Burns01]
Alan Burns and Andy Wellings
Real-Time Systems and Programming Languages
Addison Wesley, third edition, 2001

[Burns01]
Alan Burns and Andy Wellings
Real-Time Systems and Programming Languages
Addison Wesley, third edition, 2001

[Burns01]
Alan Burns and Andy Wellings
Real-Time Systems and Programming Languages
Addison Wesley, third edition, 2001

[Burns01]
Alan Burns and Andy Wellings
Real-Time Systems and Programming Languages
Addison Wesley, third edition, 2001

[Burns01]
Alan Burns and Andy Wellings
Real-Time Systems and Programming Languages
Addison Wesley, third edition, 2001

[Burns01]
Alan Burns and Andy Wellings
Real-Time Systems and Programming Languages
Addison Wesley, third edition, 2001

[Mercur97]
Clifford W. Mercer
Operating system resource reservation for real-time
and multimedia applications
Ph.D. thesis CMU-CS-97-155, June 1997, Pittsburgh,
Pennsylvania 15213-3890

[Mercur97]
Clifford W. Mercer
Operating system resource reservation for real-time
and multimedia applications
Ph.D. thesis CMU-CS-97-155, June 1997, Pittsburgh,
Pennsylvania 15213-3890

[Mercur97]
Clifford W. Mercer
Operating system resource reservation for real-time
and multimedia applications
Ph.D. thesis CMU-CS-97-155, June 1997, Pittsburgh,
Pennsylvania 15213-3890

[Mercur97]
Clifford W. Mercer
Operating system resource reservation for real-time
and multimedia applications
Ph.D. thesis CMU-CS-97-155, June 1997, Pittsburgh,
Pennsylvania 15213-3890

[Mercury97]
Clifford W. Mercer
Operating system resource reservation for real-time
and multimedia applications
Ph.D. thesis CMU-CS-97-155, June 1997, Pittsburgh,
Pennsylvania 15213-3890

[Mercury97]
Clifford W. Mercer
Operating system resource reservation for real-time
and multimedia applications
Ph.D. thesis CMU-CS-97-155, June 1997, Pittsburgh,
Pennsylvania 15213-3890

[Mercury97]
Clifford W. Mercer
Operating system resource reservation for real-time
and multimedia applications
Ph.D. thesis CMU-CS-97-155, June 1997, Pittsburgh,
Pennsylvania 15213-3890

[Mercury97]
Clifford W. Mercer
Operating system resource reservation for real-time
and multimedia applications
Ph.D. thesis CMU-CS-97-155, June 1997, Pittsburgh,
Pennsylvania 15213-3890

[Murphy2001]
C. Siva Ram Murthy, G. Manimaran
Resource Management in Real-time Systems
and Networks
MIT Press, Cambridge, Massachusetts,
London, England

[Murphy2001]
C. Siva Ram Murthy, G. Manimaran
Resource Management in Real-time Systems
and Networks
MIT Press, Cambridge, Massachusetts,
London, England

[Murphy2001]
C. Siva Ram Murthy, G. Manimaran
Resource Management in Real-time Systems
and Networks
MIT Press, Cambridge, Massachusetts,
London, England

[Murphy2001]
C. Siva Ram Murthy, G. Manimaran
Resource Management in Real-time Systems
and Networks
MIT Press, Cambridge, Massachusetts,
London, England

[Murphy2001]
C. Siva Ram Murthy, G. Manimaran
Resource Management in Real-time Systems
and Networks
MIT Press, Cambridge, Massachusetts,
London, England

[Murphy2001]
C. Siva Ram Murthy, G. Manimaran
Resource Management in Real-time Systems
and Networks
MIT Press, Cambridge, Massachusetts,
London, England

[Murphy2001]
C. Siva Ram Murthy, G. Manimaran
Resource Management in Real-time Systems
and Networks
MIT Press, Cambridge, Massachusetts,
London, England

[Murphy2001]
C. Siva Ram Murthy, G. Manimaran
Resource Management in Real-time Systems
and Networks
MIT Press, Cambridge, Massachusetts,
London, England

[Murphy2001]
C. Siva Ram Murthy, G. Manimaran
Resource Management in Real-time Systems
and Networks
MIT Press, Cambridge, Massachusetts,
London, England

[Murphy2001]
C. Siva Ram Murthy, G. Manimaran
Resource Management in Real-time Systems
and Networks
MIT Press, Cambridge, Massachusetts,
London, England

[Murphy2001]
C. Siva Ram Murthy, G. Manimaran
Resource Management in Real-time Systems
and Networks
MIT Press, Cambridge, Massachusetts,
London, England

(All references and links are available on the course page)

Resource control

Topics in real-time resource control

… from synchronization primitives and schedulers to resource management:

- Toby Bloom’s evaluation criteria for synchronization primitives
- Resource atomicity, liveliness, and double interaction
- Resource reclaiming (C. Siva Ram Murthy, G. Manimaran)
- Resource reservation schemes (Clifford W. Mercer)

(not covered here: general dead-lock prevention / avoidance / detection / recovery algorithms
☞ operating systems course)

Evaluating synchronization mechanisms

Categorizing resource/service requests
(based on Toby Bloom)

Service requests can be categorized by:

- their type
  (read requests might be treated very differently from update request)
- their time (often: by their order or relative time only)
- their attributes, parameters, and the priority of the calling process
  (this includes timing constraints)
- the synchronization state of the resource
  (states which refer to the synchronisation aspect – including timing constraints)
- the internal state of the resource
  (states which refer to the actual contents and available resources– including timing constraints)

Expressive power
☞ are all (required) forms of synchronization available?
☞ can all timing requirements be expressed?

Ease of use
☞ how error-prone are the constructs?
☞ how easy can basic methods be combined to complex resource control systems?
Evaluating synchronization mechanisms

Accepting or Avoiding?

Requests which cannot be fulfilled right now, can be handled via

**Conditional wait**
- accept all calls and suspend the threads internally
  - all threads are immediately inside the synchronized server
  - client threads are released from the server, only when the request is completed (can be overcome)

**Avoidance synchronisation**
- suspend tasks on the level of guards
  - all threads are 'at the borders' of the synchronized server
  - threads can easily revoke their requests

Handling resource requests

**Required features:**
- Handling request types by priorities ✔ (Ada95, Occam2)
- Handling threads by priorities ✔ (most rt-systems)
- Handling threads in order or by their timing constraints ✔ (Real-time Java)
- Handling requests by client-attributes ✍ (mostly: call needs to be accepted first)
- Handling requests by server state ✔ (Ada95, Occam2)

``` vandalism
WHILE TRUE
PRI ALT
ALT i=0 FOR max update [i] ? object
ALT j=0 FOR max modify [j] ? object
```
Evaluating synchronization mechanisms

Handling requests by parameters

protected body resource_control is
  entry allocate(size : instances_of_resource)
  when resources_free < size is
  begin
    resource_free := resource_free - size;
    end allocate;
  procedure free(size : instances_of_resource) is
  begin
    resource_free := resource_free + size;
    end free;
end resource_control;

☞ ‘SR’ [Andrews and Olsson 1993] allows for such a direct access
☞ in most other synchronization environments: accept all and then conditional wait or requeue

package Resource_Manager is
  Max_Resources : constant Integer := 100;
  type Resource_Range is new Integer range 1..Max_Resources;
  subtype Instances_Of_Resource is Resource_Range range 1..50;
  procedure Allocate (Size : Instances_Of_Resource);
  procedure Free (Size : Instances_Of_Resource);
end Resource_Manager;
package body Resource_Manager is
  task Manager is
    entry Sign_In  (Size : Instances_Of_Resource);
    entry Allocate (Instances_Of_Resource);
    entry Free     (Size : Instances_Of_Resource);
  end Manager;
  procedure Allocate (Size : Instances_Of_Resource) is begin
    Manager.Sign_In  (Size);
    Manager.Allocate (Size);
    end Allocate;
  procedure Free (Size : Instances_Of_Resource) is begin
    Manager.Free     (Size);
    end Free;
end Resource_Manager;

☞ double interaction is hidden
☞ Manager can apply any policy to accept the ‘Allocate’ entries
☞ Manager is informed about the request attributes first
☞ entry family

Handling requests by parameters (using wrappers)

Lack of expressive power (e.g. in Ada95) may lead to:
☞ Double Interactions
  e.g. register all requests first, then serve the individual types in a global order
  e.g. announce the parameters first, then serve the individual types based in parameters
  ☞ Requests are no longer atomic!
  ☞ Server deadlocked,
     when wrongly assuming the client is going to make the second call
  ☞ Client deadlocked,
     when wrongly assuming that the client died and is not going to make the second call
Handling requests by types, attributes, and in a global order

Lack of expressive power (e.g. in Ada95) may lead to:

☞ Double Interactions

Ways out:

- Define the double interaction by means of atomic actions and make this known to the underlying synchronization methods.
- Assume that the client will never die during a double interaction sequence.
- Eliminate the double interaction by means of a attributed, single request type and requeuing.

Server has full control over the types, parameters, and orders.

The clients are providing all information.

The protected object is arranging the suspending queues accordingly (requeue-facility).
Handling requests by types, attributes, and in a global order

```haskell
type Request_Kinds is (Allocate_Req, Expand_Req, Free_Req);
type Resource_Range is (small, medium, large);
protected Resource_Control is
  entry Resource_Request (Kind : Request_Kinds; Amount : Resource_Range);
private
  entry Allocate_Sign_In  (Amount : Resource_Range);
  entry Allocate          (Resource_Range_Groups);
  entry Expand_Sign_In    (Amount : Resource_Range);
  entry Expand            (Resource_Range_Groups);
  entry Free              (Amount : Resource_Range);
end Resource_Control;

Server has full control over the types, parameters, and orders
```

Handling requests by types and in a global order

**requeue with abort**

With a standard `requeue` statement:
- any outstanding timeout is cancelled
- the thread is not longer abortable
- clients losing control stemming from an ATC statement, or a timed entry-call
- the server can rely on the client thread no being revoked.

With a `requeue with abort` statement:
- all timeouts are maintained
- allows the client to still revoke the call
- maintains client side control

**requeue can also lead to external entries!**

Reacts need to be considered carefully

**Motivation for resource reclaiming [Murthy2001]**

1. Worst case assumptions give schedulable systems, but might leave only a few spare resources.
2. Resources might not be actually used at run-time.
3. Some aspects of reliability in a real-time system rely directly on the amount of spare resources.

Resource reclaiming may enhance the system’s reliability

**Categorizing resource/service requests**

(based on Toby Bloom)

Service requests can be categorized by:
- their **type**
- their **time** (often: by their order or relative time only)
- their **attributes**, parameters, and the **priority** of the calling process
- the **synchronization state** of the resource
- the **internal state** of the resource

The real-time perspective:
- **take special care of failing tasks (atomic actions, deadlocks)**
- **determine and handle timing constraints in resource requests**
Resource reclaiming properties

- **Correctness:**
  - maintain the feasibility!

- **Inexpensiveness:**
  - resource reclaiming overhead need to be small in comparison to the possible gains

- **Bounded complexity:**
  - resource reclaiming should be included in the task's worst case computation time 
    \( \Rightarrow \) complexity needs to be bound by a constant

- **Effectiveness:**
  - improve the system's actual reliability, 
    thus e.g. more failures can be handled by applying resource reclaiming

More terminology

- **Feasible (prerun) schedule** \( S \):
  - taking into account timing, resource, precedence constraints, and worst case computation times.

- **Postrun schedule** \( S' \):
  - starting from \( S \) and considering the actual computation times into account.

- **Start and finish times:**
  - the scheduled start \( st_i \) and finish times \( ft_i \) as from the feasible prerun schedule \( S \),
    and the actual start \( st'_i \) and finish times \( ft'_i \) as depicted in the postrun schedule \( S' \) of the task \( t_i \).

- **Correct postrun schedule:**
  - a postrun schedule is considered correct iff \( \forall t_i \in Q : (st'_i \leq st_i) \land (ft'_i \leq ft_i) \).

- **Passing tasks:**
  - a task \( t_i \) passed a task \( t_j \) iff \( (st'_i < st_j) \land (ft'_i < ft_j), \) i.e. the strict order in \( S \) is not maintained.

Resource reclaiming algorithms

Two extreme versions:

- Dispatching according to the feasible prerun schedule \( S \),
  - i.e. no reclaiming at all – resource reclaiming cost is zero.

- Global re-scheduling, whenever reclaiming is requested, or at each release of a resource,
  - i.e. optimal reclaiming
    - can be applied only, if the reclaiming cost is smaller than the gained resources

Optimal scheduling of dynamically arriving non-pre-emptive tasks on a multi-processor environment \( \Rightarrow \) NP-hard

- all practical re-scheduling algorithms are approximating. The come in two classes:
  - Algorithms without passing \( \Rightarrow \) bounded complexity
  - Algorithms with passing \( \Rightarrow \) in general: \( O(\log n) \), but bounded with restricted passing
Resource reclaiming from independent tasks

☞ trivial: apply a greedy strategy, which dispatches tasks, whenever there are runnable tasks.

Prerun schedule $S$

- Feasible prerun schedule $S$

Postrun schedule $S'$

- Postrun schedule $S'$ without resource reclaiming

Resource reclaiming from interdependent tasks

☞ greedy reclaiming

Postrun schedule $S'$

- Postrun schedule $S'$ without resource reclaiming
- Tasks $t_4 \otimes t_7; t_4 \otimes t_9; t_7 \otimes t_{12}; t_9 \otimes t_{12}$ have conflicting resource locks
Resource Reclaiming [Murthy2001]

Resource reclaiming from interdependent tasks

☞ greedy reclaiming

Runtime anomaly

- Postrun schedule $S'$ without greedy resource reclaiming
- Tasks $t_4 \otimes t_7; t_4 \otimes t_9; t_7 \otimes t_9; t_{12}; t_9 \otimes t_{12}$ have conflicting resource requests

☞ basic reclaiming: look for simultaneous idling

Postrun schedule $S'$

- Postrun schedule $S'$ without basic resource reclaiming
- Tasks $t_4 \otimes t_7; t_4 \otimes t_9; t_7 \otimes t_9; t_{12}; t_9 \otimes t_{12}$ have conflicting resource locks

Resource Reclaiming from interdependent tasks

☞ early start algorithm

- Detect overlaps in the prerun schedule $S$:
  
  $t_{<j} = \{ t_j | t_j < st_j \}$
  
  $t_{>j} = \{ t_j | t_j > st_j \}$
  
  $t_{<j} \cap \{ t_j \notin t_{<j} \}$

- Detect tasks overlapping with $t_j$ on processor $k$ and order all sets

- Allow tasks in $t_{<j}$ to be executed simultaneously and ensure that they do not overlap with tasks out of $t_{<j}$ or $t_{>j}$.

- Complexity $O(m^2)$; with $m$ processors.
Resource reclaiming from interdependent tasks

**Resource Reclaiming** [Murthy2001]

- Early start algorithm

Postrun schedule $S'$

- Postrun schedule $S'$ without resource reclaiming
- Tasks $t_4 \otimes t_7; t_4 \otimes t_9; t_7 \otimes t_9; t_7 \otimes t_{12}; t_9 \otimes t_{12}$ have conflicting resource locks

Restriction vectors (RV) algorithm (Shen, Ramamritham, Stankovic, '93)

- Restriction vector (RV):

$$RV[j][t] = \begin{cases} t_k \in t_{c_i(j)}[\neg\exists t \in t_{c_i(j)}(st_t > st_k)] & \text{if } j= \text{proc}(i) \\ t_m \in t_{c_i(j)}([t_m < t_j \lor t_m \otimes t_j]) \land (\neg\exists t \in t_{c_i(j)}(st_t > st_m) \land (t_j < t_m \lor t_j \otimes t_m)) & \text{if } j \neq \text{proc}(i) \\ - & \text{no such task} \end{cases}$$

- Completion bit matrix (CBM):

$$\text{CMB}[i][j] = \begin{cases} 1 & \text{if task } t_i \text{ has completed its scheduled execution in processor } j \\ 0 & \text{otherwise} \end{cases}$$

compute the $RV(j)$ by checking the $k$ most recent tasks in $t_{c_i(j)}$

for any task $t_j$ next to be scheduled on processor $j$:

- fetch the most recent CBM
- if $\forall t_i \in RV(j) \left[ \text{CMB}(i,j) = 1 \right]$ then start $t_j$ else idle until the next CBM update.

The algorithm is heuristic in the sense that it is only checking the $k$ most recent tasks in $t_{c_i(j)}$

The complexity is $O(m^2)$ since $m$ processors need to check $m$ RV-entries = bounded
Resource reclaiming from interdependent tasks

- Restriction vectors (RV) algorithm (Shen, Ramamritham, Stankovic, '93)

**Prerun schedule S**

- Tasks \( t_4 \otimes t_2; t_4 \otimes t_9; t_3 \otimes t_5; t_4 \otimes t_1; t_9 \otimes t_12 \) have conflicting resource locks.

- Tasks \( t_10 < t_2; t_10 < t_4; t_5 < t_6; t_8 < t_13; t_1 < t_2; t_1 < t_3; t_2 < t_12; t_3 < t_12; t_11 < t_12 \) have precedence relations.

**Postrun schedule S'**

- RVs: \( t_1: [\cdot, \cdot, \cdot]; \quad t_2: [t_1, \cdot, \cdot]; \quad t_3: [t_2, \cdot, \cdot]; \quad t_4: [t_3, t_4, t_11]; \quad t_5: [t_4, t_9, t_10]; \quad t_6: [t_5, \cdot, \cdot]; \quad t_7: [t_6, \cdot, \cdot]; \quad t_8: [t_7, t_10]; \quad t_9: [t_8, t_12]; \quad t_{10}: [t_9, \cdot, \cdot]; \quad t_{11}: [\cdot, \cdot, t_10]; \quad t_{12}: [t_3, t_7, t_11]; \quad t_{13}: [\cdot, t_6, \cdot] \)

- RVs: \( t_1: [\cdot, \cdot, \cdot]; \quad t_2: [t_1, \cdot, \cdot]; \quad t_3: [t_2, \cdot, \cdot]; \quad t_4: [t_3, t_4, t_11]; \quad t_5: [t_4, t_9, t_10]; \quad t_6: [t_5, \cdot, \cdot]; \quad t_7: [t_6, \cdot, \cdot]; \quad t_8: [t_7, t_10]; \quad t_9: [t_8, t_12]; \quad t_{10}: [t_9, \cdot, \cdot]; \quad t_{11}: [\cdot, \cdot, t_10]; \quad t_{12}: [t_3, t_7, t_11]; \quad t_{13}: [\cdot, t_6, \cdot] \)
Resource reclaiming from interdependent tasks

Restriction vectors (RV) algorithm (Shen, Ramamritham, Stankovic, ‘93)

Proof of Correctness

Lemma: Given a feasible prerun schedule $S$: if $\exists t_i [s_{t_i'} > s_{t_i}]$ then passing must have occurred.

Proof: Assuming that no passing occurred, then all $t \in t_{\prec_j}$ have been dispatched before $t_j$ and all $t \in t_{\succ_j}$ are only dispatched after $t_j$ completed. By definition of a feasible schedule all $t \in t_{\prec_j}$ do not interfere with $t_j$ and can thus by no means delay the execution of $t_j$.

Therefore $s_{t_i'} \leq s_{t_i}$.

Restriction vector (RV) algorithm with task migration (Manimaran, Murthy, ‘97)

- Restriction vector (RV) with static processor assignment:

  $$RV_{ij} = \begin{cases} 
  t_k \in t_{\prec_j}(j) \land \neg \exists t_l \in t_{\prec_j}(j) [s_{t_l} > s_{t_i}] & \text{if } j = \text{proc}(i) \\
  t_m \in t_{\prec_j}(j) \land (t_m < t_j \land t_m \land t_j) \land \neg \exists t_l \in t_{\prec_j}(j) [s_{t_l} > s_{t_m}] \land (t_j < t_i \land t_j \land t_i) & \text{if } j \neq \text{proc}(i) \\
  \text{no such task} & 
  \end{cases}$$

- Restriction vector (RV) with dynamic processor assignment:

  $$RV_{ij} = \begin{cases} 
  t_m \in t_{\prec_j}(j) \land (t_m < t_j \land t_m \land t_j) \land \neg \exists t_l \in t_{\prec_j}(j) [s_{t_l} > s_{t_m}] \land (t_j < t_i \land t_j \land t_i) & \text{if } t_m \text{ exists} \\
  \text{no such task} & 
  \end{cases}$$

Prerun schedule $S$

- RVs: $t_1: [\ldots, \ldots, \ldots]$; $t_4: [\ldots, \ldots, \ldots]$; $t_7: [\ldots, \ldots, \ldots]$; $t_{10}: [\ldots, \ldots, \ldots]$; $t_{13}: [\ldots, \ldots, \ldots]$
Resource reclaiming from interdependent tasks

Restriction vectors (RV) algorithm with task migration (Manimaran, Murthy, ‘97)

Prerun schedule $S$

Postrun schedule $S'$ with task migration

Restriction vectors (RV) algorithm with task migration (Manimaran, Murthy, ‘97)

Resource Reclaiming [Murthy2001]
Resource reclaiming from interdependent tasks

Restriction vectors (RV) algorithm with task migration (Manimaran, Murthy, '97)

Postrun schedule $S'$ with task migration

$RVs$: $t_1$: $[5, 5, \ldots]$; $t_2$: $[\ldots, 5, \ldots]$; $t_3$: $[\ldots, 5, \ldots]$; $t_4$: $[\ldots, 5, t_{10}]$; $t_5$: $[\ldots, \ldots]$; $t_6$: $[\ldots, \ldots]$; $t_7$: $[\ldots, \ldots]$; $t_8$: $[\ldots, \ldots, t_{10}]$; $t_9$: $[\ldots, t_8, t_{12}]$; $t_{10}$: $[\ldots, \ldots]$;
$t_{11}$: $[\ldots, \ldots]$; $t_{12}$: $[t_3, t_7, t_{11}]$; $t_{13}$: $[t_6, \ldots]$

Resource reclaiming from interdependent tasks

Restriction vectors (RV) algorithm with task migration (Manimaran, Murthy, '97)

Postrun schedule $S'$ with task migration

$RVs$: $t_1$: $[5, 5, \ldots]$; $t_2$: $[\ldots, 5, \ldots]$; $t_3$: $[\ldots, 5, \ldots]$; $t_4$: $[\ldots, 5, t_{10}]$; $t_5$: $[\ldots, \ldots]$; $t_6$: $[\ldots, \ldots]$; $t_7$: $[\ldots, \ldots]$; $t_8$: $[\ldots, \ldots, t_{10}]$; $t_9$: $[\ldots, t_8, t_{12}]$; $t_{10}$: $[\ldots, \ldots]$;
$t_{11}$: $[\ldots, \ldots]$; $t_{12}$: $[t_3, t_7, t_{11}]$; $t_{13}$: $[t_6, \ldots]$

Resource Reclaiming [Murthy2001]

Resource Reclaiming [Murthy2001]
Resource reclaiming from interdependent tasks

- Restriction vectors (RV) algorithm with task migration (Manimaran, Murthy, ‘97)

Correctness of the migration process

To ensure that the swapping of dispatching queues $DQ_x \leftrightarrow DQ_y$ between processor $P_x$ and $P_y$ does not interfere with the correctness of the postrun schedule $S'$, swapping is permitted only if:

$$st_i \geq ft_j$$

where task $t_i$ is not further delayed (where task $t_j$ is next to be scheduled on the idling $P_x$ and task $t_j$ is currently executing on $P_y$).

The unrestricted and executable task $t_k$, which is next to be scheduled on $P_y$ is started earlier by transferring it to the idling $P_x$.

- No task is delayed by swapping these dispatching queues.

Resource reclaiming evaluated

Some additional observables:

- Task graph density $\rho \rightarrow [0\ldots1]$, where $\rho$ indicates tasks, and independent and one a fully dependent task-set.
- $aw$-ratio: $C_i' / C_i$ (actual to worst case ratio)
- mig-attempts: number of checks on dispatch queues by the RV with migration algorithm

RC computational costs (from Manimaran, Murthy, Vijay, Ramamritham ‘97):

$$C_{RC-basic} = 1$$

$$C_{RC-early-start} = mC_{RC-basic}$$

$$C_{RC-RV} = C_{RC-early-start} + C_{RV}$$

with $C_{RV}$ the cost for the calculation of the RVs.

$$C_{RC-RV-migration} = C_{RC-RV} + f(mig-attempts, C_{RC-early-start})$$

Practical measurements:

- There is a continuous improvement in terms of gained resources by applying: basic $\rightarrow$ early-start $\rightarrow$ RV-reclaiming $\rightarrow$ RV-reclaiming-with-task-migration-algorithms.
- In case of RV reclaiming with task migration, the extended communication/synchronization overhead can reach noticeable levels.
- There need to be a high degree of dependencies in the task-set ($P_n$), in order to justify the application of RV reclaiming with task migration.

Reclaiming in the introduced sense is applicable only to real-time systems which:

- allow for earlier task start times
- allow for task migration
- and where all dependencies can be expressed in terms of the introduced formalism
Summary

Resource control

- Resource synchronization primitives
  - evaluation criteria for resource synchronization methods
  - atomicity, liveliness, and double interaction
- Resource reclaiming schemes
  - basic reclaiming, early start, and restriction vector algorithms
  - resource reclaiming with task migration
- Real-time resource control
  - policy and run-time issues to be considered

References for this chapter

- [Burns98] Alan Burns, Brian Dobbing, George Romanski
  The Ravenscar Tasking Profile for High Integrity Real-Time Programs
- [Burns01] Alan Burns and Andy Wellings
  Real-Time Systems and Programming Languages
  Addison Wesley, third edition, 2001
  Assuring Design Diversity in N-Version Software:
  A Design Paradigm for N-Version Programming
  IEEE Computer Society Press Technology Series, October 1992, pp. 45-54
  Axioms for Real-Time Logics
  Lecture Notes in Computer Science 1466, Springer-Verlag, 1999, pp. 219-236

Reliability

Based on a set of powerful and diverse tools …

… reconsidering the basic problems of:

- system identification / analysis
- fault prevention
- error detection
- fault tolerance

… building predictable / dependable systems …

... in a real-time domain!
Reliability

**Terminology**

Reliability ::= measure of success with which a system conforms to its specification or low failure rate.

Failure ::= deviation of a system from its specification

Error ::= system state which lead to failures

Fault ::= the reason for an error

---

**Faults on different levels**

- Inconsistent or inadequate specification
  - very frequent source for disastrous faults
- Software design errors
  - very frequent source for disastrous faults
- Component & communication system failures
  - rare and mostly predictable

---

**Faults in the time domain**

- Transient faults
  - many communication system failures, electric interference, etc.
- Intermittent faults
  - transient errors which occur more than once (e.g. overheating effects)
- Permanent faults
  - stay in the system until they are repaired by some means

---

**Observable failures states**

![Diagram of observable failure states]

- Time domain
  - fail never
  - fail uncontrolled
- Value domain
  - Constraint error
  - Value error
  - fail silent
  - fail stop
  - fail controlled
Achieving reliability

Fault avoidance

Fault avoidance at hardware-level:
- use reliable hardware components — consider the environmental demands!
- use an adequate (hardware) system design — shock, humidity, interference, ...
- ensure proper assembly and encapsulation — weak connectors, bad connectors, ...

Fault avoidance at software design level:
- strict system specifications (employ format methods if applicable)
- use proven software-engineering and design methodologies
- employ languages and run-time environments with reasonable support for the requirements.

Fault removal

Find and remove errors from the previous stage.
☞ Team programming methods like extreme programming or rigorous testing may help here.

but ...
- no re-evaluation method indicates the absence of faults (even formal methods cannot identify specification faults)

... and specifically for real-time and embedded systems ...:
- often: tests cannot be performed under realistic conditions ... especially exceptional conditions
- most simulation environments have a severe impact on real-time systems
- the test space for real-time system is significantly larger than for non-real-time systems

Investigate:
- static applications specifications
- physical sensors and converters constraints
- constraints of the employed controller network
- constraints of the underlying run-time system
- dynamic application specifications (requested real-time behaviour)

☞ Understanding all critical real-time requirements and issues

Fault avoidance at hardware-level:
- use reliable hardware components — consider the environmental demands!
- use an adequate (hardware) system design — shock, humidity, interference, ...
- ensure proper assembly and encapsulation — weak connectors, bad connectors, ...

Fault avoidance at software design level:
- strict system specifications (employ format methods if applicable)
- use proven software-engineering and design methodologies
- employ languages and run-time environments with reasonable support for the requirements.
Reliability

Fault prevention
(avoidance & removal)

Regardless of the rigor of fault prevention methods:

the real-time system might still fail

This is specifically critical for all non-monitored systems:

- systems which are (temporary) inaccessible
- un-manned vehicles which operate autonomously by default
- systems in remote / dangerous environments

Instead (or in addition to fault prevention): enabling a ‘safe landing’:

Fault tolerance

- Full fault tolerance
  the system continues to operate in the presence of ‘foreseeable’ error conditions without any significant failures — also this might induct a reduced operation period.
- Graceful degradation (fail soft)
  the system continues to operate in the presence of ‘foreseeable’ error conditions, accepting a partial loss of functionality or performance.
- Fail safe
  the system halts and maintains its integrity

Full fault tolerance is not maintainable for an infinite operation time!
Graceful degradation might have multiple levels of reduced functionality.

Fault tolerance

Hardware redundancy

- adding extra hardware resources:
  - for the detection of failures and the localization of faults
  - for the handling of exceptional situations and error-recovery.
  - as a functional duplication or multiplication of complete (sub-)systems in order to hot-swap or select the operational one in case of a failure in one part of the (sub-)system.
- Fault-detection and recovery hardware includes:
  - watch-dog timers, limit switches, additional physical sensors, transient-recording-systems (emergency system dump), overload-backup-systems, or even in-circuit emulators.
- Triple Modular Redundancy (TMR) or N-Modular Redundancy (NMR) assumes:
  - functionally identical components which are either:
    - static parts of the system and connected via a voting/masking/comparing system
    - or in case of a detected error-condition: dynamic parts which are swapped in.
  - the fault is based on a physical phenomenon, which applies only locally
  - or the structure of the functionally identical systems is sufficiently different

For some high-risk systems this approach is applied in forms of redundant sub-systems with:

- the same specification
- different computer systems (CPUs, buses, memory systems, drives)
- different operating systems
- different real-time languages and development environments (N-Version programming)
- and by restricting the communication between the different developer teams

not too surprisingly, the outputs from the different systems are slightly different …
Real-Time & Embedded Systems

Reliability

**Triple Modular Redundancy**
*(example)*

3 identical primary flight computers distributed in the Boeing 777, each consisting of:

- 3 processors: AMD 29050, Motorola 68040, INTEL 80486 (called ‘lanes’)
- independent power-sources and inertia measurements
- code build by 3 different Ada compilers
- the same Ada source code (‘the specification’): around 3 million lines of code, but different monitor functions

Targeted failure probability: \(< 10^{-10}/h\) (e.g. UK Seizewell B nuclear reactor (emerg.): \(< 10^{-3}/h\))

No single fault on board the 777 should occur without failure identification.
No single fault on board the 777 should cause more than the loss of one primary flight computer.

Sophisticated synchronization and communication systems.

(not a single fatal event — information from November 2001)

Real-Time & Embedded Systems

Reliability

**N-version programming**

Impacts to software diversity:

<table>
<thead>
<tr>
<th>Development teams</th>
<th>Languages</th>
<th>Tools</th>
<th>Algorithms</th>
<th>Methodologies</th>
</tr>
</thead>
<tbody>
<tr>
<td>Specification</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>Design</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>Coding</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>Testing</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

(■: highest – ■: high – ■: low – ■: lowest) (source: [Lyu92])

Real-Time & Embedded Systems

Reliability

**“The six-language project”**

Joint project between the UCLA (Dependable computing and fault-tolerance systems) and the Honeywell Commercial Flight Systems Division (1992)

- The specifications (about a flight controller) were original system description documents (SDD) by Honeywell enhanced by additional cross-checking points and included some enforced diversity elements (a 64-page document).
- The development teams were isolated and any technical discussions were strictly prohibited
- All communication and documentation is requested to follow predefined protocols (written form) defined and handled by a coordinating team.
- Specified tests were performed by the coordinating team before a version was accepted for integration.
- The N-Version paradigm was applied to all stages of the development cycle.

**Language** | **Sources (l.o.c.)** | **Test Runs** | **Errors** | **Failure-rate**
---|---|---|---|---
Ada | 2256 | 5127400 | 0 | 0
‘C’ | 1531 | — | 568 | 1.108×10⁻⁴
Modula-2 | 1562 | — | 0 | 0
Pascal | 2331 | — | 0 | 0
Prolog | 2228 | — | 680 | 1.326×10⁻⁴
T (close to Lisp) | 1568 | — | 680 | 1.326×10⁻⁴
Average | 1913 | — | 321 | 0.627×10⁻⁴
Reliability

“The six-language project”
(source: [Lyu92])

<table>
<thead>
<tr>
<th>Failure category</th>
<th>average single version failure probabilities (5127400 cases)</th>
<th>average 3-version failure probabilities (102548000 cases)</th>
<th>average 5-version failure probabilities (30764400 cases)</th>
</tr>
</thead>
<tbody>
<tr>
<td>no errors</td>
<td>0.99993733</td>
<td>0.9998409</td>
<td>0.9997807</td>
</tr>
<tr>
<td>single error</td>
<td>6.27x10^{-5}</td>
<td>13.05x10^{-5}</td>
<td>19.15x10^{-5}</td>
</tr>
<tr>
<td>two distinct errors</td>
<td>0.20x10^{-5}</td>
<td>0.23x10^{-5}</td>
<td></td>
</tr>
<tr>
<td>two coincident errors</td>
<td>2.65x10^{-5}</td>
<td>2.21x10^{-5}</td>
<td></td>
</tr>
<tr>
<td>three errors</td>
<td></td>
<td>0.34x10^{-5}</td>
<td></td>
</tr>
</tbody>
</table>

*The resulting 3-version and 5-version systems displayed lower failure rates than a ‘golden master’ reference implementation by Honeywell.*

*Coincident errors involving more than two versions were never observed.*

*a total of 93 faults were detected.*

*control problems are specifically suitable for n-version programming, since the error-detection and synchronization algorithms are relatively simple.*

*in general: diverting results do not necessarily imply any faults.*

N-version programming – Voting issues

• **Integer arithmetic:**
  
  • Integer (or any discrete sub-type) -based results will be identical → ✔

• **Real arithmetic:**
  
  • Real-valued results will usually be different → Comparisons need to consider tolerances.
  
  • If the process is not fully continuous (thresholds, quantizations, bifurcations)
    
    → Comparisons need to re-model the whole process in order to evaluate similarities
  
  • Independence ✔ re-specify the system

• **Multiple solutions:**
  
  • The solution space itself allows for multiple correct, but different solutions
    
    → ✔ re-specify the system

some issues:

• **Specification:**
  
  Assuming that a good part of software faults stem from wrong or incomplete specifications

  ⇨ N-version programming will not help in this case

• **Diversity assumption:**
  
  Diversity can be enforced and supported in some areas (demonstrated by examples), while coincident error conditions can be observed in other application domains (also documented by case-studies). The rigorous identification of adequate domains for N-version programming is currently part of active research.

• **Project costs:**
  
  Since the development costs are increasing by a factor of N plus coordination costs, it needs to be considered carefully whether a single version developed with the same effort shows perhaps a similar level of reliability.
Four constituent phases (Anderson and Lee, ‘90):

1. **Error detection**
   Detection of a precise error state is essential.

2. **Damage confinement and assessment**
   Diagnosis of the damage, which occurred between the fault and the detected error state.

3. **Error recovery**
   Sequence of operations leading from the detected error state to an operational state.

4. **Fault treatment**
   In order to prevent the same error state again, the fault itself might/should be eliminated.

---

**Dynamic redundancy — Damage diagnosis**

**Confinement:**
- How to avoid the transfer of fault-effects between system parts?
  - Modular decomposition
  - Atomic actions
  - ‘Firewalls’

**Assessment:**
- resulting from the location of the detected error state and the possible paths though the system which are all leading to this error state.
- a fine-granular system structure (error-confinement) limits the length of these possible paths.
- a very well structured system is the cornerstone of damage diagnosis.

**Dynamic redundancy — Error detection**

- Error states from the environment
  - **Hardware** ... CPU, controllers, communication systems, ...
  - **Run-time environment**
- Error states stemming from checks without the application processes
  - **Replication** – employ N-version programming to detect error states
  - **Timing** – watchdog timers and overrun detectors
  - **Reversal** – apply the reverse function and compare \( x \oplus f^{-1}(f(x)) \)
  - **Coding** – detect corrupted data via redundant information (CRC-checks, ...)
  - **Reasonableness** – check assertions (e.g. in Eiffel)
  - **Structural** – check structural integrity (e.g. lists, file-systems)
  - **Continuity** – assuming a limited difference between consecutive controller values.

**Dynamic redundancy — Error recovery**

**Backward error recovery:**
- set checkpoints and safe the system state with each passing of a checkpoint. how can system-wide consistent checkpoints be ensured?
- if an error state is detected: set back to the last consistent checkpoint.
  - applicable even if the fault itself can not be identified.
  - not applicable at all, if the system contains non-reversible or -resetable components (time, ...)

**Forward error recovery:**
- method of choice for most time critical parts of real-time and embedded systems.
  - highly application dependent.
  - may involve complex mode and priority changes (deadlines might be still relevant).
Dynamic redundancy — Fault treatment

- Localization of a hardware fault is usually easier and more precise than of a software fault.
- On-line fault treatment might be tricky and is usually limited to (hot) exchanges of complete modules (software as well as hardware).
- Granularity is usually finer than in static redundant systems.
- Exchange of faulty components is nevertheless usually an expensive and complex operation.

☞ the number of substitutable sub-systems in a dynamic redundant system is still limited.

(many systems will assume transient faults, log the event and continue operations ... )

Safety and Dependability

- **Safety**: freedom from those conditions that can cause death, injury, occupational illness, damage to (or loss of) equipment (or property), or environmental harm (Leveson, '86)

☞ are there any safe and functional systems beyond a certain complexity? ... aeroplanes? cars?

- **Dependability**:
  - Availability — ready to use
  - Reliability — absence of failures
  - Safety — absence of fatal failures
  - Confidentiality — absence of unauthorized disclosures
  - Integrity — no data corruptions
  - Maintainability — accessibility to changes and improvements

... more reliability in the design process:

Restrict, Formalise, ...?

Restrict:
- limit the tools and environments to 'safer' operations
- e.g. Esterel, High-Intensity Pearl, Ada95 Ravenscar profile, ...

Formalise:
- UML ('the object oriented approach')
- Temporal logic, Real-Time Logic (RTL) as an extension of predicate logic
- classical real-time design methods: MASCOT, JSD, MOON, HOOD, HRT-HOOD, CODARTS, ...

Ada95 Ravenscar profile (Burns, Dobbing, Romanski '98)

- **Task type and object declarations at the library level**
  - no hierarchy of tasks, and hence no exit protocols needed from blocks and sub-programs.
- **No dynamic allocation or unchecked de-allocation of protected and task objects**
  - removes the need for dynamic objects.
- **Tasks are assumed to be non-terminating**
  - this is primarily because task termination is generally considered to be an error for a real-time program which is long-running and defines all of its tasks at start-up.
- **Library level Protected objects with no entries**
  - these provide atomic updates to shared data and can be implemented simply.
- **Library level Protected objects with a single entry**
  - used for invocation signalling; but removes the overheads of a complicated exit protocol.
Real-Time & Embedded Systems

Reliability

Ada95 Ravenscar profile (Burns, Dobbing, Romanski ’98)

- Barrier consisting of a single boolean variable — no side effects are possible and exit protocol becomes simple.
- Only a single task may queue on an entry — hence no queue required; this is a static property that can easily be verified, or it can lead to a bounded error at runtime.
- No requeue — leads to complicated protocols, significant overheads and is difficult to analyse (both functionally and temporally).
- No Abort or ATC — these features lead to the greatest overhead in the run-time system due to the need to protect data structures against asynchronous task actions.
- No use of the select statement — non-deterministic behaviour is difficult to analyse, moreover the existence of protected objects has diminished the importance of the select statement to the tasking model.
- No use of task entries — not necessary to program systems that can be analysed; it follows that there is no need for the accept statement.

- “Delay until” statement but no “delay” statement — the absolute form of delay is the correct one to use for constructing periodic tasks.
- “Real-Time” package — to gain access to the real-time clock.
- No Calendar package — “Real-Time” package is sufficient.
- Atomic and Volatile pragmas — needed to enforce the correct use of shared data.
- Count attribute (but not within entry barriers) — can be useful for some algorithms and has low overhead.
- Ada.Task_Identification — can be useful for some algorithms and has low overhead, available in reduced form (no Abort_Task or task attribute functions Callable or Terminated).
- Task discriminants — can be useful for some algorithms and has low overhead.
- No user-defined task attributes — introduces a dynamic feature into the run-time that has complexity and overhead.

- No use of dynamic priorities — ensures that the priority assigned at task creation is unchanged during the task’s execution, except when the task is executing a protected operation.
- Protected procedures as interrupt handlers — required if interrupts are to be handled.

☞ commercially available and employed in hard-real-time systems

Temporal logic

- Extending predicate logic
- Adding the concepts of ordering for events and states
- Suitable for event driven system, reactive systems ☞ Esterel

© 2009 Uwe R. Zimmer, The Australian National University
Real-Time & Embedded Systems

Reliability

Temporal logic

• Assertions on sequences and orders of states
  ☞ employ predicate logic & a set of new operators:

  • A: A is true for all future states
  • A: A is eventually true
  • A: A is true for the following state

  e.g. ●((Collision_Warning ⇒ ◊ Collision_Avoidance)
    or: ●((Collision_Warning ⇒ ◇ Collision_Avoidance)

assuming that there is a sequence of distinguishable states (or ‘time’) S.

• Another temporal operator:
  AμB: A holds until the first occurrence of B, which will occur eventually.

  e.g. ●((Tasks_Waiting μ Entry_Closed) ∧ (~Tasks_Waiting μ Entry_Open))

  ⇒ Temporal logic expresses the order of events only
  and has means to express temporal scopes, deadlines, ...

Real-Time & Embedded Systems

Reliability

Real-Time Logic

• Assertions on real-time events:
  ☞ employ predicate logic & a occurrence function:

    • @E(i): denotes the time of the i-th occurrence of event (-class) E
    • the event (-class) E is strictly ordered by instance (i) and time (@).
    • all events of kind (class) E can be distinguished.
    • instance order ⇒ order in time.

  Examples:
  ▲E ((@E(i) ≤ @(▲ A, j)) ∧ (@(▼ A, j) ≤ @E(i) + d) ∧ (@(▼ A, j+1) ≤ @E(i) + p))
  ▼E, i, j ((@▼ A, i) < @((▲ B, j)) ∧ (i @((▲ B, j) < @((▲ A, i))

Interpretations:
  does ∀i @(E, i) indicate all possible, all defined, or all observed instances of E?
Reliability

**Linear Temporal Logic of Real Numbers (LTR)**

\[ \phi ::= p \lor \phi_1 \lor \phi_2 \lor \neg \phi_1 \lor \phi_2 \lor \phi_1 \lor \phi_2 \]

where

- \((\tau, t) \models p\) iff \(p \in \tau(t)\)
- \((\tau, t) \models \phi_1 \lor \phi_2\) iff \((\tau, t) \models \phi_1\) or \((\tau, t) \models \phi_2\)
- \((\tau, t) \models \neg \phi\) iff \((\tau, t) \not\models \phi\)
- \((\tau, t) \models \phi_1 \lor \phi_2\) iff \(\exists t' > t \land t' \models \phi_2\) and \(\forall t'' \in (t, t'), t'' \models \phi_1 \lor \phi_2\)
- \((\tau, t) \models \phi_1 \lor \phi_2\) iff \(\exists t' < t \land t' \models \phi_2\) and \(\forall t'' \in (t', t), t'' \models \phi_1 \lor \phi_2\)

\(\phi\) is satisfiable iff \(\exists (\tau, t) \models \phi\) — \(\phi\) is valid iff \(\forall (\tau, t) \models \phi\)

**Event-Clock Temporal Logic**

\[ \phi ::= p \lor \phi_1 \lor \phi_2 \lor \neg \phi_1 \lor \phi_2 \lor \phi_1 \lor \phi_2 < \phi \lor \phi_1 \lor \phi_2 \]

where

- \((\tau, t) \models p\) iff \(p \in \tau(t)\)
- \((\tau, t) \models \phi_1 \lor \phi_2\) iff \((\tau, t) \models \phi_1\) or \((\tau, t) \models \phi_2\)
- \((\tau, t) \models \neg \phi\) iff \((\tau, t) \not\models \phi\)
- \((\tau, t) \models \phi_1 \lor \phi_2\) iff \(\exists t' > t \land t' \models \phi_2\) and \(\forall t'' \in (t, t'), t'' \models \phi_1 \lor \phi_2\)
- \((\tau, t) \models \phi_1 \lor \phi_2\) iff \(\exists t' < t \land t' \models \phi_2\) and \(\forall t'' \in (t', t), t'' \models \phi_1 \lor \phi_2\)
- \((\tau, t) \models \phi_1 \lor \phi_2\) iff \(\exists t' < t \land t' \models \phi_2\) and \(\forall t'' \in (t - 1, t), t'' \not\models \phi\)
- \((\tau, t) \models \phi_1 \lor \phi_2\) iff \(\exists t' > t \land t' \models \phi_2\) and \(\forall t'' \in (t + 1, t), t'' \not\models \phi\)

**Summary**

- Terminology
  - Faults, Errors, Failures – Reliability
- Faults
  - Fault avoidance, removal, prevention \(\Rightarrow\) Fault tolerance
- Redundancy
  - Static (TMR, NMR) and dynamic redundancy
  - N-version programming, and dynamic redundancy in software design
- Reduce & Formalise
  - Ada95 Ravenscar profile
  - Real-time Logic